

- Spec ial guests to day:
* Informa tic s students:
- Benji Sc hwa rtz-G ilb ert
- Ryan Musgrave
- Devyn J ones

- Why is BYTE spelled with a Y?
* The Engineers at IBM were looking for a word for a quantity of memory between a bit and a word (usually 32 bits).
- They liked bite but too close to bit
- Typing errors could confuse the two
- Changed the ito a y to make them distinct


Everyone knows computers use bits and bytes ... but what are they?


Information must be in a form that * Humans can understand and * Computers can manipulate

Digitizing bridges the gap


- Digitize: Represent information with digits (noma lly base-10 numerals 0 through 9)
- Limita tion of Digits
* Altemative Representation: Any set of symbols c ould represent phone number digits, as long as the keypad is labeled accordingly

- Symbols, Briefly
* Digits ha ve the adva nta ge of having short names (easy to say)
* But c omputer professiona ls a re shortening symbol names
- Period is "dot"
- Exclamation point is pronounced "bang"



Figure 8.1. Three symbol assignments for a telephone keypad.


- Digits for enc oding info
* Can list items in numerical order
- To use othersymbols, we need an ordering system (collating sequence)
* Agreed order from sma llest to la rgest value
- In choosing symb ols for enc oding, consider how symbols interact with things being encoded

- Analog is continuous data/information * Sound waves


Time


- Digital is disc rete data/information * Many distinct samples of data
* Stored in binary (0's a nd 1's)
- All data in a computer is represented in binary


Time


- The funda mental pattems used in ITcome when the physical world meets the logical world
- The most funda mental form of information is the presence orabsence ofa physical phenomenon
- In the logical world, the concepts of true and false a re importa nt
* Associate true with presence of a phenomenon and false with its absence, we use the physical world to implement the logical world, and produce information technology
- PandA is the mnemonic for "presence and absence"
- It is disc rete (distinct or separable)the phenomenon is present or it is not (true or false; 1 or 0). There in no continuous gradation in between.

- Two pattems make a bina ry system
* Base 2 (0's or 1's)
- The basic binary unit is known as a "bit" (short for bina ry digit)
- 8 bits a re grouped together to form a byte
* Memory accessed by byte addresses
- We can give any na mes to these two pattems as long aswe are consistent
* PandA (Presence and Absence can represent 1 and 0, respectively)


| Present | Absent |
| :--- | :--- |
| 1 | 0 |
| On | Off |
| Yes | No |
| + | - |
| Black | White |
| For | Against |
| Yang | Yin |
| Lisa | Bart |





- Memory is a rranged inside a computer in a very long sequence of bits
* Bits = places where a phenomenon can be set and detected
- Analogy: Sidewalk "Memory"
* Each sidewalk squa re represents a memory slot, or bit, and stones represent the presence orabsence
* If a stone is on the square, the value is 1 , if not the value is 0


Figure 8.2 Sidewalk sections as a sequence of bits (1010 0010).


- Altemate waysto encode two states using physical phenomena
* Use stones on all squa res, but black stones for one state and white for the other
* Use multiple stones of two colors per square, saying more blackthan white means 0 and more white than black means 1
* Stone in center for one state, off-centerfor the other
* etc.

- Since we only have two pattems, we must combine them into sequences to create enough symbols to encode necessary information
- Bina ry (Pa nd A) has 2 pattems, a rranging them into n-length sequences, we can create $2^{n}$ symbols

- Number of symbols when the number of possible pattems is two (0 and 1)

| $n$ | $2^{\text {n }}$ | Symbols |
| :---: | :---: | :---: |
| 1 | $2^{1}$ | 2 |
| 2 | $2^{2}$ | 4 |
| 3 | $2^{3}$ | 8 |
| 4 | $2^{4}$ | 16 |
| 5 | $2^{5}$ | 32 |
| 6 | $2^{6}$ | 61 |
| 7 | $2^{7}$ | 128 |
| 8 | $2^{8}$ | 256 |
| 9 | $2^{9}$ | 512 |
| 10 | $2^{10}$ | 1,024 |



- Rec all in Cha pter 4, we specified custom colors in HTML using hex digits
* e.g., <p bgColor ="\#FF8E2A">
* Hex is short forhexadecimal (base 16)
-Why use hex?
* Writing the sequence of bits is long, tedious, and error-prone

- Sixteen values can be represented perfectly by 4-bit sequences ( $2^{4}=16$ )
- Changing hexdigits to bits and back a ga in:

| Binary | 0101 | 1100 |
| :--- | :--- | :--- |
| Hex | 5 | C |
| Hex | 3 | G |
| Binary |  |  |



| Decimal |
| :---: |
| 0 |
| 1 |
| 2 |
| 3 |
| 4 |
| 5 |
| 6 |
| 7 |
| 8 |
| 9 |
| 10 |
| 11 |
| 12 |
| 13 |
| 14 |
| 15 |


| Hex |
| :--- |
| 0 |
| 1 |
| 2 |
| 3 |
| 3 |
| 4 |
| 5 |
| 6 |
| 7 |
| 8 |
| 9 |
| A |
| B |
| C |
| D |
| E |
| F |

Binary
0000
0001
0010
0011
0100
0101
0100
0111
1000
1001
1010
1011
1100
1101
1110
1111


- Early binary representation-1 and 0encoded numbers and keyboard characters
- Now representation for sound, video, a nd othertypes of information
- For encoding text, what symbols should be included?
* We want to keep the list small enough to use fewer bits, but we don't want to leave out critical characters


| Characters | Quantity |
| :--- | :--- |
| Uppercase letters | 26 |
| Lowercase letters | 26 |
| Arithmetic characters (0-9) | 10 |
| Punctuation characters <br> (including space) | 20 |
| Non-Printable characters | 3 |
|  | Total | 95 



| 0 | NUL | 1 | SOH | 2 | STX | 3 | ETX | 4 | EOT | 5 | ENQ | 6 | ACK | 7 | BEL |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 8 | BS | 9 | HT | 10 | NL | 11 | VT | 12 | NP | 13 | CR | 14 | So | 15 | SI |
| 16 | DLE | 17 | DC1 | 18 | DC2 | 19 | DC3 | 20 | DC4 | 21 | NAK | 22 | SYN | 23 | ETB |
| 24 | CAN | 25 | EM | 26 | SUB | 27 | ESC | 28 | FS | 29 | GS | 30 | RS | 31 | US |
| 32 | SP | 33 | ! | 34 | " | 35 | \# | 36 | \$ | 37 | \% | 38 | \& | 39 |  |
| 40 | ( | 41 | ) | 42 | * | 43 | + | 44 |  | 45 | - | 46 |  | 47 | / |
| 48 | 0 | 49 | 1 | 50 | 2 | 51 | 3 | 52 | 4 | 53 | 5 | 54 | 6 | 55 | 7 |
| 56 | 8 | 57 | 9 | 58 | : | 59 | ; | 60 | < | 61 | = | 62 | > | 63 | ? |
| 64 | @ | 65 | A | 66 | B | 67 | C | 68 | D | 69 | E | 70 | F | 71 | G |
| 72 | H | 73 | I | 74 | J | 75 | K | 76 | L | 77 | M | 78 | N | 79 | 0 |
| 80 | P | 81 | Q | 82 | R | 83 | S | 84 | T | 85 | U | 86 | V | 87 | W |
| 88 | X | 89 | Y | 90 | Z | 91 | [ | 92 | \} | 93 | ] | 94 | $\wedge$ | 95 | - |
| 96 |  | 97 | a | 98 | b | 99 | c | 100 | d | 101 | e | 102 | f | 103 | g |
| 104 | h | 105 | i | 106 | j | 107 | k | 108 | 1 | 109 | m | 110 | n | 111 | o |
| 112 | p | 113 | q | 114 | r | 115 | S | 116 | t | 117 | u | 118 | v | 119 | W |
| 120 | x | 121 | y | 122 | z | 123 | \{ | 124 | \| | 125 | \} | 126 | $\sim$ | 127 | DEL |



## Hexadecimal - Character

| 00 | NUL | 01 | SOH | 02 | STX | 03 | ETX | 04 | EOT | 05 | ENQ | 06 | ACK |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | $0^{07}$ BEL



- By the mid-1960's, it became clear that 7bit ASC II was not enough to represent text from la nguages other than English
- IBM extended ASC II to 8 bits (256 symbols)
- Called "Extended ASC II," the first half is original ASCII with a 0 added at the beginning of each group of bits
- Ha nd les most Westem la ng ua ges a nd additional useful symbols


Character Binary

| $\#$ | 0010 | 0011 |
| :--- | :--- | :--- |
| $\odot$ | 1010 | 1001 |
| $\dot{e}$ |  | 1110 |


| ASCII | $\begin{aligned} & 0 \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \\ & 0 \\ & 1 \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \\ & 1 \\ & 0 \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & 0 \\ & 1 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & 0 \\ & 1 \\ & 0 \\ & 1 \end{aligned}$ | $\begin{aligned} & 0 \\ & 1 \\ & 1 \\ & 0 \end{aligned}$ | $\begin{aligned} & 0 \\ & 1 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & 1 \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & 1 \\ & 0 \\ & 0 \\ & 1 \end{aligned}$ | $\begin{aligned} & 1 \\ & 0 \\ & 1 \\ & 0 \end{aligned}$ | $\begin{aligned} & 1 \\ & 0 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & 1 \\ & 1 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & 1 \\ & 1 \\ & 0 \\ & 1 \end{aligned}$ | $\begin{aligned} & 1 \\ & 1 \\ & 1 \\ & 0 \end{aligned}$ | 1 1 1 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0000 | $\mathrm{Nu}_{u}$ | $\mathrm{s}_{\mathrm{H}}$ | ${ }^{\text {s }} \mathrm{x}$ | ${ }^{5} \times$ | $\mathrm{E}_{\mathrm{T}}$ | $\mathrm{E}_{0}$ | ${ }^{\text {A }}$ | ${ }_{\text {B }}$ | ${ }^{\text {B }}$ | ${ }_{\text {H }}$ | ${ }_{\text {L }}$ | $\mathrm{v}_{\mathrm{T}}$ | $\mathrm{F}_{\mathrm{F}}$ | $\mathrm{C}_{\mathrm{R}}$ | so | $\mathrm{s}_{1}$ |
| 0001 | $\mathrm{D}_{\mathrm{L}}$ | ${ }_{1}$ | $\mathrm{D}_{2}$ | $\mathrm{O}_{3}$ | $\mathrm{D}_{4}$ | ${ }^{\text {N }}$ | $\mathrm{s}_{\mathrm{r}}$ | $\mathrm{E}_{\Sigma}$ | $\mathrm{c}_{\mathrm{N}}$ | $\mathrm{E}_{\mathrm{M}}$ | $\mathrm{s}_{\mathrm{B}}$ | ${ }_{\text {E }}^{\text {c }}$ | $\mathrm{F}_{\mathrm{s}}$ | ${ }^{6}$ | $\mathrm{R}_{\mathrm{s}}$ | $u_{\text {s }}$ |
| 0010 |  | ！ | 11 | \＃ | \＄ | \％ | \＆ | ＇ | （ | ） | ＊ | ＋ | ， | － | ． | ／ |
| 0011 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | ： | ； | ＜ | ＝ | ＞ | ？ |
| 0100 | ＠ | A | B | C | D | E | F | G | H | I | J | K | L | M | N | 0 |
| 0101 | P | Q | R | S | T | U | V | W | X | Y | Z | ［ | 1 | ］ | $\wedge$ | － |
| 0110 | － | a | b | C | d | e | f | 9 | h | i | j | k | 1 | m | n | $\bigcirc$ |
| 0111 | p | q | r | S | t | u | V | W | X | Y | z | \｛ |  | \} | ～ | ${ }^{\text {D }}$ |
| 1000 | ${ }_{8}^{8}$ | ${ }^{8} 1$ | $8_{2}$ | 83 | ${ }^{\text {in }}$ | $\mathrm{N}_{\mathrm{L}}$ | $\mathrm{s}_{\text {s }}$ | $\mathrm{E}_{\text {S }}$ | $\mathrm{H}_{\text {s }}$ | $\mathrm{H}_{3}$ | ${ }^{\text {r }}$ | $P_{0}$ | $\mathrm{P}_{\mathrm{v}}$ | ${ }_{\text {R }}$ | $\mathrm{s}_{2}$ | $\mathrm{s}_{3}$ |
| 1001 | ${ }^{\text {c }}$ | $\mathrm{P}_{1}$ | $\mathrm{P}_{\mathrm{z}}$ | $\mathrm{s}_{\mathrm{E}}$ | $\mathrm{c}_{\mathrm{c}}$ | $\mathrm{m}_{\mathrm{M}}$ | $\mathrm{s}_{\mathrm{p}}$ | $\mathrm{E}_{\mathrm{p}}$ | $\mathrm{o}_{8}$ | $0_{0}$ | ${ }_{\text {a }}$ | $\mathrm{c}_{\mathrm{s}}$ | ${ }_{\text {s }}^{\text {T }}$ | $\mathrm{o}_{\mathrm{s}}$ | ${ }^{\text {P }}$ | $A_{p}$ |
| 1010 | $A_{0}$ | $i$ | ¢ | £ |  | 玨 | $1$ | § | ．． | （C） | $\bigcirc$ | ［ | ᄀ | － |  | － |
| 1011 | － | $\pm$ | 2 | ${ }^{3}$ | － | $\mu$ | ¢ | － | ， | 1 | $\sigma^{*}$ | 3 | 1／4 | 1／2 | 3／4 | ¿ |
| 1100 | A | Á | Â | Ã | À | A | 厄 | Ç | E | É | $\hat{\mathrm{E}}$ | E | I | İ | $\hat{\mathrm{I}}$ | $\stackrel{\text { I }}{ }$ |
| 1101 | Đ | N | O | O | Ô | Õ | O | $\times$ | $\varnothing$ | U̇ | Ú | U | Ü | Y＇ | P | $\beta$ |
| 1110 | à | á | â | ã | a | $\stackrel{\circ}{\circ}$ | æ | C | è | é | e | e | 1 | 1 | ̂̂ | 1 |
| 1111 | ठ | ñ | ò | ó | ô | õ | O | $\div$ | $\varnothing$ | ù | ú | ut | ü | 立 | P | $\ddot{\mathrm{y}}$ |



Digits in Phone number

| 888 | 0011 <br> 0011 <br> 0000 <br> 0011 | 1000 |
| :--- | :--- | :--- |
| 5 | 0011 0101 <br> 0011 0101 <br> 0011 0101 |  |
|  | 0011 | 0001 |
| 1 | 0011 | 0010 |
| 2 | 0011 | 0001 |
| 1 | 0011 | 0010 |
| 2 |  |  |


| ASCII | $\begin{aligned} & 0 \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \\ & 0 \\ & 1 \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \\ & 1 \\ & 0 \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & 0 \\ & 1 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & 0 \\ & 1 \\ & 0 \\ & 1 \end{aligned}$ | $\begin{aligned} & 0 \\ & 1 \\ & 1 \\ & 0 \end{aligned}$ | $\begin{aligned} & 0 \\ & 1 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & 1 \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & 1 \\ & 0 \\ & 0 \\ & 1 \end{aligned}$ | $\begin{aligned} & 1 \\ & 0 \\ & 1 \\ & 0 \end{aligned}$ | $\begin{aligned} & 1 \\ & 0 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & 1 \\ & 1 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & 1 \\ & 1 \\ & 0 \\ & 1 \end{aligned}$ | $\begin{aligned} & 1 \\ & 1 \\ & 1 \\ & 0 \end{aligned}$ | 1 1 1 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0000 | $\mathrm{N}_{\mathrm{u}}$ | ${ }^{\text {S }}$ | ${ }^{\text {s }} \times$ | ${ }^{\text {E }}$ ¢ | $\mathrm{E}_{\mathrm{T}}$ | $\varepsilon_{0}$ | ${ }^{\text {A }}$ | ${ }^{\text {B }}$ | ${ }_{5}$ | ${ }_{\text {H }}$ | ${ }_{\text {L }}$ | ${ }_{\text {T }}$ | $\mathrm{F}_{\mathrm{F}}$ | $\mathrm{c}_{\mathrm{B}}$ | so | $\mathrm{s}_{\mathrm{I}}$ |
| 0001 | $\mathrm{D}_{\mathrm{L}}$ | $\mathrm{o}_{1}$ | $\mathrm{O}_{2}$ | $\mathrm{D}_{3}$ | $\mathrm{D}_{4}$ | ${ }^{\text {K }}$ | ${ }_{\text {s }}{ }^{\text {r }}$ | $\mathrm{E}_{\mathrm{E}}$ | ${ }^{c_{N}}$ | ${ }_{\text {E }}^{\text {M }}$ | $\mathrm{S}_{\mathrm{B}}$ | ${ }_{\text {E }}$ | $\mathrm{F}_{\mathrm{s}}$ | ${ }^{\text {G }}$ | ${ }^{\text {R }}$ | $u_{s}$ |
| 0010 |  | ! | " | \# | \$ | \% | \& | ' | 1 | ) | * | + | , | - | . | / |
| 0011 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | : | ; | < | = | $>$ | ? |
| 0100 | @ | A | B | C | D | E | F | G | H | I | J | K | L | M | N | 0 |
| 0101 | P | Q | R | S | T | U | V | W | X | Y | Z | [ | $\backslash$ | ] | $\wedge$ |  |
| 0110 | - | a | b | C | d | e | f | 9 | h | i | j | k | 1 | m | n | - |
| 0111 | p | q | r | S | t | u | V | W | X | Y | z | \{ |  | \} | $\sim$ | ${ }^{\text {d }}$ |
| 1000 | ${ }_{8}^{8}$ | ${ }^{8} 1$ | ${ }_{8}$ | ${ }^{8}$ | ${ }^{\text {IN }}$ | NL | $\mathrm{s}_{\mathrm{s}}$ | $\mathrm{E}_{\text {S }}$ | ${ }_{\text {H }}$ | ${ }^{\text {J }}$ | ${ }^{\text {r }}$ | ${ }_{\text {P }}$ | $\mathrm{P}_{\mathrm{v}}$ | ${ }_{\text {R }}^{\text {I }}$ | $\mathrm{S}_{2}$ | $\mathrm{s}_{3}$ |
| 1001 | ${ }^{\text {D }}$ | $\mathrm{P}_{1}$ | $\mathrm{P}_{\mathrm{z}}$ | $\mathrm{S}_{\mathrm{E}}$ | ${ }^{\text {c }}$ c | $M_{M}$ | $\mathrm{s}_{\mathrm{p}}$ | $\mathrm{E}_{\mathrm{p}}$ | $\mathrm{a}_{8}$ | ${ }^{\circ}{ }_{0}$ | ${ }^{\circ}{ }_{\text {A }}$ | $\mathrm{c}_{\text {s }}$ | ${ }_{\text {st }}$ | ${ }^{\circ}$ s | ${ }^{\text {P }}$ M | $A_{p}$ |
| 1010 | ${ }^{\text {a }}$ | i | ¢ | £ |  | ¥ | ! | § | .. | (C) | $\bigcirc$ | [ | $\neg$ | - | (R) |  |
| 1011 | - | $\pm$ | 2 | ${ }^{3}$ | - | $\mu$ | \\| | . | , | 1 | ${ }^{*}$ | \} | 1/4 | 1/2 | 3/4 |  |
| 1100 | A | Á | Â | Ã | Ȧ | A | 厌 | Ç | E | E | E | E | I | İ | $\hat{\text { I }}$ | İ |
| 1101 | Đ | N | O | O | Ô | Õ | Ö | $\times$ | $\varnothing$ | U̇ | Ú | U | Ü | Ý | P | $\beta$ |
| 1110 | à | á | â | ã | ä | $\stackrel{\circ}{\text { a }}$ | æ | ¢ | è | é | $\hat{e}$ | è | 1 | 1 | ̂ | i |
| 1111 | ð | ñ | ò | ó | ô | o | ○ | $\div$ | $\varnothing$ | ù | ú | û | ü | ý | P | \#̈ |



Digits in Phone number

| 8 | 0011 | 1000 |
| :--- | :--- | :--- |
| 5 | 0011 | 0101 |
| 1 | 0011 | 0001 |
| 2 | 0011 | 0010 |


| ASCII | $\begin{aligned} & 0 \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & \hline 0 \\ & 0 \\ & 0 \\ & 1 \end{aligned}$ | $\begin{array}{\|l\|} \hline 0 \\ 0 \\ 1 \\ 0 \end{array}$ | $\begin{aligned} & 0 \\ & 0 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{array}{\|l} \hline 0 \\ 1 \\ 0 \\ 0 \end{array}$ | $\begin{aligned} & 0 \\ & 1 \\ & 0 \\ & 1 \end{aligned}$ | $\begin{aligned} & 0 \\ & 1 \\ & 1 \\ & 0 \end{aligned}$ | $\begin{aligned} & 0 \\ & 1 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & 1 \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{array}{\|l\|} \hline 1 \\ 0 \\ 0 \\ 1 \end{array}$ | $\begin{aligned} & 1 \\ & 0 \\ & 1 \\ & 0 \end{aligned}$ | $\begin{aligned} & 1 \\ & 0 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & \hline 1 \\ & 1 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & 1 \\ & 1 \\ & 0 \\ & 1 \end{aligned}$ | $\begin{aligned} & 1 \\ & 1 \\ & 1 \\ & 0 \end{aligned}$ | 1 1 1 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0000 | $\mathrm{Nu}^{0}$ | $\mathrm{s}_{\mathrm{H}}$ | ${ }^{\text {s }} \times$ | ${ }_{\text {E }}$ | $\mathrm{E}_{\mathrm{T}}$ | ${ }_{0}$ | ${ }^{\text {a }}$ | ${ }_{\text {B }}$ | ${ }^{8}$ | ${ }_{\text {H }}$ | ${ }_{\text {L }}$ | ${ }_{\text {r }}^{T}$ | $\mathrm{F}_{\mathrm{F}}$ | $\mathrm{c}_{\text {R }}$ | $\mathrm{s}_{0}$ | $\mathrm{s}_{1}$ |
| 0001 | $\mathrm{D}_{\mathrm{L}}$ | ${ }_{0}$ | $\mathrm{D}_{2}$ | $\mathrm{D}_{3}$ | $\mathrm{D}_{4}$ | ${ }^{\text {NK}}$ | $\mathrm{s}_{\mathrm{r}}$ | $\mathrm{E}_{\mathrm{\Sigma}}$ | ${ }^{\text {c }}$ N | ${ }_{\text {E }}^{\text {M }}$ | $\mathrm{s}_{\mathrm{B}}$ | ${ }_{\text {E }}$ | ${ }^{\text {F }}$ | ${ }^{\text {G }}$ | ${ }^{\text {R }}$ | $u_{s}$ |
| 0010 |  | $!$ | " | \# | \$ | \% | \& | ' |  | ) | * | + | , | - | . | / |
| 0011 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | : | ; | $<$ | = | > | ? |
| 0100 | @ | A | B | C | D | E | F | G | H | I | J | K | L | M | N | 0 |
| 0101 | P | Q | R | S | T | U | V | W | X | Y | Z | [ | $\backslash$ | ] | $\wedge$ |  |
| 0110 | - | a | b | C | d | e | f | 9 | h | i | j | k | 1 | m | n |  |
| 0111 | p | q | r | S | t | u | V | W | x | Y | z | \{ |  | \} | ~ |  |
| 1000 | ${ }_{8}^{8}$ | ${ }_{1} 1$ | 8 | ${ }^{8}$ | ${ }^{\text {I }}$ N | NL | $\mathrm{s}_{\mathrm{s}}$ | ${ }^{\text {E }}$ | $\mathrm{H}_{\text {s }}$ | ${ }^{\text {J }}$ | ${ }^{\text {s }}$ | $\mathrm{P}_{\mathrm{D}}$ | $\mathrm{P}_{\mathrm{v}}$ | ${ }_{\text {R }}$ | $\mathrm{S}_{2}$ | $s_{3}$ |
| 1001 | $\mathrm{D}_{\mathrm{c}}$ | ${ }^{\text {P }} 1$ | $\mathrm{P}_{\mathrm{z}}$ | $\mathrm{S}_{\mathrm{E}}$ | ${ }^{\text {c }}$ c | M ${ }_{\text {M }}$ | $\mathrm{s}_{\mathrm{p}}$ | $\mathrm{E}_{\mathrm{p}}$ | $\mathrm{O}_{8}$ | ${ }^{\circ}{ }_{0}$ | ${ }^{\circ}{ }_{A}$ | $\mathrm{c}_{\mathrm{s}}$ | ${ }_{\text {st }}$ | ${ }^{\circ}$ | ${ }^{\text {P }}$ M | ${ }^{\text {A }}$ |
| 1010 | ${ }^{\text {a }}$ | i | ¢ | £ |  | 装 | ! | § | .. | (C) | $\bigcirc$ | [1 | ᄀ | - |  |  |
| 1011 | - | $\pm$ | 2 | ${ }^{3}$ | - | $\mu$ | d |  |  | 1 | $\sigma^{*}$ | \} | 1/4 | 1/2 | 3/4 |  |
| 1100 | A | Á | Â | Ã | A | A | E | Ç | E | É | E | E | I | I | I | I |
| 1101 | Đ | N | O | O | Ô | Õ | O | $\times$ | $\varnothing$ | U̇ | U' | $\hat{U}$ | Ü | Y' | P | $\beta$ |
| 1110 | à | á | â | ã | à | $\stackrel{\circ}{\text { a }}$ | æ | ¢ | è | é | ê | è | 1 | 1 | ̂̂ |  |
| 1111 | ठ | n | ò | ó | ô | o | 0 | $\div$ | $\varnothing$ | ù | ú | û | ü | ý | P | У̀ |



- Several languages around the world have more than 256 individual characters
- Unic ode uses 16 bits; $2^{16}=65536$ characters
* $1^{\ddagger} 7$ bits (128 chars) are ASCII chars * Different locales-different characters beyond $1^{\text {tt }} 7$ bits

- The code forbroadcast communication is purposefully ineffic ient, to be distinctive when spoken a mid noise

Table 8.4 NATO broadcast alphabet designed not to be minimal

| A | Alpha | H | Hotel | O | Oscar | V | Victor |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| B | Bravo | I | India | P | Papa | W | Whiskey |
| C | Charlie | J | Juliet | Q | Quebec | X | X-ray |
| D | Delta | K | Kilo | R | Romeo | Y | Yankee |
| E | Echo | L | Lima | S | Sierra | Z | Zulu |
| F | Foxtrot | M | Mike | T | Tango |  |  |
| G | Golf | N | November | U | Uniform |  |  |



- Extended ASCII encodesletters and characters well, but most documents conta in more than just text.
* Format information like font, font size, justification
- Formatting characters could be added to ASC II, but that mixes the c ontent with the description of its form (metadata)
* Metadata is "data about data"
- ${ }_{8,35}$ Metadata is represented using tags, asin ${ }^{8,35} \mathrm{HTML}$


Bits and bytes encode the information, but that's not all

* Tags encode format and some structure in word processors
* Tags enc ode format and some structure in HTML
* In the Oxford English Dictionary tags encode structure and some fomatting

- Oxford Eng lish Dictionary (OED) printed version is 20 volumes
- We could type the entire contents as ASC II characters (in a bout 120 years), but searching would be diffic ult
* Suppose you search for the word "set." It is included in many other words like closet, ho rseta il, settle, etc
* How will the software know what characters comprise the definition of set?
- Inc orporate meta data

- Special set of tags was developed to specify OED's structure
* <uw>means headword, the word being defined
* Othertagslabel pronunciation $\langle p r$, phonetic notation <ph>, parts of speech <ps>
- The tags do not print. They are there only to specify structure so the computer knows what part of the dictionary it is looking at

byte (baIt). Computers. [Arbitrary, prob. influenced by bit sb. ${ }^{4}$ and bite sb.] A group of eight consecutive bits operated on as a unit in a computer. 1964 Blaauw \& Brooks in IBM Systems JrnI. III. 122 An 8bit unit of information is fundamental to most of the formats [of the System/360]. A consecutive group of $n$ such units constitutes a field of length $n$. Fixed-length fields of length one, two, four, and eight are termed bytes, halfwords, words, and double words respectively. 1964 IBM Jrnl. Res. \& Developm. VIII. 97/1 When a byte of data appears from an I/O device, the CPU is seized, dumped, used and restored. 1967 P. A. Stark Digital Computer Programming xix. 351 The normal operations in fixed point are done on four bytes at a time. 1968 Dataweek 24 Jan. $1 / 1$ Tape reading and writing is at from 34,160 to 192,000 bytes per second.

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<e><hg><hw>byte</hw> <pr><ph>baIt</ph></pr></hg>. <la>Computers</la>.
<etym>Arbitrary, prob. influenced by <xr><x>bit</x></xr>
<ps>n.<hm>4</hm></ps>and <xr><x>bite</x> <ps>n.</ps> </xr></etym>
<s4>A group of eight consecutive bits operated on as a unit in a
computer.</s4> <qp><q><qd>1964</qd><a>Blaauw</a> &amp. <a>Brooks</a>
<bib>in</bi.b> <w>IBM Systems Jrnl.</w> <lc>III. 122</lc> <qt>An 8-bit
unit of information is fundamental to most of the formats <ed>of the
System/360</ed>.&es.A consecutive group of <i>n</i> such units
constitutes a field of length <i>n</i>.&es.Fixed-length fields of
length one, two, four, and eight are termed bytes, halfwords, words,
and double words respectively. </qt></q><q><qd>1964</qd> <w>IBM Jrnl.
Res. &amp. Developm.</w> <lc>VIII. 97/1</lc> <qt>When a byte of data
appears from an I/O device, the CPU is seized, dumped, used and
restored.</qt></q> <q><qd>1967</qd> <a>P. A. Stark</a> <w>Digital
```

Table 8.3 Sixteen symbols of the 4-bit PandA representation

| Symbol | Binary | Physical Bits | Hex | Symbol | Binary | Physical Bits | Hex |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| AAAA | 0000 |  | 0 | PAAA | 1000 |  | 8 |
| AAAP | 0001 |  | 1 | PAAP | 1001 |  | 9 |
| AAPA | 0010 | $2+25$ | 2 | PAPA | 1010 |  | A |
| AAPP | 0011 |  | 3 | PAPP | 1011 |  | B |
| APAA | 0100 |  | 4 | PPAA | 1100 |  | C |
| APAP | 0101 |  | 5 | PPAP | 1101 |  | D |
| APPA | 0110 |  | 6 | PPPA | 1110 |  | E |
| APPP | 0111 |  | 7 | PPPP | 1111 |  | F |



Course Web site:

* http://www.cs.washington.edu/educat ion/courses/100/07au/

Address munging: http://www.addressmunger.com


IT joins physic al \& logic al doma ins so physical devices do our logical work

* Symbols represent things 1-to-1
* Create symbols by grouping pattems
* PandA representation is fundamental
* Bit, a place where 2 pattems set/detect * ASC Il is a byte encoding of Latin abet
* In addition to content, encode structure with tags

