CSE 142 Computer Programming I

Sorting

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Overview

Sorting defined Algorithms for sorting Selection Sort algorithm Efficiency of Selection Sort

Q-2

Sorting

The problem: put things in order
Usually smallest to largest: "ascending"
Could also be largest to smallest:
"descending"

Lots of applications!
ordering hits in web search engine
preparing lists of output
merging data from multiple sources
to help solve other problems
faster search (allows binary search)
too many to mention!

Sorting: More Formally

Given an array b[0], b[1], ... b[n-1], reorder entries so that b[0] <= b[1] <= ... <= b[n-1]

Shorthand for these slides: the notation <code>array[i..k]</code> means all of the elements <code>array[i],array[i+1]...array[k]</code>
Using this notation, the entire array would be: <code>b[0..n-1]</code>

P.S.: This is not C syntax!

Sorting Algorithms

Sorting has been intensively studied for decades Many different ways to do it!

We'll look at only one algorithm, called

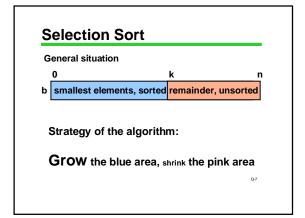
"Selection Sort"

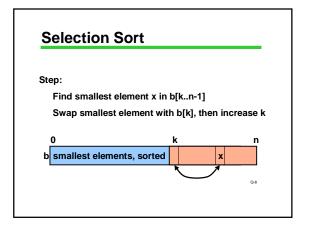
Other algorithms you might hear about in other courses include Bubble Sort, Insertion Sort, QuickSort, and MergeSort. And that's only the beginning!

Q-5

Sorting Problem What we want: Data sorted in order o n sorted: b[0]<=b[1]<=...<=b[n-1] Initial conditions

unsorted



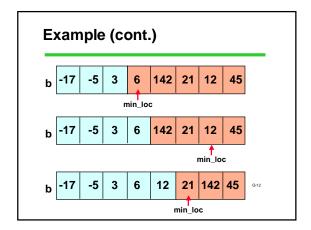



```
Code for Selection Sort

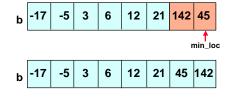
/* Sort b[0..n-1] in non-decreasing order (rearrange elements in b so that b[0]<=b[1]<=...<=b[n-1]) */

void sel_sort (int b[], int n) {
    int k, m;
    for (k = 0; k < n - 1; k = k + 1) {
        m = min_loc(b,k,n);
        swap(&b[k], &b[m]);
    }
}
```

```
Example
                             -17
                                 45
                           min_loc
                6
                   142
                         21
                             3
          min_loc
       -5
           12
                6
                   142
                         21
                             3
                                 45
b
                           min_loc
```



Example (concluded)



Sorting Analysis

How many steps are needed to sort n things?

For each swap, we have to search the remaining array

length is proportional to original array length *n*Need about *n* search/swap passes

Total number of steps proportional to *n*²

Conclusion: selection sort is pretty expensive (slow) for large *n*

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Can We Do Better Than n2?

Sure we can!
Selection, insertion, bubble sorts are all proportional to n²
Other sorts are proportional to n log n
Mergesort
Quicksort

etc.

 $log\ n$ is considerably smaller than n, especially as n gets larger

As the size of our problem grows, the time to run a n^2 sort will grow much faster than an $n \log n$ one.

Any better than *n log n*?

In general, no. But in special cases, we can do better

Example: Sort exams by score: drop each exam in one of 101 piles; work is proportional to *n*

Curious fact: efficiency can be studied and predicted mathematically, without using a computer at all!

Comments about Efficiency

Efficiency means doing things in a way that saves resources

Usually measured by *time* or *memory*

Many small programming details have little or no measurable effect on efficiency The big differences comes with the right choice of algorithm and/or data structure

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Summary

Sorting means placing things in order Selection sort is one of many algorithms

At each step, finds the smallest remaining value

Selection sort requires on the order of n² steps

There are sorting algorithms which are greatly more efficient

It's the algorithm that makes the difference, not the coding details

Q-18