CSE 142
Computer Programming I

## Sorting

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## Sorting

The problem: put things in order Usually smallest to largest: "ascending"
Could also be largest to smallest:
"descending"
Lots of applications!
ordering hits in web search engine
preparing lists of output
merging data from multiple sources
to help solve other problems
faster search (allows binary search) too many to mention!

## Sorting Algorithms

Sorting has been intensively studied for decades Many different ways to do it!
We'll look at only one algorithm, called "Selection Sort"

Other algorithms you might hear about in other courses include Bubble Sort, Insertion Sort, QuickSort, and MergeSort. And that's only the beginning!

## Overview

Sorting defined
Algorithms for sorting
Selection Sort algorithm Efficiency of Selection Sort

## Sorting: More Formally

Given an array b[0], b[1], ... b[n-1],
reorder entries so that
$\mathrm{b}[0]<=\mathrm{b}[1]<=. . .<=\mathrm{b}[\mathrm{n}-1]$
Shorthand for these slides: the notation array[i..k] means all of the elements
array[i],array[i+1]...array[k]
Using this notation, the entire array would be:
b[0..n-1]
P.S.: This is not C syntax!
0.4

| Sorting Algorithms |
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## Subproblem: Find Smallest

/* Find location of smallest element in b[k..n-1] */
/* Assumption: $\mathbf{k}$ < $\mathbf{n}$ */
/* Returns index of smallest, does not return the smallest value itself */
int min_loc (int b[ ], int k, int n) \{ int $\mathrm{j}, \overline{\text { pos }}$; /* $\mathrm{b}[$ pos] is smallest element */ pos $=\mathbf{k}$;
for ( $\mathrm{j}=\mathrm{k}, \mathrm{k}+\mathrm{i} ; \mathrm{j}<\mathrm{n} ; \mathrm{j}=\mathrm{j}+1$ )
if (b[j] < b[pos])
pos $=\mathrm{j}$;
return pos;
\}

## Code for Selection Sort

/* Sort b[0..n-1] in non-decreasing order (rearrange elements in $b$ so that $\mathrm{b}[0]<=\mathrm{b}[1]<=\ldots<=\mathrm{b}[\mathrm{n}-1]$ ) */
void sel_sort (int b[ ], int n) \{
int $k$, m ;
for $(k=0 ; k<n-1 ; k=k+1)\{$ m = min_loc(b,k,n); swap(\&b[k], \&b[m]);
\}
\}

## Example (cont.)


b



Example (concluded)


## Can We Do Better Than $\boldsymbol{n}^{\mathbf{2}}$ ?

## Sure we can!

Selection, insertion, bubble sorts are al proportional to $n^{2}$
Other sorts are proportional to $n \log n$
Mergesort
Quicksort
etc.
$\log \mathrm{n}$ is considerably smaller than n , especially as n gets larger

As the size of our problem grows, the time to run a $n^{2}$ sort will grow much faster than an $n \log n$ one.

## Comments about Efficiency

Efficiency means doing things in a way that saves resources

Usually measured by time or memory used
Many small programming details have little or no measurable effect on efficiency The big differences comes with the right choice of algorithm and/or data structure

## Sorting Analysis

How many steps are needed to sort n things?
For each swap, we have to search the remaining array
length is proportional to original array length $n$
Need about $n$ search/swap passes
Total number of steps proportional to $n^{2}$
Conclusion: selection sort is pretty expensive (slow) for large $n$

## Any better than $\boldsymbol{n} \log \boldsymbol{n}$ ?

In general, no. But in special cases, we can do better
Example: Sort exams by score: drop each exam in one of 101 piles; work is proportional to $n$

Curious fact: efficiency can be studied and predicted mathematically, without using a computer at all!

## Summary

Sorting means placing things in order Selection sort is one of many algorithms

At each step, finds the smallest remaining value
Selection sort requires on the order of $\mathrm{n}^{2}$ steps

There are sorting algorithms which are greatly more efficient
It's the algorithm that makes the difference, not the coding details

