## **CSE 311: Foundations of Computing**

#### Lecture 13: Number theory & modular arithmetic



- Arithmetic over a finite domain
- In computing, almost all computations are over a finite domain

#### Number Theory (and applications to computing)

- Branch of Mathematics with direct relevance to computing
- Many significant applications
  - Cryptography
  - Hashing
  - Security
- Important tool set

# I'm ALIVE!

# I'm ALIVE!

```
public class Test {
   final static int SEC IN YEAR = 364*24*60*60*100;
   public static void main(String args[]) {
       System.out.println(
          "I will be alive for at least " +
          SEC IN YEAR * 101 + " seconds."
      );
         ----jGRASP exec: java Test
        I will be alive for at least -186619904 seconds.
          ----jGRASP: operation complete.
```

## Divisibility

#### **Definition: "a divides b"**

For 
$$a \in \mathbb{Z}, b \in \mathbb{Z}$$
 with  $a \neq 0$ :  
 $a \mid b \leftrightarrow \exists k \in \mathbb{Z} \ (b = ka)$ 

Check Your Understanding. Which of the following are true?

# Divisibility

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#### **Division Theorem**

For  $a \in \mathbb{Z}$ ,  $d \in \mathbb{Z}$  with d > 0there exist *unique* integers q, r with  $0 \le r < d$ such that a = dq + r.

To put it another way, if we divide *d* into *a*, we get a unique quotient  $q = a \operatorname{div} d$ and non-negative remainder r = a % d

> Note:  $r \ge 0$  even if a < 0. Not quite the same as in Java

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```
public class Test2 {
    public static void main(String args[]) {
        int a = -5;
        int d = 2;
        System.out.println(a % d);
    }
} Note: r ≥ 0 even if a < 0.
Not quite the same as in Java</pre>
```

Definition: "a is congruent to b modulo m"

For  $a, b, m \in \mathbb{Z}$  with m > 0 $a \equiv b \pmod{m} \leftrightarrow m \mid (a - b)$ 

Check Your Understanding. What do each of these mean? When are they true?

 $-1 \equiv 19 \pmod{5}$ 

 $y \equiv 2 \pmod{7}$ 

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 $x \equiv 0 \pmod{2}$ 

This statement is the same as saying "x is even"; so, any x that is even (including negative even numbers) will work.

 $-1 \equiv 19 \pmod{5}$   $5 - 1 - 19 \equiv 5 - 20 \equiv 3k$ . -20 = 5k

This statement is true. 19 - (-1) = 20 which is divisible by 5

$$y \equiv 2 \pmod{7}$$
  $7/y-2 = 3k$ .  $4-7=7k = 3k$ .  $y=2-7k$ 

This statement is true for y in  $\{ ..., -12, -5, 2, 9, 16, ... \}$ . In other words, all y of the form 2+7k for k an integer.

# The % m function vs the $\equiv \pmod{m}$ predicate

- % is a function (operator) with two arguments. The result is an integer
- ≡ ... (mod m) is a predicate
  - "a is equivalent, modulo m, to b"
  - "a is equivalent to b (modulo m)"
  - $-a \equiv b \pmod{m}$

## Arithmetic, mod 7



#### **Modular Arithmetic: A Property**

Let a, b, m be integers with m > 0. Then,  $a \equiv b \pmod{m}$  if and only if a % m = b % m.

Suppose that  $a \equiv b \pmod{m}$ .

Suppose that a % m = b % m.

### **Modular Arithmetic: A Property**

Let a, b, m be integers with m > 0. Then,  $a \equiv b \pmod{m}$  if and only if a % m = b % m.

Suppose that  $a \equiv b \pmod{m}$ . Then,  $m \mid (a - b)$  by definition of congruence. So, a - b = km for some integer k by definition of divides. Therefore, a = b + km. Taking both sides modulo m we get: a % m = (b + km)% m = b % m.

Suppose that a % m = b % m.

By the division theorem, a = mq + (a % m) and

b = ms + (b % m) for some integers q,s.

Then, a - b = (mq + (a % m)) - (ms + (b % m))= m(q - s) + (a % m - b % m)= m(q - s) since a % m = b % mTherefore,  $m \mid (a - b)$  and so  $a \equiv b \pmod{m}$ . The % m function vs the  $\equiv \pmod{m}$  predicate

- What we have just shown
  - The % *m* function takes any  $a \in \mathbb{Z}$  and maps it to a remainder  $a \% m \in \{0, 1, ..., m 1\}$ .
  - Imagine grouping together all integers that have the same value of the % m function That is, the same remainder in  $\{0, 1, ..., m - 1\}$ .
  - The  $\equiv \pmod{m}$  predicate compares  $a, b \in \mathbb{Z}$ . It is true if and only if the % *m* function has the same value on *a* and on *b*.

That is, *a* and *b* are in the same group.

#### **Modular Arithmetic: Basic Property**

```
Let m be a positive integer.
If a \equiv b \pmod{m} and b \equiv c \pmod{m},
then a \equiv c \pmod{m}
```

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```

Suppose that  $a \equiv b \pmod{m}$  and  $b \equiv c \pmod{m}$ . Then, by the previous property, we have a % m = b % m and b % m = c % m.

Putting these together, we have a % m = c % m, which says that  $a \equiv c \pmod{m}$ , by definition.

So " $\equiv$ " behaves like "=" in that sense. And that is not the only similarity...

#### **Modular Arithmetic: Addition Property**

Let *m* be a positive integer. If  $a \equiv b \pmod{m}$  and  $c \equiv d \pmod{m}$ , then  $a + c \equiv b + d \pmod{m}$ 

#### **Modular Arithmetic: Addition Property**

Let *m* be a positive integer. If  $a \equiv b \pmod{m}$  and  $c \equiv d \pmod{m}$ , then  $a + c \equiv b + d \pmod{m}$ 

Suppose that  $a \equiv b \pmod{m}$  and  $c \equiv d \pmod{m}$ . Unrolling definitions gives us some k such that a - b = km, and some j such that c - d = jm.

Adding the equations together gives us (a + c) - (b + d) = m(k + j). Now, re-applying the definition of congruence gives us  $a + c \equiv b + d \pmod{m}$ .

### **Modular Arithmetic: Multiplication Property**

Let *m* be a positive integer. If  $a \equiv b \pmod{m}$  and  $c \equiv d \pmod{m}$ , then  $ac \equiv bd \pmod{m}$ 

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Let *m* be a positive integer. If  $a \equiv b \pmod{m}$  and  $c \equiv d \pmod{m}$ , then  $ac \equiv bd \pmod{m}$ 

Suppose that  $a \equiv b \pmod{m}$  and  $c \equiv d \pmod{m}$ . Unrolling definitions gives us some k such that a - b = km, and some j such that c - d = jm.

Then, a = km + b and c = jm + d. Multiplying both together gives us  $\underline{ac} = (km + b)(jm + d) = \underline{kjm^2 + kmd + bjm} + \underline{bd}$ .

Re-arranging gives us ac - bd = m(kjm + kd + bj). Using the definition of congruence gives us  $ac \equiv bd \pmod{m}$ .

# Lecture 13 Activity

You will be assigned to breakout rooms. Please:

- Introduce yourself
- Choose someone to share their screen, showing this PDF
- Consider the statement:

For all  $a, b, c, m \in \mathbb{Z}, m > 0$  one has  $a \equiv b \pmod{m} \rightarrow a + c \equiv b + c \pmod{m}$ .

- Discuss what the statement means.
- Prove the statement.

D	efinition: "a is congruent to b modulo m"									
	For $a, b, m \in \mathbb{Z}$ with $m > 0$ $a \equiv h \pmod{m} \Leftrightarrow m \mid (a - b)$									
	Definition: "a divides b"									

For  $a \in \mathbb{Z}$ ,  $b \in \mathbb{Z}$  with  $a \neq 0$ :

 $a \mid b \leftrightarrow \exists k \in \mathbb{Z} \ (b = ka)$ 

Fill out the poll everywhere for Activity Credit! Go to <u>pollev.com/philipmg</u> and login with your UW identity

# Lecture 13 Activity

You will be assigned to breakout rooms. Please:

- Introduce yourself
- · Choose someone to share their screen, showing this PDF
- Consider the statement:

For all  $a, b, c, m \in \mathbb{Z}, m > 0$  one has  $a \equiv b \pmod{m} \rightarrow a + c \equiv b + c \pmod{m}$ .

#### Proof.

Let  $a, b, c \in \mathbb{Z}$  be arbitrary and let m > 0. Assume that  $a \equiv b \pmod{m}$ . Then  $m \mid a - b$  and hence there is an integer x with mx = a - b. Then (a + c) - (b + c) = a - b = mxand so  $m \mid (a + c) - (b + c)$ . Then  $a + c \equiv b + c \pmod{m}$  by definition of mod.

Definition: "a is congruent to b modulo m"									
For $a, b, m \in \mathbb{Z}$ with $m > 0$ $a \equiv b \pmod{m} \leftrightarrow m \mid (a - b)$									
	Definition: "a divides b"								
	For $a \in \mathbb{Z}, b \in \mathbb{Z}$ with $a \neq 0$ :								

 $a \mid b \leftrightarrow \exists k \in \mathbb{Z} \ (b = ka)$ 

Let's start by looking a a small example:

 $0^2 = 0 \equiv 0 \pmod{4}$   $1^2 = 1 \equiv 1 \pmod{4}$   $2^2 = 4 \equiv 0 \pmod{4}$   $3^2 = 9 \equiv 1 \pmod{4}$  $4^2 = 16 \equiv 0 \pmod{4}$ 

Case 1 (n is even):

Let's start by looking a a small example:

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It looks like

$$\label{eq:new_relation} \begin{split} n &\equiv 0 \;(mod\;2) \rightarrow n^2 \equiv 0 \;(mod\;4) \text{, and} \\ n &\equiv 1 \;(mod\;2) \rightarrow n^2 \equiv 1 \;(mod\;4) \text{.} \end{split}$$

Case 1 (*n* is even): Suppose *n* is even. Then, n = 2k for some integer *k*. So,  $n^2 = (2k)2 = 4k^2$ . So, by definition of congruence, we have  $n^2 \equiv 0 \pmod{4}$ .

Let's start by looking a a small example:

 $0^{2} = 0 \equiv 0 \pmod{4}$   $1^{2} = 1 \equiv 1 \pmod{4}$   $2^{2} = 4 \equiv 0 \pmod{4}$   $3^{2} = 9 \equiv 1 \pmod{4}$  $4^{2} = 16 \equiv 0 \pmod{4}$ 

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Case 1 (n is even): Done.

Case 2 (n is odd):

Let's start by looking a a small example:

 $0^2 = 0 \equiv 0 \pmod{4}$   $1^2 = 1 \equiv 1 \pmod{4}$   $2^2 = 4 \equiv 0 \pmod{4}$   $3^2 = 9 \equiv 1 \pmod{4}$  $4^2 = 16 \equiv 0 \pmod{4}$ 

It looks like

$$\label{eq:new_relation} \begin{split} n &\equiv 0 \;(mod\;2) \rightarrow n^2 \equiv 0 \;(mod\;4) \text{, and} \\ n &\equiv 1 \;(mod\;2) \rightarrow n^2 \equiv 1 \;(mod\;4) \text{.} \end{split}$$

Let *n* be an integer. Prove that  $n^2 \equiv 0 \pmod{4}$  or  $n^2 \equiv 1 \pmod{4}$ Let's start by looking a a small example: Case 1 (*n* is even): Done.  $0^2 = 0 \equiv 0 \pmod{4}$  $1^2 = 1 \equiv 1 \pmod{4}$ Case 2 (*n* is odd):  $2^2 = 4 \equiv 0 \pmod{4}$ Suppose *n* is odd.  $3^2 = 9 \equiv 1 \pmod{4}$ Then, n = 2k + 1 for some integer k.  $4^2 = 16 \equiv 0 \pmod{4}$ So,  $n^2 = (2k + 1)^2$  $=4k^{2}+4k+1$ It looks like  $= 4(k^2 + k) + 1.$  $n \equiv 0 \pmod{2} \rightarrow n^2 \equiv 0 \pmod{4}$ , and So, by the earlier property of mod,  $n \equiv 1 \pmod{2} \rightarrow n^2 \equiv 1 \pmod{4}$ . we have  $n^2 \equiv 1 \pmod{4}$ .

Result follows by "proof by cases": n is either even or not even (odd)

# n-bit Unsigned Integer Representation

- Represent integer *x* as sum of powers of 2:
  - If  $\sum_{i=0}^{n-1} b_i 2^i$  where each  $b_i \in \{0,1\}$

then representation is  $b_{n-1}...b_2 b_1 b_0$ 

- For n = 8:
  - 99: 0110 0011
  - 18: 0001 0010

*n*-bit signed integers Suppose that  $-2^{n-1} < x < 2^{n-1}$ First bit as the sign, n-1 bits for the value 99 = 64 + 32 + 2 + 118 = 16 + 2For n = 8: 99: 0110 0011 -18: 1001 0010

Any problems with this representation?

# **Two's Complement Representation**

n bit signed integers, first bit will still be the sign bit

Suppose that  $0 \le x < 2^{n-1}$ , *x* is represented by the binary representation of *x* Suppose that  $0 \le x \le 2^{n-1}$ , -*x* is represented by the binary representation of  $2^n - x$ 

**Key property:** Twos complement representation of any number y is equivalent to y, mod  $2^n$  so arithmetic works mod  $2^n$ 

#### Sign-Magnitude vs. Two's Complement

	-7	-6	-5	-4	-3	-2	-1	0	1	2	3	4	5	6	7
	1111	1110	1101	1100	1011	1010	1001	0000	0001	0010	0011	0100	0101	0110	0111
Sign-bit															
-8	-7	-6	-5	-4	-3	-2	-1	0	1	2	3	4	5	6	7
1000	1001	1010	1011	1100	1101	1110	1111	0000	0001	0010	0011	0100	0101	0110	0111
	Two's complement														

## **Two's Complement Representation**

- For  $0 < x \le 2^{n-1}$ , -x is represented by the binary representation of  $2^n x$ 
  - That is, the two's complement representation of any number y has the same value as y modulo  $2^n$ .

• To compute this: Flip the bits of x then add 1: - All 1's string is  $2^n - 1$ , so Flip the bits of  $x \equiv$  replace x by  $2^n - 1 - x$ Then add 1 to get  $2^n - x$ 

# **Basic Applications of mod**

- Hashing
- Pseudo random number generation
- Simple cipher

#### Scenario:

Map a small number of data values from a large domain  $\{0, 1, \dots, M - 1\}$  ...

...into a small set of locations  $\{0,1, ..., n-1\}$  so one can quickly check if some value is present

- hash(x) = x % p for p a prime close to n- or hash(x) = (ax + b)% p
- Depends on all of the bits of the data
  - helps avoid collisions due to similar values
  - need to manage them if they occur

**Linear Congruential method** 

$$x_{n+1} = (ax_n + c) \% m$$

#### Choose random $x_0$ , a, c, m and produce a long sequence of $x_n$ 's

- Caesar cipher, A = 1, B = 2, ...
  HELLO WORLD
- Shift cipher
  - -f(p) = (p + k) % 26 $-f^{-1}(p) = (p - k) \% 26$
- More general

-f(p) = (ap + b) % 26