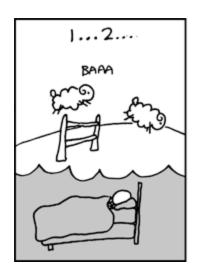
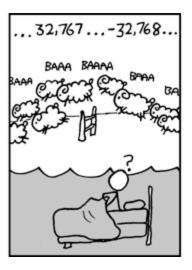
CSE 311: Foundations of Computing

Lecture 13: Number theory & modular arithmetic









Number Theory (and applications to computing)

Branch of Mathematics with direct relevance to computing

- Many significant applications
 - Cryptography
 - Hashing
 - Security

Important tool set

Modular Arithmetic

Arithmetic over a finite domain

In computing, almost all computations are over a finite domain

I'm ALIVE!

I'm ALIVE!

```
public class Test {
   final static int SEC IN YEAR = 364*24*60*60*100;
   public static void main(String args[]) {
       System.out.println(
          "I will be alive for at least " +
          SEC IN YEAR * 101 + " seconds."
      );
         ----jGRASP exec: java Test
        I will be alive for at least -186619904 seconds.
          ----jGRASP: operation complete.
```

Divisibility

Definition: "a divides b"

For
$$a \in \mathbb{Z}$$
, $b \in \mathbb{Z}$ with $a \neq 0$:
 $a \mid b \leftrightarrow \exists k \in \mathbb{Z} \ (b = ka)$

Check Your Understanding. Which of the following are true?

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Check Your Understanding. Which of the following are true?

$$25 \mid 5$$
 iff $\exists k. 5 = 25k$



$$5 \mid 0$$

$$\exists k. 0 = 5k$$

$$0 \mid 5$$
 iff $\exists k. 5 = 0k$

$$5 \mid 0$$
 $3 \mid 2$
0 = 5k iff 3k, 2 = 3k

$$2 \mid 3$$
 iff $\exists k. 3 = 2k$

Division Theorem

Division Theorem

For $a \in \mathbb{Z}$, $d \in \mathbb{Z}$ with d > 0there exist *unique* integers q, r with $0 \le r < d$ such that a = dq + r.

To put it another way, if we divide d into a, we get a unique quotient $q = a \operatorname{div} d$ and non-negative remainder r = a % d

Note: $r \ge 0$ even if a < 0. Not quite the same as in Java

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```
public class Test2 {
   public static void main(String args[]) {
      int a = -5;
      int d = 2;
      System.out.println(a % d);
   }
      Note: r ≥ 0 even if a < 0.</pre>
```

Not quite the same as in Java

Modular Arithmetic

Definition: "a is congruent to b modulo m"

For
$$a, b, m \in \mathbb{Z}$$
 with $m > 0$
 $a \equiv b \pmod{m} \leftrightarrow m \mid (a - b)$

Check Your Understanding. What do each of these mean? When are they true?

$$x \equiv 0 \pmod{2}$$

$$-1 \equiv 19 \pmod{5}$$

$$y \equiv 2 \pmod{7}$$

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Check Your Understanding. What do each of these mean? When are they true?

$$x \equiv 0 \pmod{2}$$

This statement is the same as saying "x is even"; so, any x that is even (including negative even numbers) will work.

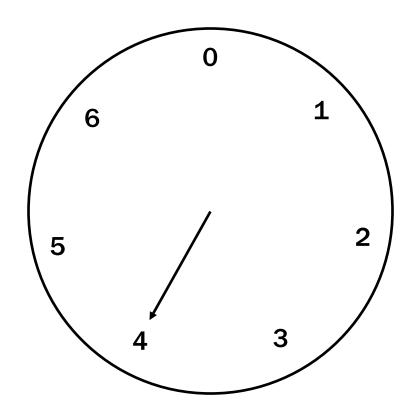
$$-1 \equiv 19 \pmod{5}$$

This statement is true. 19 - (-1) = 20 which is divisible by 5

$$y \equiv 2 \pmod{7}$$

This statement is true for y in { ..., -12, -5, 2, 9, 16, ...}. In other words, all y of the form 2+7k for k an integer.

Arithmetic, mod 7



Modular Arithmetic: A Property

Let a, b, m be integers with m > 0. Then, $a \equiv b \pmod{m}$ if and only if a % m = b % m.

Suppose that $a \equiv b \pmod{m}$.

Suppose that a % m = b % m.

Modular Arithmetic: A Property

Let a, b, m be integers with m > 0. Then, $a \equiv b \pmod{m}$ if and only if a % m = b % m. Suppose that $a \equiv b \pmod{m}$. Then, $m \mid (a - b)$ by definition of congruence. So, a - b = km for some integer k by definition of divides. Therefore, a = b + km. Taking both sides modulo m we get: a % m = (b + km)% m = b % m.Suppose that a % m = b % m. By the division theorem, a = mq + (a % m) and

$$b = ms + (b \% m) \text{ for some integers } q,s.$$
Then, $a - b = (mq + (a \% m)) - (ms + (b \% m))$

$$= m(q - s) + (a \% m - b \% m)$$

$$= m(q - s) \text{ since } a \% m = b \% m$$
Therefore, $m \mid (a - b) \text{ and so } a \equiv b \text{ (mod } m).$

The % m function vs the $\equiv \pmod{m}$ predicate

- What we have just shown
 - The % m function takes any $a \in \mathbb{Z}$ and maps it to a remainder a % $m \in \{0,1,...,m-1\}$.
 - Imagine grouping together all integers that have the same value of the % m function That is, the same remainder in $\{0,1,\ldots,m-1\}$.
 - The $\equiv \pmod{m}$ predicate compares $a, b \in \mathbb{Z}$. It is true if and only if the % m function has the same value on a and on b.

That is, a and b are in the same group.

Modular Arithmetic: Basic Property

```
Let m be a positive integer.

If a \equiv b \pmod{m} and b \equiv c \pmod{m},

then a \equiv c \pmod{m}
```

Modular Arithmetic: Basic Property

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Let m be a positive integer.

If a \equiv b \pmod{m} and b \equiv c \pmod{m},

then a \equiv c \pmod{m}
```

Suppose that $a \equiv b \pmod{m}$ and $b \equiv c \pmod{m}$. Then, by the previous property, we have a % m = b % m and b % m = c % m.

Putting these together, we have a % m = c % m, which says that $a \equiv c \pmod{m}$, by definition.

So "≡" behaves like "=" in that sense. And that is not the only similarity...

Modular Arithmetic: Addition Property

Let m be a positive integer. If $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$, then $a + c \equiv b + d \pmod{m}$

Modular Arithmetic: Addition Property

Let m be a positive integer. If $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$, then $a + c \equiv b + d \pmod{m}$

Suppose that $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$. Unrolling definitions gives us some k such that a - b = km, and some j such that c - d = jm.

Adding the equations together gives us (a+c)-(b+d)=m(k+j). Now, re-applying the definition of congruence gives us $a+c\equiv b+d\pmod{m}$.

Modular Arithmetic: Multiplication Property

Let m be a positive integer. If $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$, then $ac \equiv bd \pmod{m}$

Modular Arithmetic: Multiplication Property

Let m be a positive integer. If $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$, then $ac \equiv bd \pmod{m}$

Suppose that $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$. Unrolling definitions gives us some k such that a - b = km, and some j such that c - d = jm.

Then, a = km + b and c = jm + d. Multiplying both together gives us $ac = (km + b)(jm + d) = kjm^2 + kmd + bjm + bd$.

Re-arranging gives us ac - bd = m(kjm + kd + bj). Using the definition of congruence gives us $ac \equiv bd \pmod{m}$.

Lecture 13 Activity

You will be assigned to breakout rooms. Please:

- Introduce yourself
- Choose someone to share their screen, showing this PDF
- Consider the statement:

```
For all a, b, c, m \in \mathbb{Z}, m > 0 one has a \equiv b \pmod{m} \rightarrow a + c \equiv b + c \pmod{m}.
```

- Discuss what the statement means.
- Prove the statement.

Definition: "a is congruent to b modulo m"

For
$$a, b, m \in \mathbb{Z}$$
 with $m > 0$
 $a \equiv b \pmod{m} \leftrightarrow m \mid (a - b)$

Fill out the poll everywhere for Activity Credit!
Go to pollev.com/philipmg and login with your UW identity

Definition: "a divides b"

For $a \in \mathbb{Z}$, $b \in \mathbb{Z}$ with $a \neq 0$: $a \mid b \leftrightarrow \exists k \in \mathbb{Z} \ (b = ka)$

Let n be an integer.

Prove that
$$n^2 \equiv 0 \pmod{4}$$
 or $n^2 \equiv 1 \pmod{4}$

Let's start by looking a a small example:

$$0^2 = 0 \equiv 0 \pmod{4}$$

$$1^2 = 1 \equiv 1 \pmod{4}$$

$$2^2 = 4 \equiv 0 \pmod{4}$$

$$3^2 = 9 \equiv 1 \pmod{4}$$

$$4^2 = 16 \equiv 0 \pmod{4}$$

Let n be an integer.

Prove that $n^2 \equiv 0 \pmod{4}$ or $n^2 \equiv 1 \pmod{4}$

Case 1 (n is even):

Let's start by looking a a small example:

$$0^2 = 0 \equiv 0 \pmod{4}$$

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It looks like

$$n \equiv 0 \pmod{2} \rightarrow n^2 \equiv 0 \pmod{4}$$
, and

$$n \equiv 1 \pmod{2} \rightarrow n^2 \equiv 1 \pmod{4}$$
.

```
Let n be an integer.
```

Prove that
$$n^2 \equiv 0 \pmod{4}$$
 or $n^2 \equiv 1 \pmod{4}$

Case 1 (n is even):

Suppose n is even.

Then, n = 2k for some integer k.

So, $n^2 = (2k)2 = 4k^2$.

So, by definition of congruence,

we have $n^2 \equiv 0 \pmod{4}$.

Let's start by looking a a small example:

 $0^2 = 0 \equiv 0 \pmod{4}$

 $1^2 = 1 \equiv 1 \pmod{4}$

 $2^2 = 4 \equiv 0 \pmod{4}$

 $3^2 = 9 \equiv 1 \pmod{4}$

 $4^2 = 16 \equiv 0 \pmod{4}$

It looks like

 $n \equiv 0 \pmod{2} \rightarrow n^2 \equiv 0 \pmod{4}$, and $n \equiv 1 \pmod{2} \rightarrow n^2 \equiv 1 \pmod{4}$.

Let n be an integer.

Prove that $n^2 \equiv 0 \pmod{4}$ or $n^2 \equiv 1 \pmod{4}$

Case 1 (n is even): Done.

Case 2 (n is odd):

Let's start by looking a a small example:

$$0^2 = 0 \equiv 0 \pmod{4}$$

$$1^2 = 1 \equiv 1 \pmod{4}$$

$$2^2 = 4 \equiv 0 \pmod{4}$$

$$3^2 = 9 \equiv 1 \pmod{4}$$

$$4^2 = 16 \equiv 0 \pmod{4}$$

It looks like

$$n \equiv 0 \pmod{2} \rightarrow n^2 \equiv 0 \pmod{4}$$
, and $n \equiv 1 \pmod{2} \rightarrow n^2 \equiv 1 \pmod{4}$.

```
Let n be an integer.
Prove that n^2 \equiv 0 \pmod 4 or n^2 \equiv 1 \pmod 4
```

```
Let's start by looking a a small example:
Case 1 (n is even): Done.
                                                         0^2 = 0 \equiv 0 \pmod{4}
                                                         1^2 = 1 \equiv 1 \pmod{4}
Case 2 (n is odd):
                                                         2^2 = 4 \equiv 0 \pmod{4}
    Suppose n is odd.
                                                         3^2 = 9 \equiv 1 \pmod{4}
     Then, n = 2k + 1 for some integer k.
                                                         4^2 = 16 \equiv 0 \pmod{4}
    So, n^2 = (2k + 1)^2
         =4k^2+4k+1
                                            It looks like
          =4(k^2+k)+1.
                                                 n \equiv 0 \pmod{2} \rightarrow n^2 \equiv 0 \pmod{4}, and
    So, by the earlier property of mod, n \equiv 1 \pmod{2} \rightarrow n^2 \equiv 1 \pmod{4}.
    we have n^2 \equiv 1 \pmod{4}.
```

Result follows by "proof by cases": n is either even or not even (odd)

n-bit Unsigned Integer Representation

• Represent integer x as sum of powers of 2:

If
$$\sum_{i=0}^{n-1} b_i 2^i$$
 where each $b_i \in \{0,1\}$
then representation is $b_{n-1}...b_2$ b_1 b_0

$$99 = 64 + 32 + 2 + 1$$

 $18 = 16 + 2$

• For n = 8:

99: 0110 0011

18: 0001 0010

Sign-Magnitude Integer Representation

n-bit signed integers

Suppose that $-2^{n-1} < x < 2^{n-1}$ First bit as the sign, n-1 bits for the value

$$99 = 64 + 32 + 2 + 1$$

 $18 = 16 + 2$

For n = 8:

99: 0110 0011

-18: 1001 0010

Any problems with this representation?

Two's Complement Representation

n bit signed integers, first bit will still be the sign bit

```
Suppose that 0 \le x < 2^{n-1}, x is represented by the binary representation of x. Suppose that 0 \le x \le 2^{n-1}, -x is represented by the binary representation of 2^n - x.
```

Key property: Twos complement representation of any number y is equivalent to y, $mod 2^n$ so arithmetic works $mod 2^n$

$$99 = 64 + 32 + 2 + 1$$

 $18 = 16 + 2$

For n = 8:

99: 0110 0011 -18: 1110 1110

Sign-Magnitude vs. Two's Complement

-7 -5 -3 -2 -1 Sign-bit

Two's complement

Two's Complement Representation

- For $0 < x \le 2^{n-1}$, -x is represented by the binary representation of $2^n x$
 - That is, the two's complement representation of any number y has the same value as y modulo 2^n .

- To compute this: Flip the bits of x then add 1:
 - All 1's string is $2^n 1$, so

 Flip the bits of $x \equiv \text{replace } x \text{ by } 2^n 1 x$ Then add 1 to get $2^n x$

Basic Applications of mod

- Hashing
- Pseudo random number generation
- Simple cipher

Hashing

Scenario:

Map a small number of data values from a large domain $\{0, 1, ..., M - 1\}$...

...into a small set of locations $\{0,1,\ldots,n-1\}$ so one can quickly check if some value is present

- hash(x) = x % p for p a prime close to n
 - or hash(x) = (ax + b)% p
- Depends on all of the bits of the data
 - helps avoid collisions due to similar values
 - need to manage them if they occur

Pseudo-Random Number Generation

Linear Congruential method

$$x_{n+1} = (ax_n + c) \% m$$

Choose random x_0 , a, c, m and produce a long sequence of x_n 's

Simple Ciphers

- Caesar cipher, A = 1, B = 2, . . .
 - HELLO WORLD
- Shift cipher

$$- f(p) = (p + k) \% 26$$

$$-f^{-1}(p) = (p - k) \% 26$$

More general

$$- f(p) = (ap + b) \% 26$$