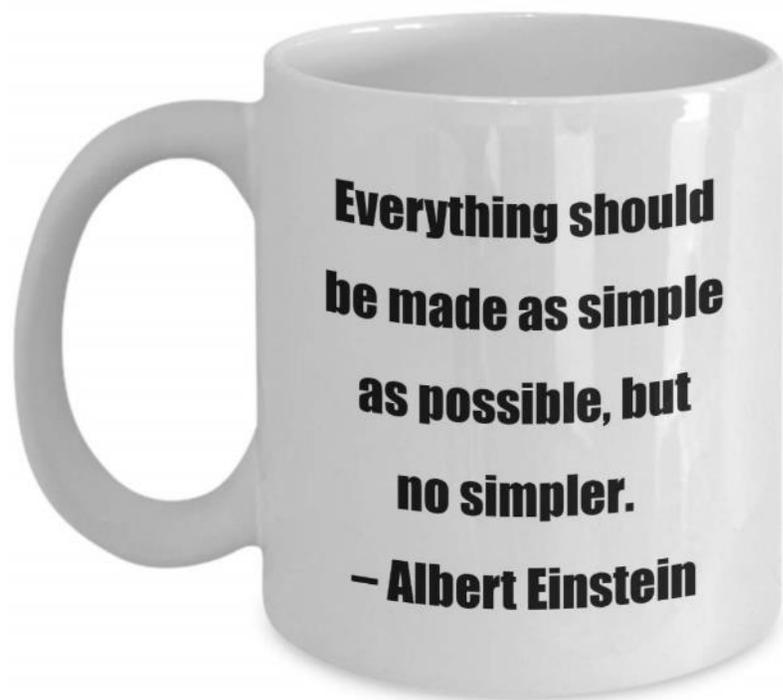


# CSE 311: Foundations of Computing

---

## Lecture 23: FSMs with Output, Minimization





# Vending Machine

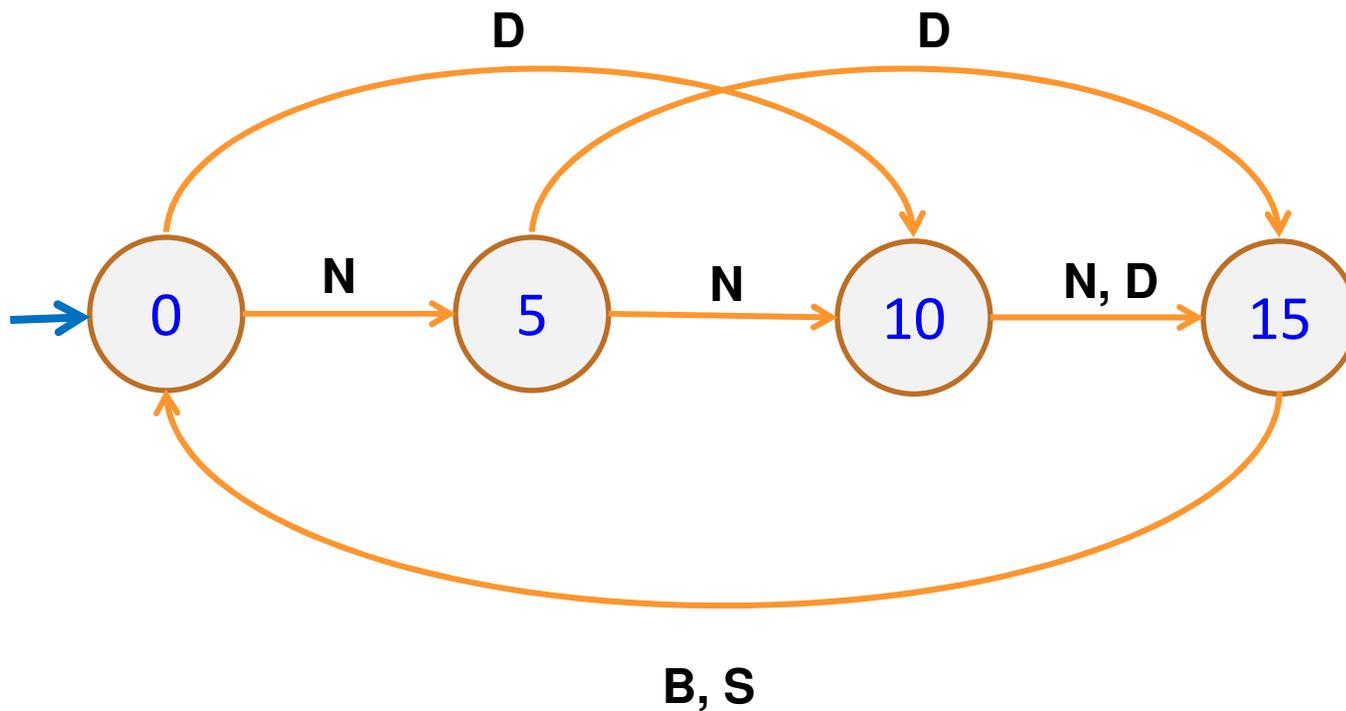


Enter 15 cents in dimes or nickels  
Press S or B for a candy bar



# Vending Machine, v0.1

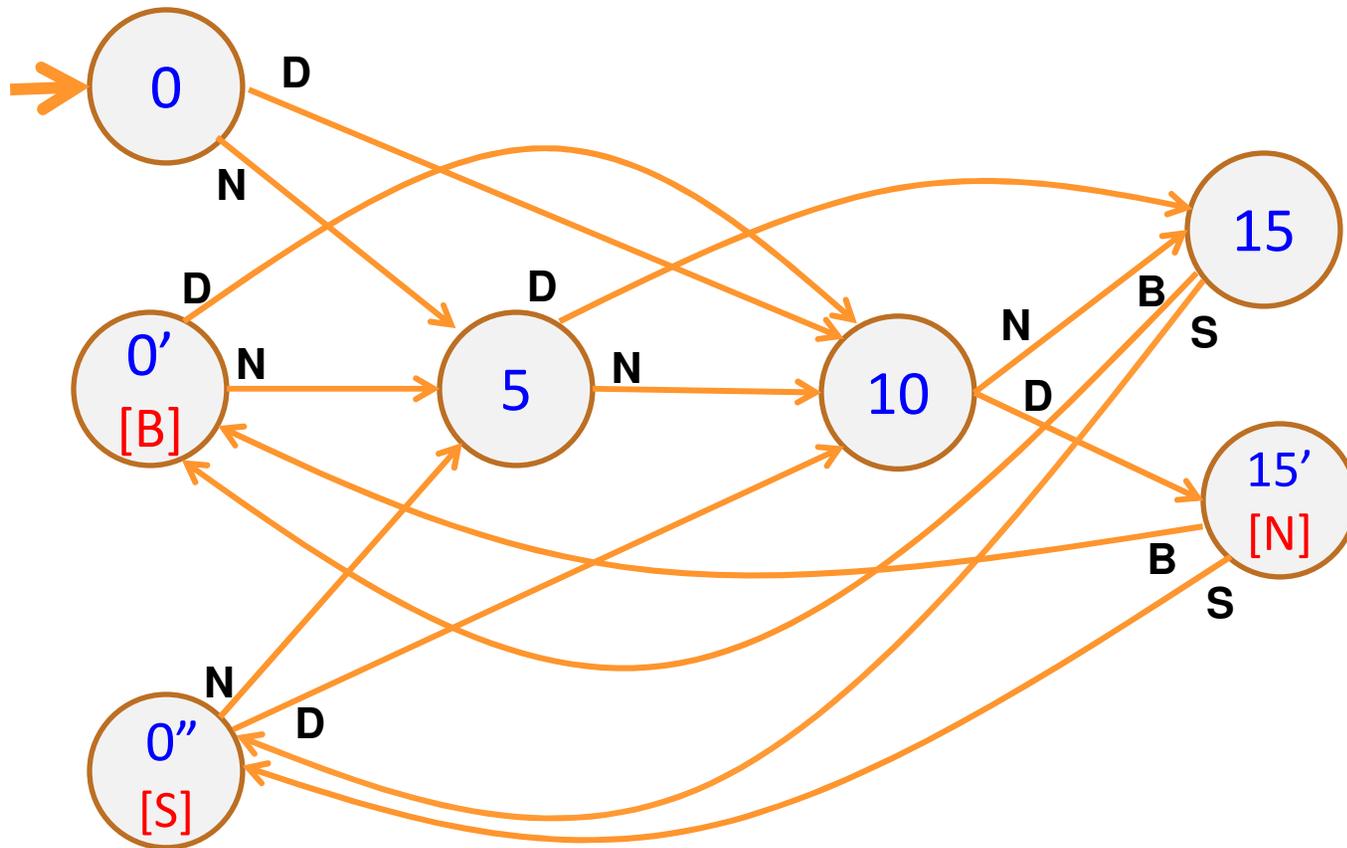
---



Basic transitions on **N** (nickel), **D** (dime), **B** (butterfinger), **S** (snickers)

# Vending Machine, v0.2

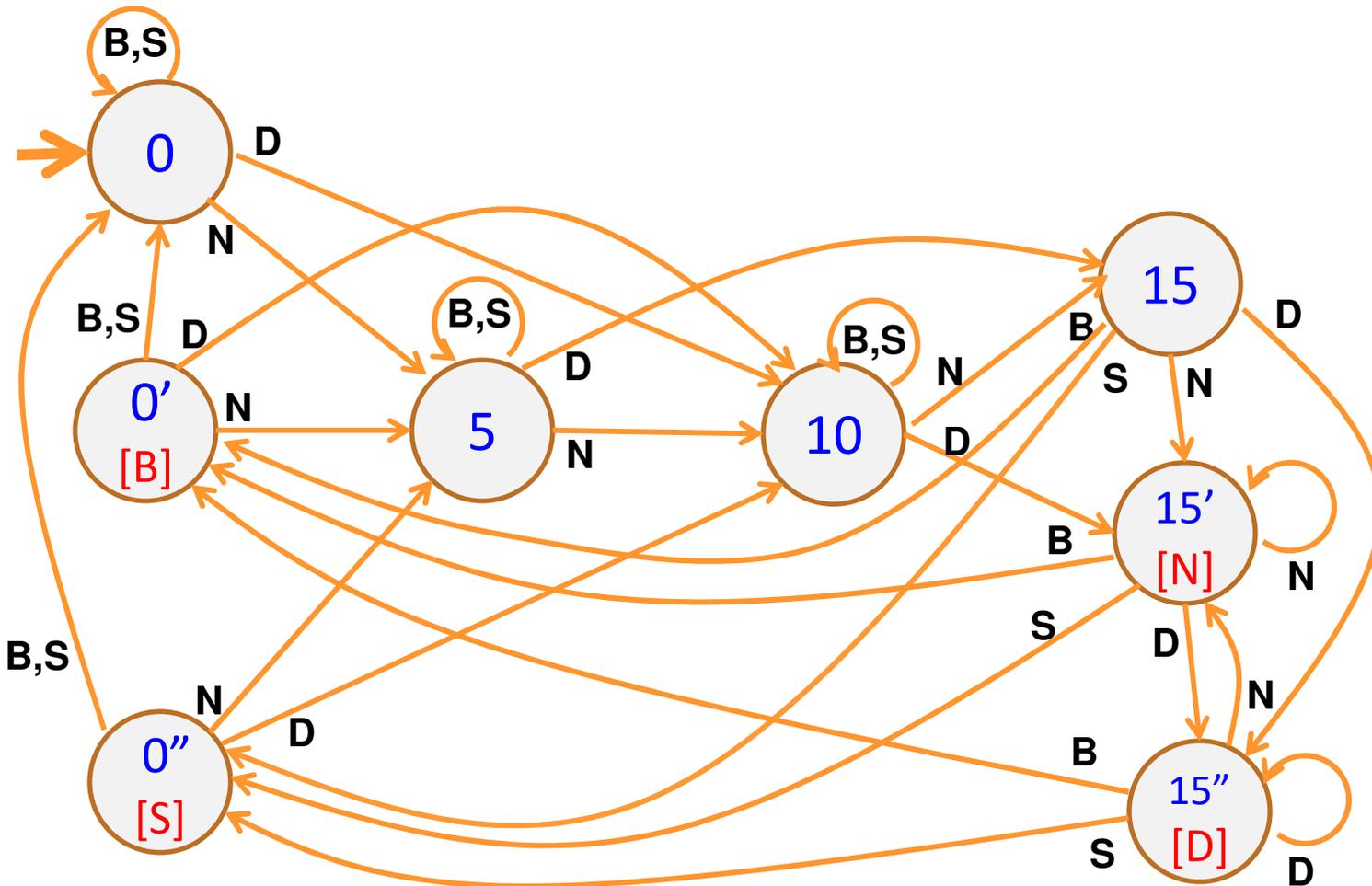
---



Adding output to states: **N** – Nickel, **S** – Snickers, **B** – Butterfinger

# Vending Machine, v1.0

---



Adding additional “unexpected” transitions to cover all symbols for each state

# State Minimization

---

- **Many different FSMs (DFAs) for the same problem**
- **Take a given FSM and try to reduce its state set by combining states**
  - **Algorithm will always produce the unique minimal equivalent machine (up to renaming of states) but we won't prove this**

# State Minimization Algorithm

---

- Put states into groups
- Try to find groups that can be collapsed into one state
  - states can keep track of information that isn't necessary to determine whether to accept or reject
- Group states together until we can *prove* that collapsing them can change the accept/reject result
  - find a specific string  $x$  such that:
    - starting from state A, following edges according to  $x$  ends in accept
    - starting from state B, following edges according to  $x$  ends in reject
  - (algorithm below could be modified to show these strings)

# State Minimization Algorithm

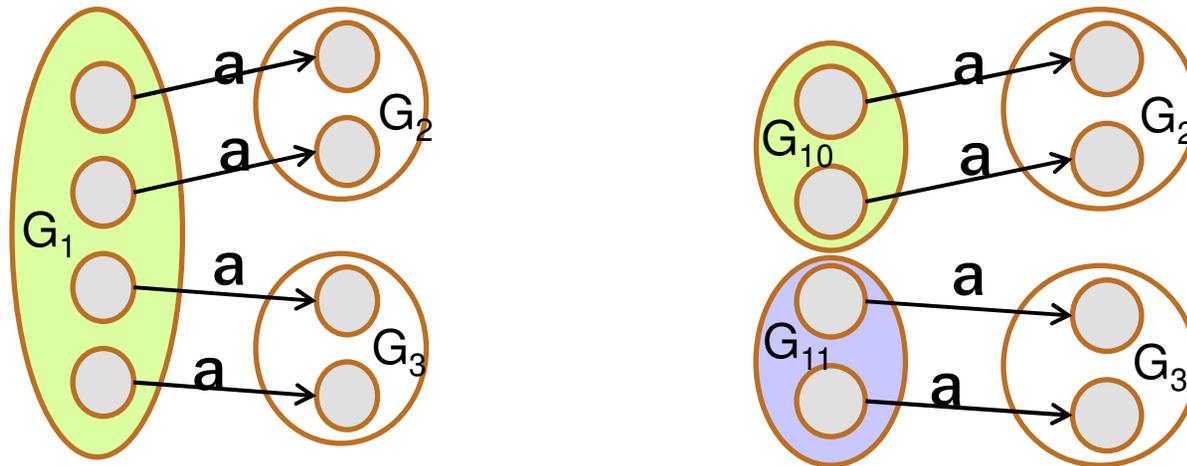
---

1. Put states into groups based on their outputs (whether they accept or reject)

# State Minimization Algorithm

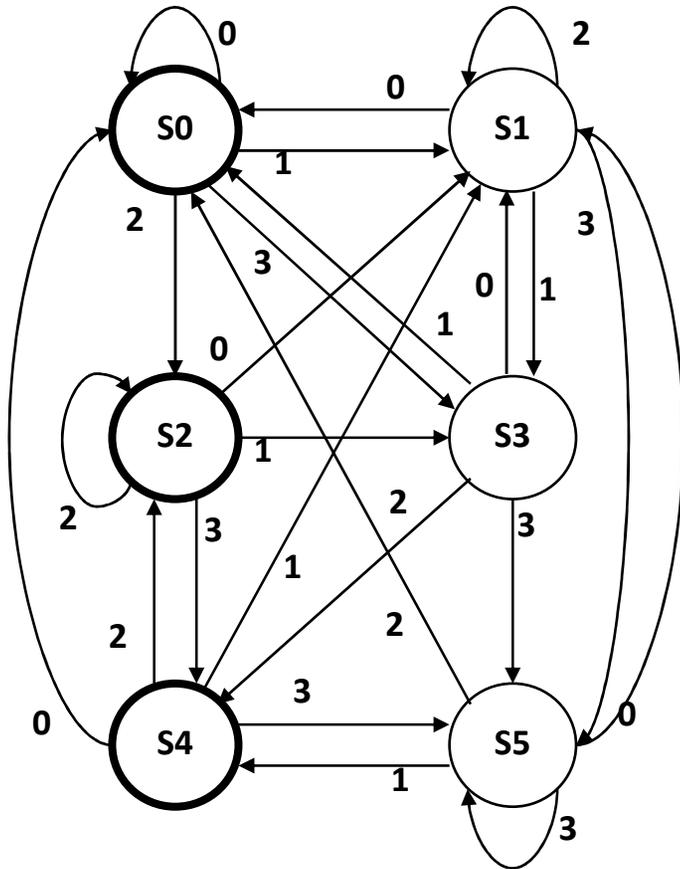
---

1. Put states into groups based on their outputs (whether they accept or reject)
2. Repeat the following until no change happens
  - a. If there is a symbol **a** so that not all states in a group **G** agree on which group **a** leads to, split **G** into smaller groups based on which group the states go to on **a**



3. Finally, convert groups to states

# State Minimization Example

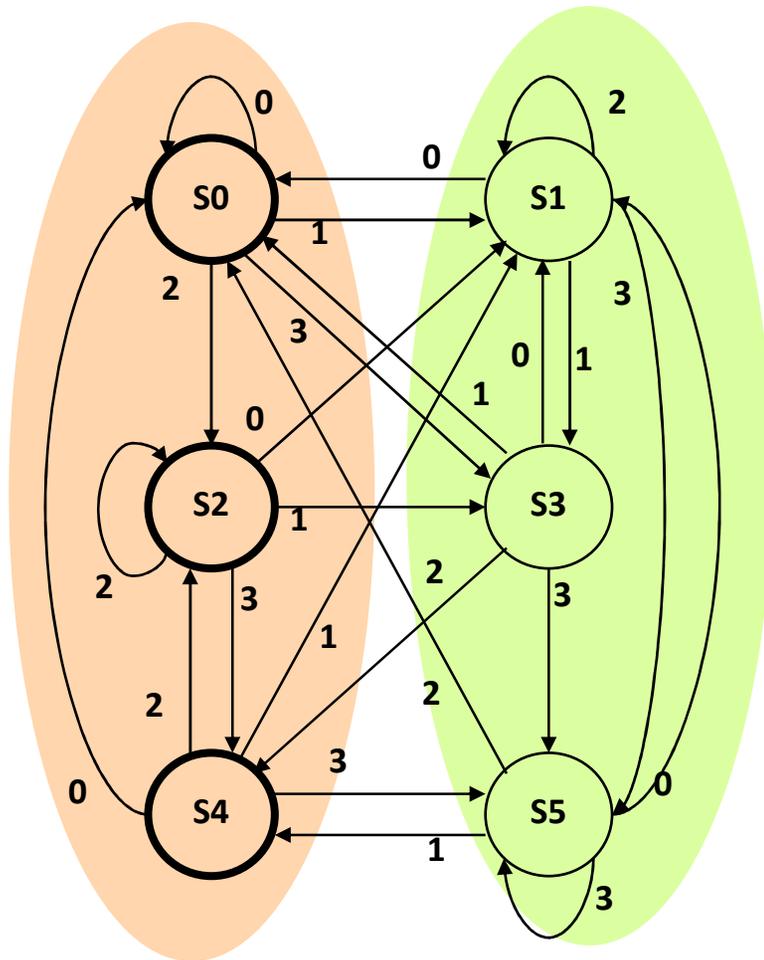


present state	next state				output
	0	1	2	3	
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
S4	S0	S1	S2	S5	1
S5	S1	S4	S0	S5	0

state transition table

Put states into groups based on their outputs (or whether they accept or reject)

# State Minimization Example

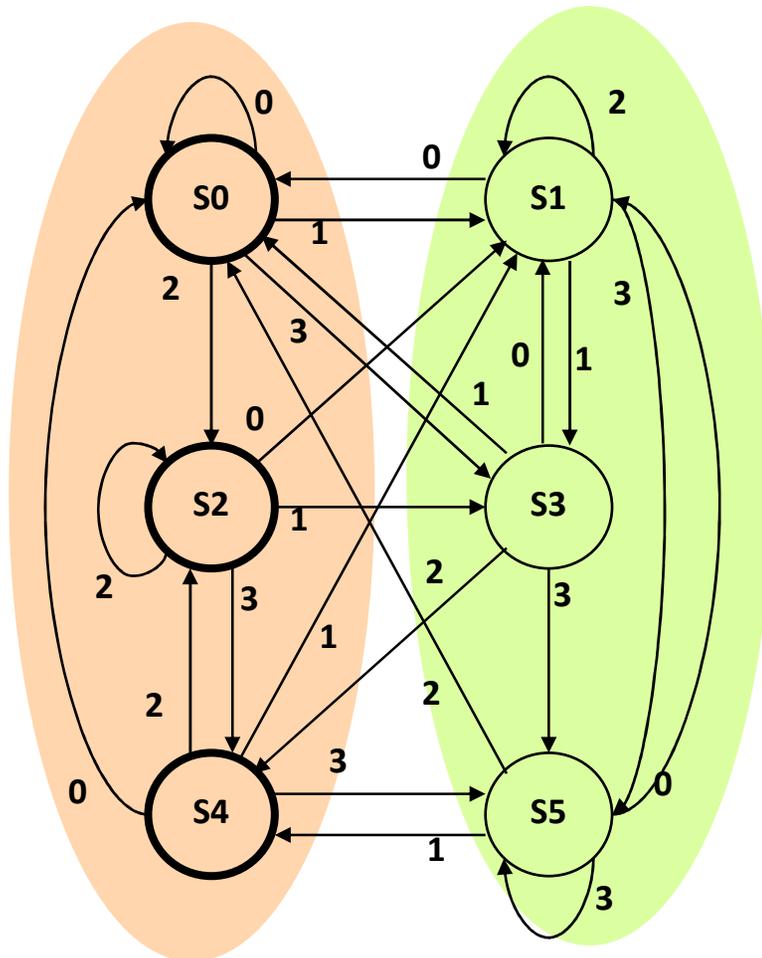


present state	next state				output
	0	1	2	3	
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
S4	S0	S1	S2	S5	1
S5	S1	S4	S0	S5	0

state transition table

Put states into groups based on their outputs (or whether they accept or reject)

# State Minimization Example



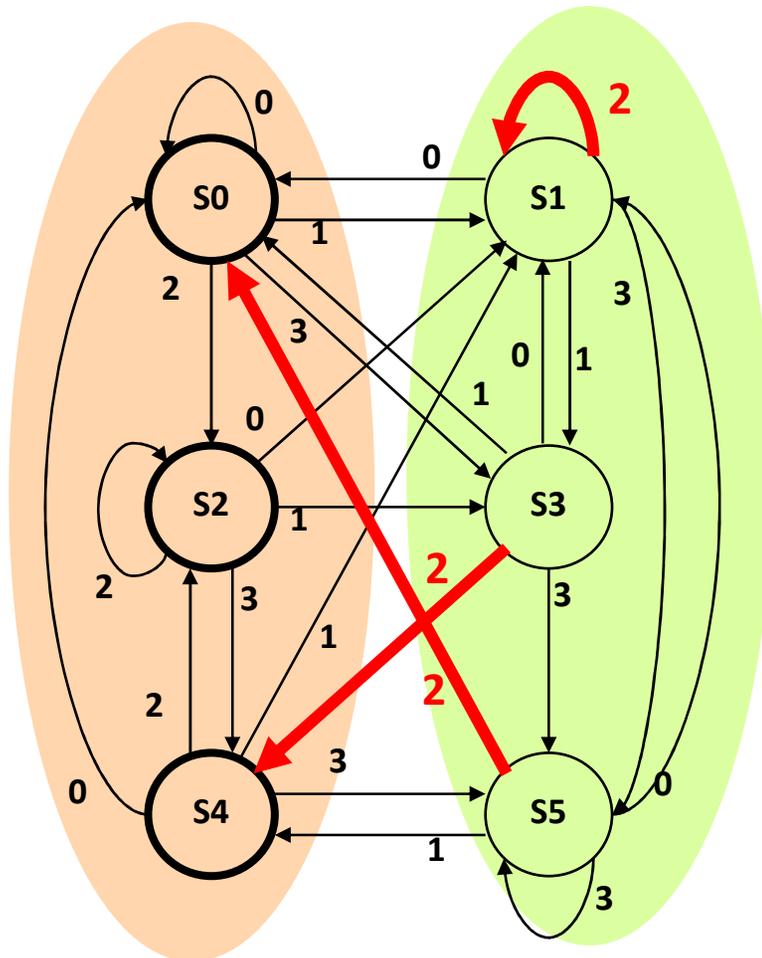
present state	next state				output
	0	1	2	3	
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
S4	S0	S1	S2	S5	1
S5	S1	S4	S0	S5	0

state transition table

Put states into groups based on their outputs (or whether they accept or reject)

If there is a symbol **a** so that not all states in a group **G** agree on which group **a** leads to, split **G** based on which group the states go to on **a**

# State Minimization Example



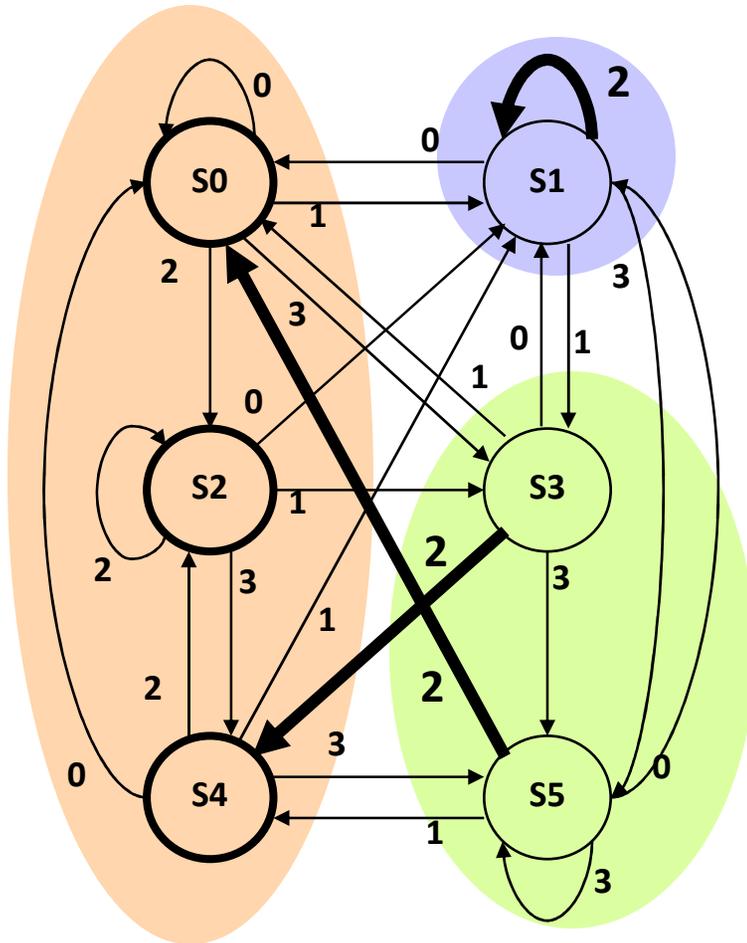
present state	next state				output
	0	1	2	3	
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
S4	S0	S1	S2	S5	1
S5	S1	S4	S0	S5	0

state transition table

Put states into groups based on their outputs (or whether they accept or reject)

If there is a symbol **a** so that not all states in a group **G** agree on which group **a** leads to, split **G** based on which group the states go to on **a**

# State Minimization Example



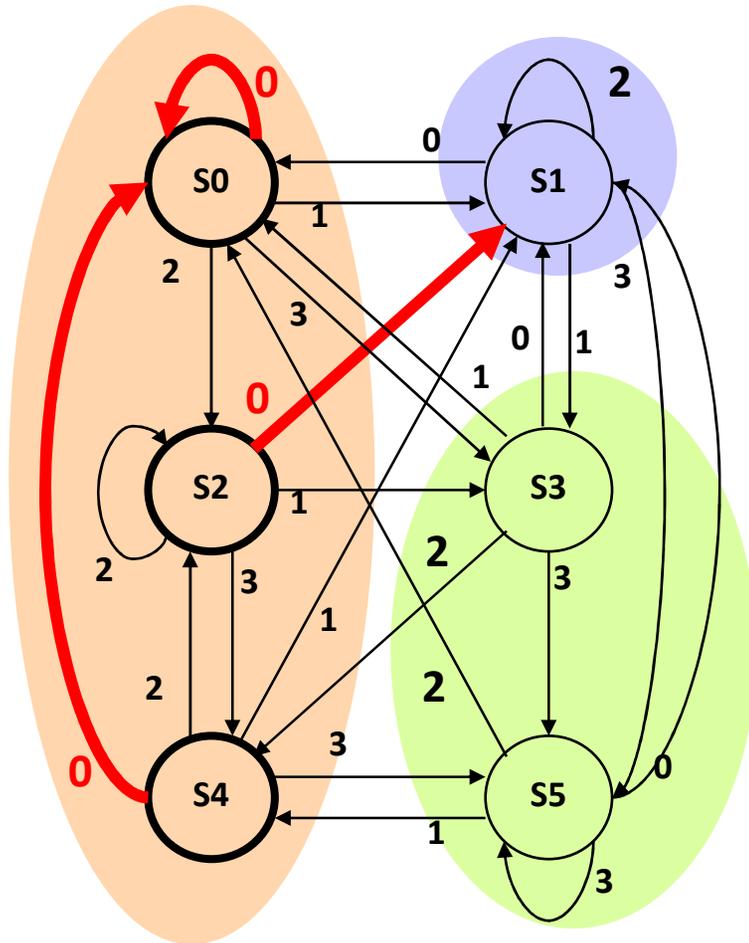
present state	next state				output
	0	1	2	3	
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
S4	S0	S1	S2	S5	1
S5	S1	S4	S0	S5	0

state transition table

Put states into groups based on their outputs (or whether they accept or reject)

If there is a symbol **a** so that not all states in a group **G** agree on which group **a** leads to, split **G** based on which group the states go to on **a**

# State Minimization Example



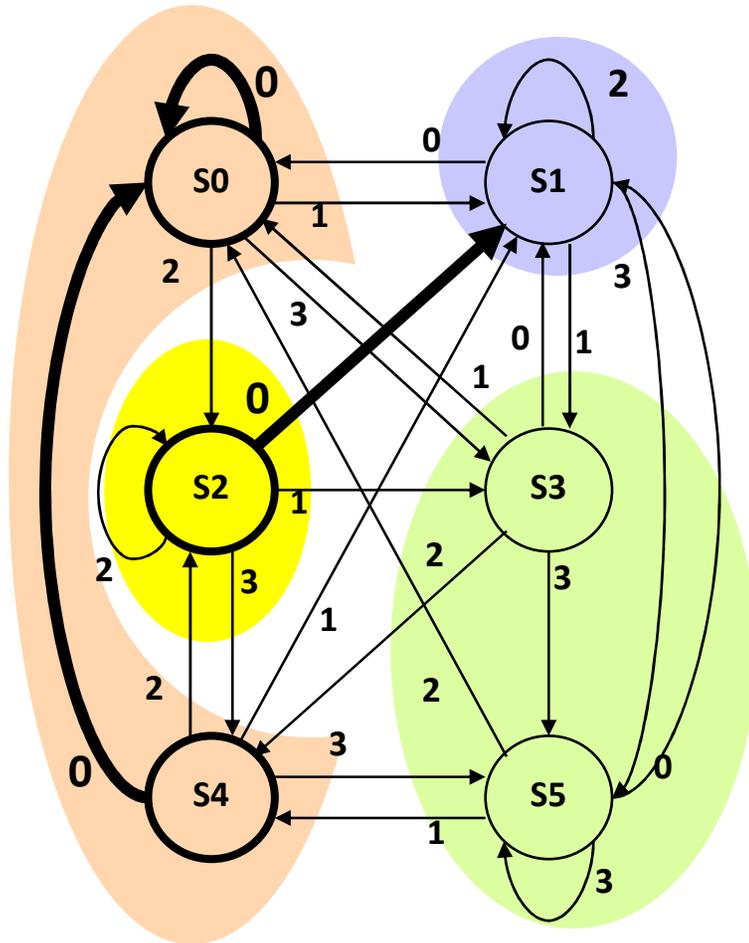
present state	next state				output
	0	1	2	3	
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
S4	S0	S1	S2	S5	1
S5	S1	S4	S0	S5	0

state transition table

Put states into groups based on their outputs (or whether they accept or reject)

If there is a symbol **a** so that not all states in a group **G** agree on which group **a** leads to, split **G** based on which group the states go to on **a**

# State Minimization Example



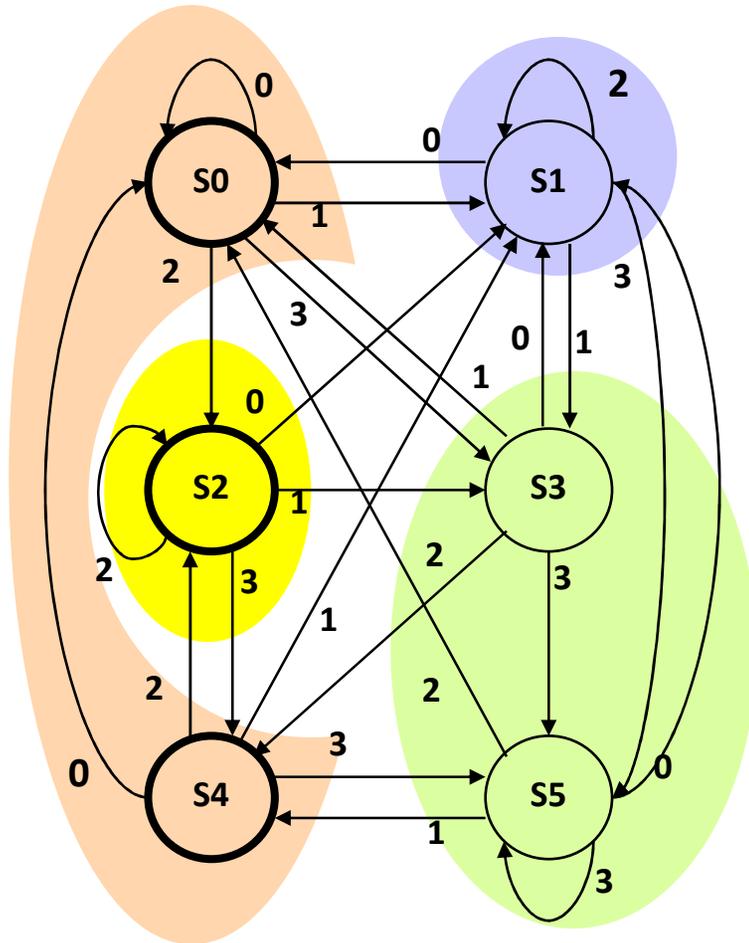
present state	next state				output
	0	1	2	3	
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
S4	S0	S1	S2	S5	1
S5	S1	S4	S0	S5	0

state transition table

Put states into groups based on their outputs (or whether they accept or reject)

If there is a symbol **a** so that not all states in a group **G** agree on which group **a** leads to, split **G** based on which group the states go to on **a**

# State Minimization Example



present state	next state				output
	0	1	2	3	
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S5	0
S2	S1	S3	S2	S4	1
S3	S1	S0	S4	S5	0
S4	S0	S1	S2	S5	1
S5	S1	S4	S0	S5	0

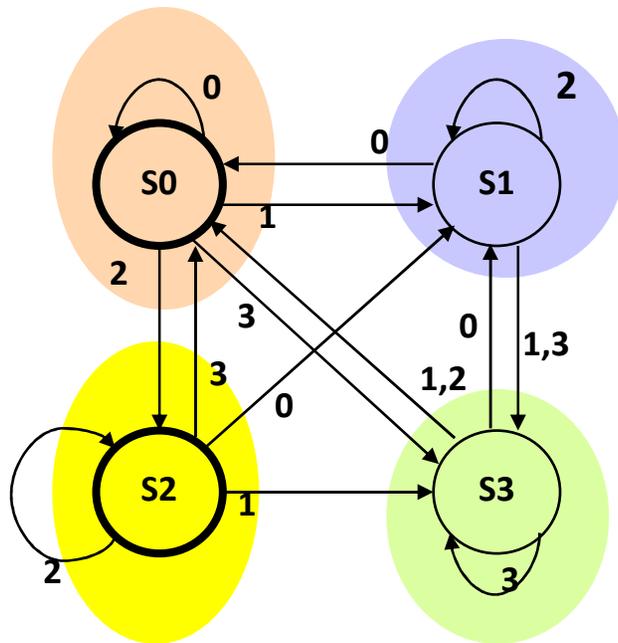
state transition table

Finally convert groups to states:

Can combine states S0-S4 and S3-S5.

In table replace all S4 with S0 and all S5 with S3

# Minimized Machine

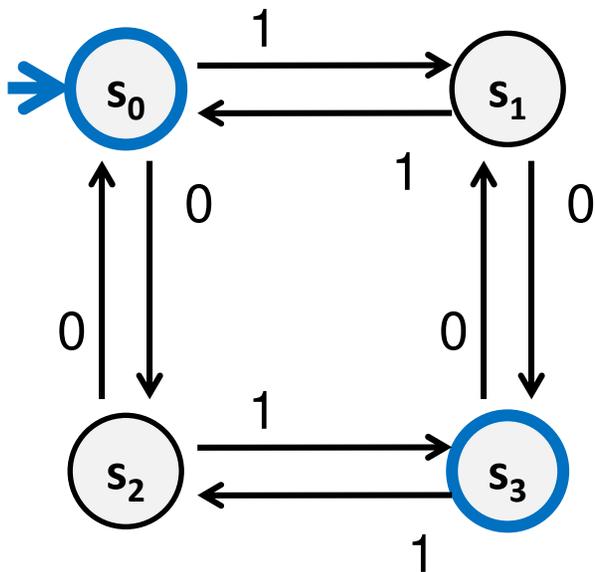


present state	next state				output
	0	1	2	3	
S0	S0	S1	S2	S3	1
S1	S0	S3	S1	S3	0
S2	S1	S3	S2	S0	1
S3	S1	S0	S0	S3	0

state transition table

# A Simpler Minimization Example

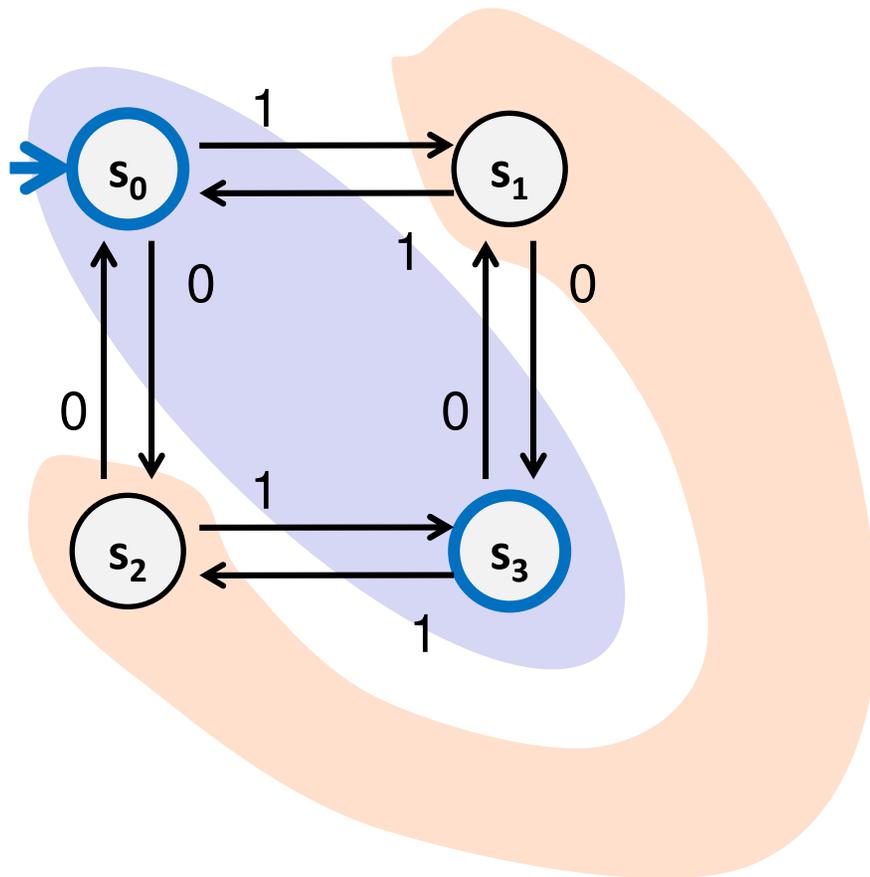
---



The set of all binary strings with  $\#$  of 1's  $\equiv$   $\#$  of 0's (mod 2).

# A Simpler Minimization Example

---

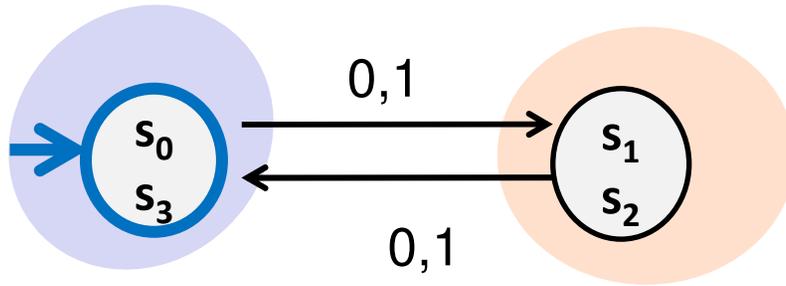


**Split states into  
accept/reject groups**

**Every symbol causes  
the DFA to go from one  
group to the other so  
neither group needs to  
be split**

# Minimized DFA

---



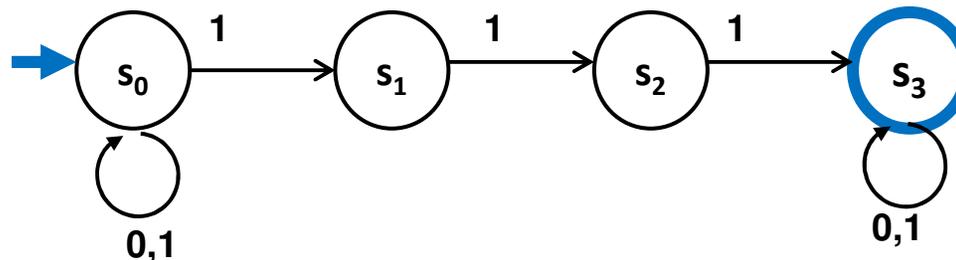
The set of all binary strings with # of 1's  $\equiv$  # of 0's (mod 2).

= The set of all binary strings with even length.

# Nondeterministic Finite Automata (NFA)

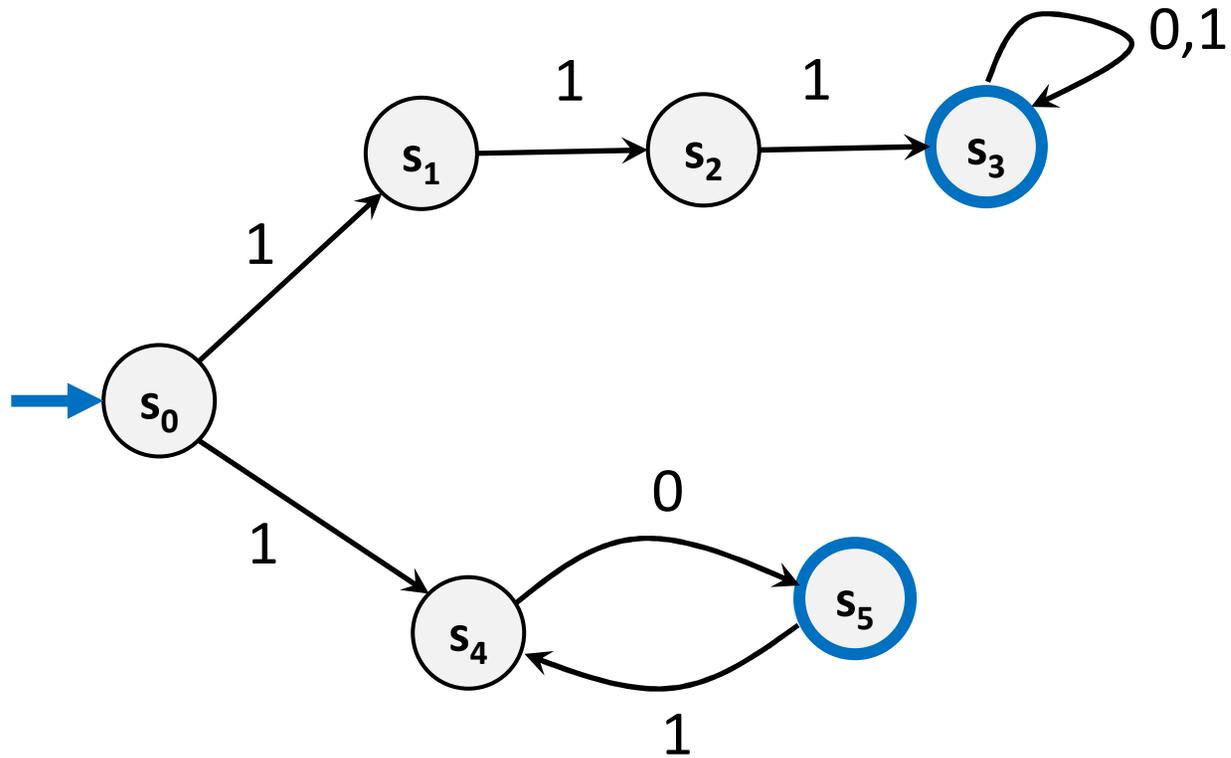
---

- Graph with start state, final states, edges labeled by symbols (like DFA) but
  - Not required to have exactly 1 edge out of each state labeled by each symbol— can have 0 or  $>1$
  - Also can have edges labeled by empty string  $\epsilon$
- **Definition:**  $x$  is in the language recognized by an NFA if and only if some valid execution of the machine gets to an accept state



## Consider This NFA

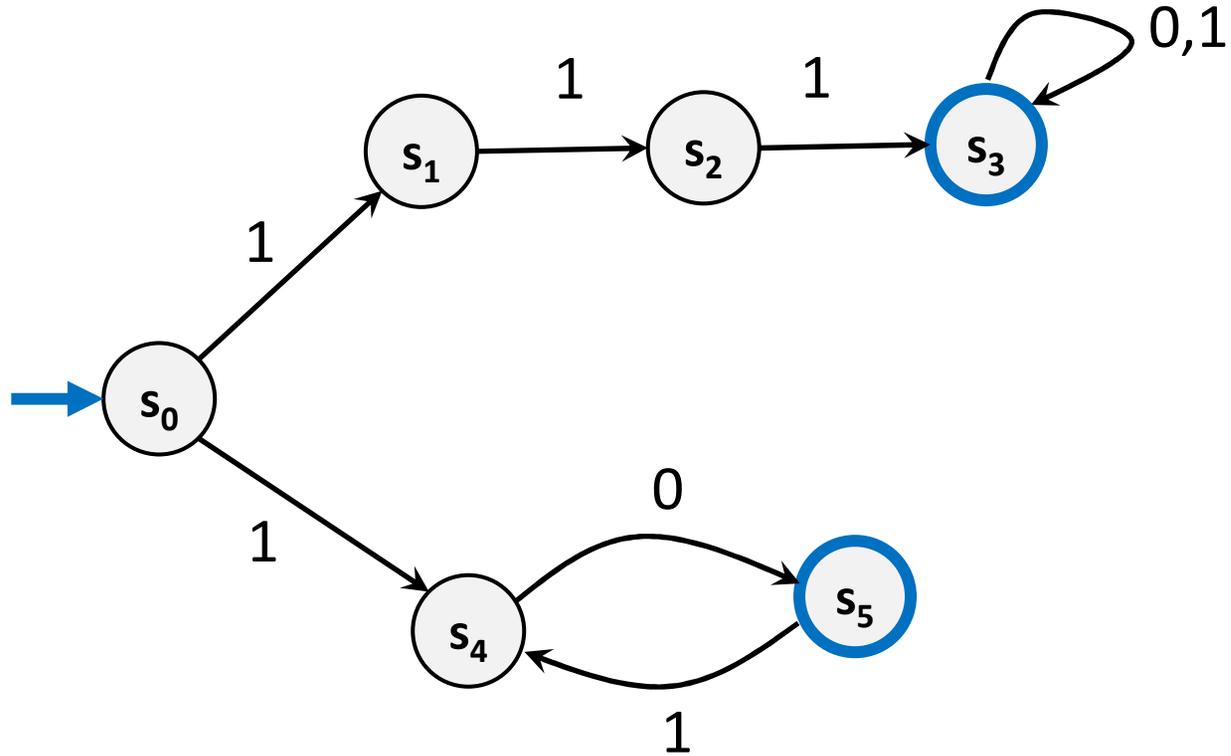
---



What language does this NFA accept?

## Consider This NFA

---

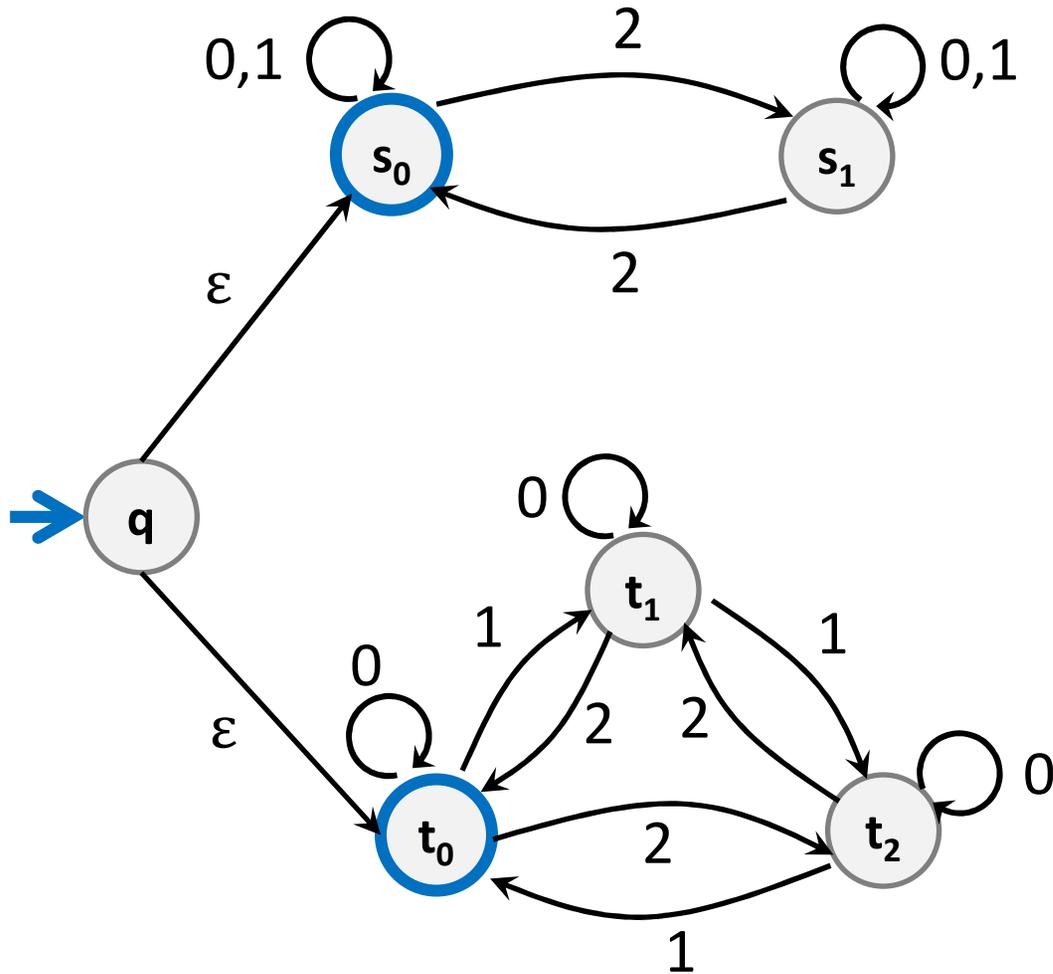


What language does this NFA accept?

$$10(10)^* \cup 111(0 \cup 1)^*$$

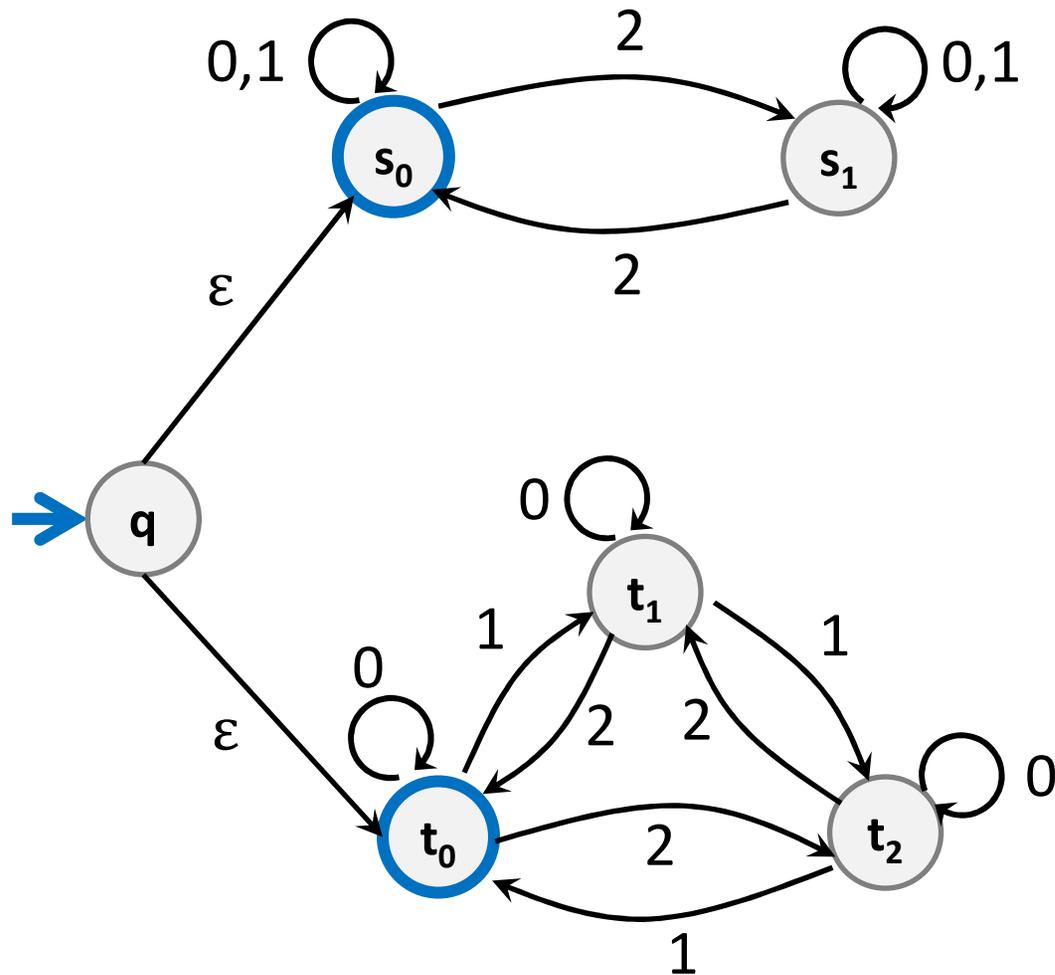
# NFA $\epsilon$ -moves

---



# NFA $\epsilon$ -moves

Strings over  $\{0,1,2\}$  w/even # of 2's OR sum to 0 mod 3

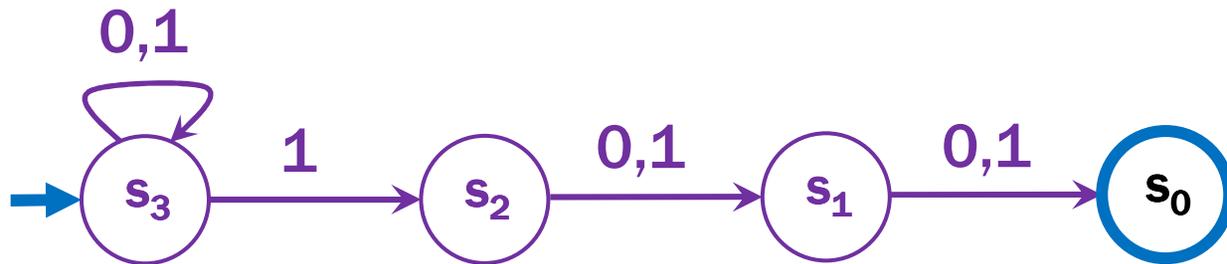


**NFA for set of binary strings with a 1 in the 3<sup>rd</sup> position from the end**

---

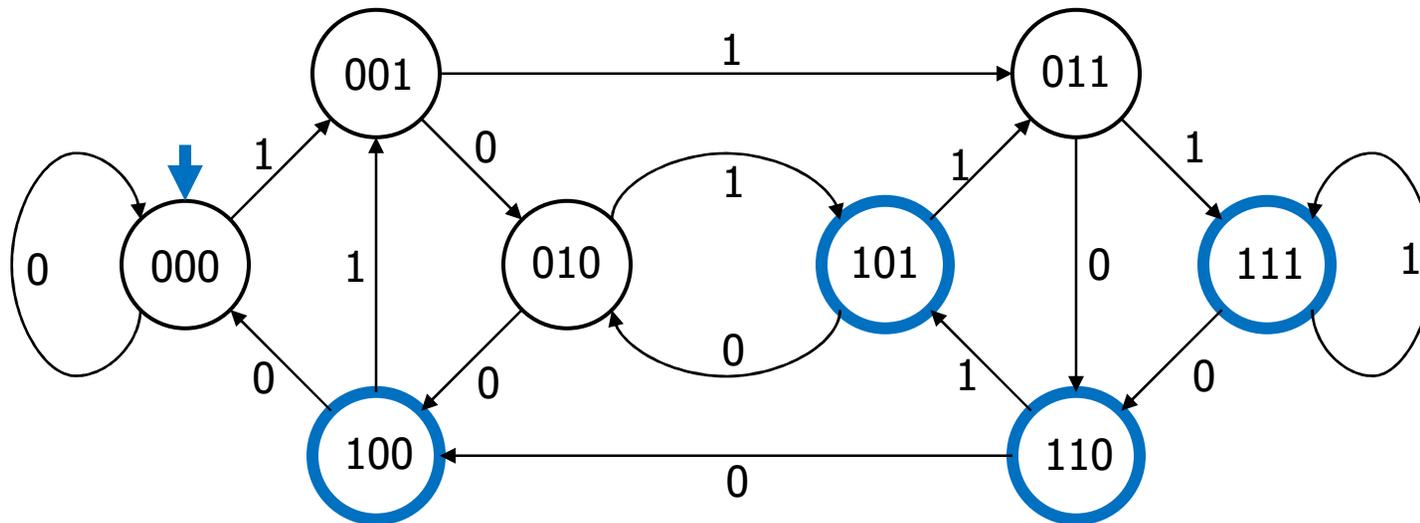
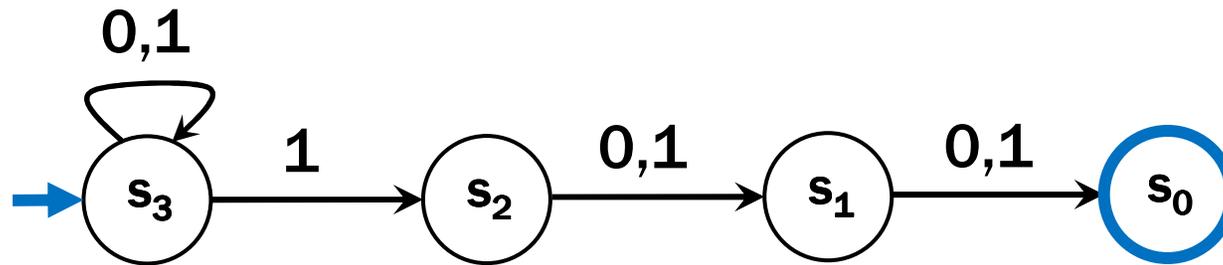
# NFA for set of binary strings with a 1 in the 3<sup>rd</sup> position from the end

---



# Compare with the smallest DFA

---



# Summary of NFAs

---

- **Generalization of DFAs**
  - drop two restrictions of DFAs
  - every DFA is an NFA
- ***Seem* to be more powerful**
  - designing is easier than with DFAs
- ***Seem* related to regular expressions**