## CSE 312 Foundations of Computing II

Lecture 20: Tail Bounds

## Review Joint PMFs and Joint Range

Definition. Let $X$ and $Y$ be discrete random variables. The Joint PMF of $X$ and $Y$ is

$$
p_{X, Y}(a, b)=P(X=a, Y=b)
$$

Definition. The joint range of $p_{X, Y}$ is

$$
\Omega_{X, Y}=\left\{(c, d): p_{X, Y}(c, d)>0\right\} \subseteq \Omega_{X} \times \Omega_{Y}
$$

Note that

$$
\sum_{(s, t) \in \Omega_{X, Y}} p_{X, Y}(s, t)=1
$$

## Review Continuous distributions on $\mathbb{R} \times \mathbb{R}$

Definition. The joint probability density function (PDF) of continuous random variables $X$ and $Y$ is a function $f_{X, Y}$ defined on $\mathbb{R} \times \mathbb{R}$ such that

- $f_{X, Y}(x, y) \geq 0$ for all $x, y \in \mathbb{R}$
- $\int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f_{X, Y}(x, y) \mathrm{d} x \mathrm{~d} y=1$
for $A \subseteq \mathbb{R} \times \mathbb{R}$ the probability that $(X, Y) \in A$ is $\iint_{A} f_{X, Y}(x, y) \mathrm{d} x \mathrm{~d} y$
The (marginal) PDFs $f_{X}$ and $f_{Y}$ are given by

$$
\begin{aligned}
& -f_{X}(x)=\int_{-\infty}^{\infty} f_{X, Y}(x, y) \mathrm{d} y \\
& -f_{Y}(y)=\int_{-\infty}^{\infty} f_{X, Y}(x, y) \mathrm{d} x
\end{aligned}
$$

## Review Law of Total Expectation

Law of Total Expectation (event version). Let $X$ be a random variable and let events $A_{1}, \ldots, A_{n}$ partition the sample space. Then,

$$
\mathbb{E}[X]=\sum_{i=1}^{n} \mathbb{E}\left[X \mid A_{i}\right] \cdot P\left(A_{i}\right)
$$

Law of Total Expectation (random variable version). Let $X$ be a random variable and $Y$ be a discrete random variable. Then,

$$
\mathbb{E}[X]=\sum_{y \in \Omega_{Y}} \mathbb{E}[X \mid Y=y] \cdot P(Y=y)
$$

## Agenda

- Covariance
- Markov's Inequality
- Chebyshev's Inequality

Covariance: How correlated are $X$ and $Y$ ?

Recall that if $X$ and $Y$ are independent, $\mathbb{E}[X Y]=\mathbb{E}[X] \cdot \mathbb{E}[Y]$

Definition: The covariance of random variables $X$ and $Y$,

$$
\operatorname{Cov}(X, Y)=\mathbb{E}[X Y]-\mathbb{E}[X] \cdot \mathbb{E}[Y]
$$

Unlike variance, covariance can be positive or negative. It has has value 0 if the random variables are independent.

Two Covariance examples:

$$
\operatorname{Cov}(X, Y)=\mathbb{E}[X Y]-\mathbb{E}[X] \cdot \mathbb{E}[Y]
$$

Suppose $X \sim \operatorname{Bernoulli}(p)$

If random variable $Y=X$ then

$$
\operatorname{Cov}(X, Y)=\mathbb{E}\left[X^{2}\right]-\mathbb{E}[X]^{2}=\operatorname{Var}(X)=p(1-p)
$$

If random variable $Z=-X$ then

$$
\begin{aligned}
\operatorname{Cov}(X, Z) & =\mathbb{E}[X Z]-\mathbb{E}[X] \cdot \mathbb{E}[Z] \\
& =\mathbb{E}\left[-X^{2}\right]-\mathbb{E}[X] \cdot \mathbb{E}[-X] \\
& =-\mathbb{E}\left[X^{2}\right]+\mathbb{E}[X]^{2}=-\operatorname{Var}(X)=-p(1-p)
\end{aligned}
$$

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- Markov’s Inequality
- Chebyshev’s Inequality


## Tail Bounds (Idea)

Bounding the probability that a random variable is far from its mean. Usually statements of the form:

$$
\begin{gathered}
P(X \geq a) \leq b \\
P(|X-\mathbb{E}[X]| \geq a) \leq b
\end{gathered}
$$

Useful tool when

- An approximation that is easy to compute is sufficient
- The process is too complex to analyze exactly


## Markov's Inequality

Theorem. Let $X$ be a random variable taking only non-negative values. Then, for any $t>0$,

$$
P(X \geq t) \leq \frac{\mathbb{E}[X]}{t} .
$$

(Alternative form) For any $k \geq 1$,

$$
P(X \geq k \cdot \mathbb{E}[X]) \leq \frac{1}{k}
$$

Incredibly simplistic - only requires that the random variable is non-negative and only needs you to know expectation. You don't need to know anything else about the distribution of $X$.

## Markov's Inequality - Proof I

Theorem. Let $X$ be a (discrete) random variable taking only non-negative values. Then, for any $t>0$,

$$
\mathbb{P}(X \geq t) \leq \frac{\mathbb{E}[X]}{t}
$$

$$
\begin{array}{rlrl}
\mathbb{E}[X] & =\sum_{x} x \cdot P(X=x) & \begin{array}{l}
\geq 0 \text { because } x \geq 0 \\
\text { whenever } P(X=x) \geq 0 \\
(X \text { takes only non-negative } \\
\text { values })
\end{array} \\
& =\sum_{x \geq t} x \cdot P(X=x)+\sum_{x<t} x \cdot P(X=x) \\
& \geq \sum_{x \geq t} x \cdot P(X=x) &
\end{array}
$$

## Markov's Inequality - Proof II

Theorem. Let $X$ be a (continuous) random variable taking only non-negative values. Then, for any $t>0$,

$$
\mathbb{P}(X \geq t) \leq \frac{\mathbb{E}[X]}{t}
$$

$$
\begin{aligned}
\mathbb{E}[X] & =\int_{0}^{\infty} x \cdot f_{X}(x) \mathrm{d} x \\
& =\int_{t}^{\infty} x \cdot f_{X}(x) \mathrm{d} x+\int_{0}^{t} x \cdot f_{X}(x) \mathrm{d} x \\
& \geq \int_{t}^{\infty} x \cdot f_{X}(x) \mathrm{d} x \\
& \geq \int_{t}^{\infty} t \cdot f_{X}(x) \mathrm{d} x=t \cdot \int_{t}^{\infty} f_{X}(x) \mathrm{d} x=t \cdot P(X \geq t)
\end{aligned}
$$

$$
\text { so } P(X \geq t) \leq \mathbb{E}[X] / t \text { as before }
$$

## Example - Geometric Random Variable

## Let $X$ be geometric RV with parameter $p$

$$
P(X=i)=(1-p)^{i-1} p \quad \mathbb{E}[X]=\frac{1}{p}
$$

" $X$ is the number of times Alice needs to flip a biased coin until she sees heads, if heads occurs with probability $p$ ?

What is the probability that $X \geq 2 \mathbb{E}[X]=2 / p$ ?
Markov's inequality: $P(X \geq 2 \mathbb{E}[X]) \leq \frac{1}{2}$
Can we do better?

## Example

Suppose that the average number of ads you will see on a website is 25 . Give an upper bound on the probability of seeing a website with 75 or more ads.

```
Poll: pollev.com/stefanotessaro617
a. 0}\leqp<0.2
b. }0.25\leqp<0.
c. }0.5\leqp<0.7
d. 0.75 \leqp
e. Unable to compute
```


## Example

Suppose that the average number of ads you will see on a website is 25 . Give an upper bound on the probability of seeing a website with 20 or more ads.

```
Poll: pollev.com/stefanotessaro617
a. 0}\leqp<0.2
b. }0.25\leqp<0.
c. }0.5\leqp<0.7
d. 0.75 \leqp
e. Unable to compute
```


## Brain Break



## Agenda

- Covariance
- Markov's Inequality
- Chebyshev’s Inequality


## Chebyshev's Inequality

Theorem. Let $X$ be a random variable. Then, for any $t>0$,

$$
P(|X-\mathbb{E}[X]| \geq t) \leq \frac{\operatorname{Var}(X)}{t^{2}}
$$

## Proof: Define $Z=X-\mathbb{E}[X]$

Definition of Variance

$$
P(|Z| \geq t)=P\left(Z^{2} \geq t^{2}\right) \leq \frac{\mathbb{E}\left[Z^{2}\right]}{t^{2}}=\frac{\mathbb{E}\left[(X-\mathbb{E}[X])^{2}\right]}{t^{2}}=\frac{\operatorname{Var}(X)}{t^{2}}
$$

$$
|Z| \geq t \text { iff } Z^{2} \geq t^{2} \quad \text { Markov's inequality }\left(Z^{2} \geq 0\right)
$$

## Example - Geometric Random Variable

Let $X$ be geometric RV with parameter $p$

$$
P(X=i)=(1-p)^{i-1} p \quad \mathbb{E}[X]=\frac{1}{p} \quad \operatorname{Var}(X)=\frac{1-p}{p^{2}}
$$

What is the probability that $X \geq 2 \mathbb{E}(X)=2 / p$ ?
Markov: $P(X \geq 2 \mathbb{E}[X]) \leq \frac{1}{2}$
Chebyshev: $P(X \geq 2 \mathbb{E}[X]) \leq P(|X-\mathbb{E}[X]| \geq \mathbb{E}[X]) \leq \frac{\operatorname{Var}(X)}{\mathbb{E}[X]^{2}}=1-p$ Better if $p>1 / 2$ ©

Suppose that the average number of ads you will see on a website is 25 and the standard deviation of the number of ads is 4 . Give an upper bound on the probability of seeing a website with 30 or more ads.

```
Poll: pollev.com/stefanotessaro617
a. 0}\leqp<0.2
b. }0.25\leqp<0.
c. }0.5\leqp<0.7
d. 0.75 \leqp
e. Unable to compute
```


## Chebyshev's Inequality - Repeated Experiments

"How many times does Alice need to flip a biased coin until she sees heads $n$ times, if heads occurs with probability $p$ ?
$X=\#$ of flips until $n$ times "heads"
$X_{i}=\#$ of flips between $(i-1)$-st and $i$-th "heads"

$$
X=\sum_{i=1}^{n} X_{i}
$$

Note: $X_{1}, \ldots, X_{n}$ are independent and geometric with parameter $p$

$$
\mathbb{E}[X]=\mathbb{E}\left[\sum_{i=1}^{n} X_{i}\right]=\sum_{i=1}^{n} \mathbb{E}\left[X_{i}\right]=\frac{n}{p} \quad \operatorname{Var}(X)=\sum_{i=1}^{n} \operatorname{Var}\left(X_{i}\right)=\frac{n(1-p)}{p^{2}}
$$

## Chebyshev's Inequality - Coin Flips

"How many times does Alice need to flip a biased coin until she sees heads $n$ times, if heads occurs with probability $p$ ?
$\mathbb{E}[X]=\mathbb{E}\left[\sum_{i=1}^{n} X_{i}\right]=\sum_{i=1}^{n} \mathbb{E}\left[X_{i}\right]=\frac{n}{p} \quad \operatorname{Var}(X)=\sum_{i=1}^{n} \operatorname{Var}\left(X_{i}\right)=\frac{n(1-p)}{p^{2}}$
What is the probability that $X \geq 2 \mathbb{E}[X]=2 n / p$ ?
Markov: $P(X \geq 2 \mathbb{E}[X]) \leq \frac{1}{2}$
Chebyshev: $P(X \geq 2 \mathbb{E}[X]) \leq P(|X-\mathbb{E}[X]| \geq \mathbb{E}[X]) \leq \frac{\operatorname{Var}(X)}{\mathbb{E}[X]^{2}}=\frac{1-p}{n}$

## Tail Bounds

Useful for approximations of complex systems. How good the approximation is depends on the actual distribution and the context you are using it in.

- Very often loose upper-bounds are okay when designing for the worst case

Generally (but not always) making more assumptions about your random variable leads to a more accurate upper-bound.

