

CSE 321: Discrete Structures

Solution to Practice Mid-term:

1. (a)

- $\forall x \forall y \forall z [D(x, y) \rightarrow D(x, yz)]$.

It's correct. Since when x is a divisor of y , x is also a divisor of yz .

- $\forall x \forall y \forall z \forall w [D(xy, zw) \rightarrow (D(x, z) \vee D(y, w))]$.

It's not correct. For example, $x = 3, y = 8, z = 4, w = 6$.

(b) Denote " $a > b$ " by $E(a, b)$, and " $a \neq b$ " by $F(a, b)$.

- "p is a prime number"

$$\forall x [D(x, p) \rightarrow (\neg F(x, p) \vee \neg F(x, 1))]$$

- "q has at least two positive divisors greater than 1"

$$\exists x [(D(x, q) \wedge E(x, 1) \wedge E(q, x))]$$

2. (a) True

(b) True

(c) True. Since q is independent to x .

(d) True

(e) True

(f) True.

(g) True

(h) False

(i) False

(j) True

(k) False. For example, $A = \{1, 2\}, B = \{1, 2, 3\}, f(x) = x$ and $g(x) = 3 - x$.

(l) True

(m) True

3. By mathematical induction.

Basis step: $n=1$. We have $1 \cdot 2^0 = 1 = (1 - 1) \cdot 2^1 + 1$.

Inductive hypothesis: Assume the argument is true for $n = k$, *i.e.*,

$$1 \cdot 2^0 + \cdots + k \cdot 2^{k-1} = (k - 1) \cdot 2^k + 1.$$

Inductive step: When $n = k + 1$, we have

$$\begin{aligned} 1 \cdot 2^0 + \cdots + k \cdot 2^{k-1} + (k + 1) \cdot 2^k &= (k - 1) \cdot 2^k + 1 + (k + 1) \cdot 2^k \\ &= 2k \cdot 2^k - 2^k + 1 + 2^k \\ &= k \cdot 2^{k+1} + 1 \end{aligned}$$

which completes the proof.

4. (i) It uses strong induction. Since in the inductive hypothesis, we assume each natural number less than n is a prime or a perfect square.

(i) Only (a) is correct. The reason why (b) is not correct is that our induction hypothesis is w.r.t all natural numbers less than n , and we only need to prove the case of n in the inductive step. The reason why (c) is not correct is that “If n is prime, then we are done” is one of two cases of our proof, and if this holds, we are done.