

Defn

$$a \bmod m = r \text{ st. } 0 \leq r < m \text{ and } a = qm + r$$

function $\mathbb{Z} : \{0 \dots m-1\}$

Defn

$$a \equiv b \pmod{m}$$

relation

$$m \mid (a-b)$$

$$a = b + km \text{ for some } k$$

Then $a \equiv b \pmod{m}$ & $c \equiv d \pmod{m}$

Then $a + c \equiv b + d \pmod{m}$

$$ac \equiv bd \pmod{m}$$

$$m \mid a-b$$

$$m \mid c-d$$

$$\therefore m \mid (a-b) + (c-d)$$

$$\therefore m \mid (a+c) - (b+d)$$

$$\therefore a+c \equiv b+d \pmod{m}$$

if $ac \equiv bc \pmod{m}$

is $a \equiv b \pmod{m}$?

$$14 \equiv 8 \pmod{6}$$

$$7 \not\equiv 4 \pmod{6}$$

but always true if $\gcd(m, c) = 1$

PT

$$m \mid ac - bc$$

$$m \mid c(a - b)$$

$$\text{since } \gcd(m, c) = 1$$

$$\therefore m \mid a - b$$

$$\therefore a \equiv b \pmod{m}$$

when/how solve

$$ax \equiv b \pmod{m}$$

if \exists we had \bar{a} st $\bar{a} \cdot a \equiv 1 \pmod{m}$

Then if $\gcd(a, m) = 1$

then $\exists \bar{a}$ st $a\bar{a} \equiv 1 \pmod{m}$

[multiplicative inverse of a]

Furthermore \bar{a} is unique
up to $\text{mod } m$]

proof (existence only)

$$\exists s, t \text{ st } s \cdot a + t \cdot m = 1$$

$$t \cdot m \equiv 0 \pmod{m}$$

$$- t \cdot m \equiv 0 \pmod{m}$$

$$s \cdot a + t \cdot m \equiv 1 \pmod{m}$$

$$\underline{- t \cdot m \equiv 0 \pmod{m}}$$

$$s \cdot a \equiv 1 \pmod{m}$$

(can find $\bar{a} = s$ via Euclid)

Chinese Remainder Theorem (CRT)

$$x \equiv a_i \pmod{m_i} \quad \text{for } i=1, 2, \dots, n$$

$$\gcd(m_i, m_j) = 1 \quad \text{for } i \neq j$$

\exists unique $0 \leq x < M = \prod m_i$
satisfying these equations.

Proof

$$\text{let } M_k = M / m_k$$

$$y_k \cdot M_k \equiv 1 \pmod{m_k}$$

$$\exists \text{ since } \gcd(M_k, m_k) = 1$$

$$x = \underbrace{\sum_i a_i M_i y_i}_{\pmod{m_j}}$$

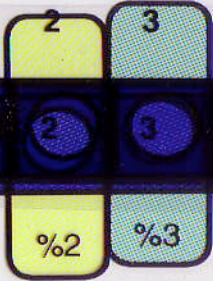
$$\equiv a_j M_j y_j \pmod{m_j}$$

$$\equiv a_j \pmod{m_j}$$

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17



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Fermat's Little Theorem

if p is prime $p \nmid a$

Then $a^{p-1} \equiv 1 \pmod{p}$

$$[\forall a \quad a^p \equiv a \pmod{p}]$$

$a^i \pmod{n}$ for $a=5, n=10$

$$5 \quad 25 \pmod{10} = 5$$

$$5^3 = 125 \pmod{10} = 5$$

$$\vdots$$
$$5^9 \equiv 5 \pmod{10}$$

fact $2^{n-1} \not\equiv 1 \pmod{n}$

for all but 22 composite numbers $n < 1000$

Fermat's Little Theorem

If p is prime and $p \nmid a$ then

$$a^{p-1} \equiv 1 \pmod{p}$$

And for all a

$$a^p \equiv a \pmod{p}$$

Proof

$$\gcd(a, p) = 1$$

$f: \mathbb{Z}_p \rightarrow \mathbb{Z}_p$ $f(i) = ai \pmod{p}$ is bijection

$$[ia \equiv ja \pmod{p} \Rightarrow i \equiv j \pmod{p}]$$

$$\prod_{i=1}^{p-1} i = \prod_{i=1}^{p-1} f(i)$$

$$(p-1)! \equiv a^{p-1} (p-1)! \pmod{p}$$

$$\gcd((p-1)!, p) = 1$$

$$1 \equiv a^{p-1} \pmod{p}$$

See also Rosen
3.7 #17

RSA - A Public Key Cryptosystem

Alice:

- ① Privately chooses two primes p, q of, say, 500 bits each, and an e rel. prime to $(p-1)(q-1)$.
- ② Privately computes $n = p \cdot q$ and d such that $ed \equiv 1 \pmod{(p-1)(q-1)}$
- ③ Publishes n and e in the phone book
(Keep p, q, d private.)

Bob (or anyone else):

Sends her an ^{encrypted} message M by looking up her n, e and sending $c = M^e \pmod n$

Alice decrypts by computing

$$c^d \pmod n = M.$$

Issues:

- do e, d always exist?
- how hard to compute?
- why $(M^e)^d \pmod n = M$?
- how secure?