

## Final Examination

Name: \_\_\_\_\_

Score: \_\_\_\_\_

This examination is open book. You may consult your class notes, homework assignments, handouts, homework solutions, and your textbook. You may cite Theorems, Lemmas, and Examples from the book; you may cite exercises from the book that were assigned as homework. You may **not** cite exercises from the book that were not on a homework assignment.

Keep your proofs succinct but complete. Write legibly. Use the back of pages if necessary.

There are 6 questions on 7 pages. The entire examination is worth 100 points. The values of individual questions and subquestions are marked.

1. [2 × 15 points] Consider the following language over the alphabet  $\Sigma = \{a, b, c\}$ .

$$L = \{a^i b^j c^k \mid i \neq j \text{ or } j \neq k \text{ or both}\}$$

- (a) [15 points] Prove that  $L$  is not regular.

- (b) [15 points] Prove that  $L$  is context-free.

2. [2 × 10 points] For any language  $L \subseteq \{0,1\}^*$ , consider the following definitions:

- $\text{EQUAL}(L) = \{x \in L \mid x \text{ has an equal number of 0's and 1's}\}.$
- $\text{DOUBLE}(L) = \{x \in L \mid x \text{ ends with } 00 \text{ or } 11\}.$

Prove or disprove the following closure properties of regular languages. If you prove your answer by construction, you need not prove your construction is correct.

(a) [10 points] For any regular language  $L$ ,  $\text{EQUAL}(L)$  is regular.

(b) [10 points] For any regular language  $L$ ,  $\text{DOUBLE}(L)$  is regular.

3. [10 points] Recall that the class of regular languages is closed under Kleene star. Prove that the converse of this closure property is false (i.e.  $L^*$  regular does not imply  $L$  regular) by giving a counter example.

4. [15 points] Consider the context-free grammar  $G = (V, \Sigma, P, S)$ , where  $V = \{S, A, B\}$ ,  $\Sigma = \{a, b\}$ , and  $P$  contains the following productions:

$$\begin{aligned} S &\rightarrow aB \mid bA \\ A &\rightarrow a \mid aS \mid bAA \\ B &\rightarrow b \mid bS \mid aBB. \end{aligned}$$

For the string  $aaabbabbba$  show (i) a rightmost derivation, (ii) a leftmost derivation, and (iii) a parse-tree.

5. [15 points] State and prove whether or not the following language is context-free:

$$L = \{a^{2^i} \mid i \geq 2\}$$

If you give an affirmative answer with a proof by construction, you need not prove that your construction is correct.

6. [10 points] In class I presented an *encoding scheme* for Turing machines with  $\Sigma = \{0, 1\}$ . This is a method by which any such Turing machine can be encoded as a string in  $\{0, 1\}^*$ . The method is such that each string in  $\{0, 1\}^*$  can be *decoded* as a unique Turing machine or as an invalid encoding. For the purposes of this problem, assume that we have a similar encoding for a *pair* of Turing machines. That is, any string  $\langle M_1, M_2 \rangle$  in  $\{0, 1\}^*$  that is in fact a valid encoding, can be uniquely decoded as a pair of Turing machines. The details of these encoding schemes are irrelevant to this problem and you should not concern yourself with them.

Consider the following language  $L = \{\langle M_1, M_2 \rangle \mid L(M_1) \cap L(M_2) \text{ is Turing-recognizable}\}$ .

State whether  $L$  is

- decidable,
- Turing-recognizable but not decidable, or
- not Turing-recognizable.

Argue briefly but convincingly for your answer.