

CSE 322  
Winter Quarter 2001  
Assignment 7  
Due Friday, February 23, 2001

All solutions should be neatly written or type set. All major steps in proofs and algorithms must be justified.

1. (10 points) Design a context-free grammar that generates the set of regular expressions (fully parenthesized) over an alphabet  $\Sigma$ . In your definition you will want to differentiate the symbol  $\varepsilon$  from the empty string.
2. (10 points) Given context-free grammars  $G_1 = (V_1, \Sigma_1, R_1, S_1)$  and  $G_2 = (V_2, \Sigma_2, R_2, S_2)$ , design context-free grammars  $G$  such that:
  - (a)  $L(G) = L(G_1)L(G_2)$  (concatenation),
  - (b)  $L(G) = L(G_1)^*$  (Kleene star),
  - (c)  $L(G) = L(G_1)^R$  (reversal).
3. (10 points) Consider the context free-grammar:

$$\begin{aligned} S &\rightarrow ASA|A|\varepsilon \\ A &\rightarrow 011|\varepsilon \end{aligned}$$

Use the method described in class to convert the grammar into Chomsky normal form. In this method do the steps in the following order: (i) add a new start symbol if  $\varepsilon$  is generated by the grammar, (ii) shorten productions whose right hand sides are longer than 2, (iii) remove  $\varepsilon$ -rules, (iv) remove unit rules, (v) make all right hand sides of length 2 into nonterminals.

4. (10 points) A context-free grammar  $G = (V, \Sigma, R, S)$  is *right-linear* if every production in  $R$  has one of the following forms (i)  $A \rightarrow wB$  where  $A, B \in V$  and  $w \in \Sigma^*$  or (ii)  $A \rightarrow w$  where  $A \in V$  and  $w \in \Sigma^*$ .
  - (a) Consider the right-linear grammar  $G = (V, \Sigma, R, S)$  where

$$\begin{aligned} V &= \{S, A, B\} \\ \Sigma &= \{a, b\} \\ R &= \{S \rightarrow aaS, S \rightarrow abA, S \rightarrow \varepsilon, A \rightarrow bA, A \rightarrow baB, \\ &\quad B \rightarrow aB, B \rightarrow bbS\} \end{aligned}$$

Design an NFA that accepts  $L(G)$ .

- (b) Prove that the language generated by a right-linear grammar is always regular.