

CSE 322: Regular Expressions and Finite Automata

◆ **Last Time:** Definition of a Regular Expression

⇒ R is a regular expression iff

R is a string over $\Sigma \cup \{ \epsilon, \emptyset, (,), \cup, * \}$ and R is:

1. Some symbol $a \in \Sigma$, or
2. ϵ , or
3. \emptyset , or
4. $(R_1 \cup R_2)$ where R_1 and R_2 are regular exps., or
5. $R_1 R_2 = R_1 \circ R_2$ where R_1 and R_2 are reg. exps., or
6. R_1^* where R_1 is a regular expression.

◆ **Precedence:** Evaluate * first, then \circ , then \cup

⇒ E.g. $0 \cup 11^* = 0 \cup (1^\circ (1^*)) = \{0\} \cup \{1, 11, 111, \dots\}$

Examples

◆ What is R for each of the following languages?

1. $L(R) = \{w \mid w \text{ contains exactly two } 0\text{'s}\}$
2. $L(R) = \{w \mid w \text{ contains at least two } 0\text{'s}\}$
3. $L(R) = \{w \mid w \text{ contains an even number of } 0\text{'s}\}$
4. $L(R) = \{w \mid w \text{ does not contain } 00\}$
5. $L(R) = \{w \mid w \text{ is a valid identifier in C}\}$ (or in Java)
6. $L(R) = \{w \mid w \text{ is a word heard on the MTV show "The Osbournes"}\}$

Are u saying our language is regular??



Regular Expressions and Finite Automata

- ◆ What is the relationship between regular expressions and DFAs/NFAs?
- ◆ Specifically:
 1. **R \rightarrow NFA**? Given a reg. exp. R, can we create an NFA N such that $L(R) = L(N)$?
 2. **NFA \rightarrow R**? Given an NFA N (or its equivalent DFA M), can we come up with a reg. exp. R such that $L(M) = L(R)$?



I think so...do you??

From Regular Expressions to NFAs

- ◆ **Problem:** Given *any* regular expression R , how do we construct an NFA N such that $L(N) = L(R)$?
- ◆ **Soln.:** Use the multi-part definition of regular expressions!!
 - ⇨ Show how to construct an NFA for each possible case in the definition: $R = a$, or $R = \epsilon$, or $R = \emptyset$, or $R = (R1 \cup R2)$, or $R = R1 \circ R2$, or $R = R1^*$.



Told ya 'twas possible!

- ◆ **Example:** Draw NFA for $10\Sigma^*01$

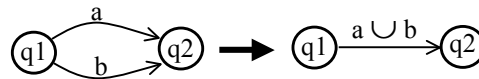
From NFAs/DFAs to Regular Expressions

- ◆ **Problem:** Given *any* NFA (or DFA) N , how do we construct a regular expression R such that $L(N) = L(R)$?
- ◆ **Solution:**
 - ⇨ **Idea:** Collapse 2 or more edges in N labeled with single symbols to a *new edge* labeled with an *equivalent regular expression*
 - ⇨ This results in a **“generalized” NFA (GNFA)**
 - ⇨ **Our goal:** Get a GNFA with 2 states (start and accept) connected by a single edge labeled with the required regular expression R

From NFAs/DFAs to Regular Expressions

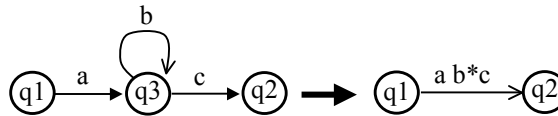
- ◆ Steps for extracting regular expressions from NFAs/DFAs:
 1. Add new start state connected to old one via an ϵ -transition
 2. Add new accept state receiving ϵ -transitions from all old ones
 3. Keep applying 2 rules until only start and accept states remain:

1. Collapse Parallel Edges:



Note: Also applies when $q1 = q2$

2. Remove “loopy” states:



Note: Also applies when $q1 = q2$