

CSE 322
Intro to Formal Models in CS
Homework #6
Due: Friday 16 Nov 07

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9 Nov 07

1. Give an informal description of a pushdown automaton recognizing the language defined by the following CFG:

$$E \rightarrow E + T \mid T$$

$$T \rightarrow T * F \mid F$$

$$F \rightarrow (E) \mid a$$

You do not need to follow either of the CFG-to-PDA constructions, nor give a formal correctness proof, but do give an informal argument/explanation for why/how it works.

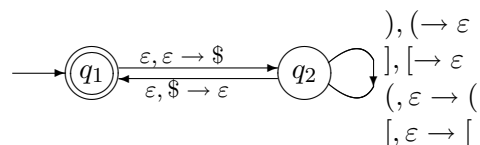
2. Let G be the following CFG:

$$S \rightarrow AA \mid B$$

$$A \rightarrow 0A \mid A0 \mid 1$$

$$B \rightarrow 00B0 \mid 1$$

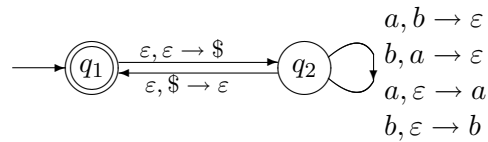
- (a) Describe $L(G)$ in English.
 - (b) Give an informal proof that your description is correct. In particular, include definitions of the strings of terminals generated by A and B .
 - (c) As in problem 1, give a PDA M recognizing the same language and informally argue correctness.
 - (d) Show that M is nondeterministic by giving a (short!) input on which it has two computations. Show the two computations.
 - (e) By analogy to the definition of an ambiguous CFG, I define a PDA to be *ambiguous* if there is some input on which it has two different accepting computations. Is M ambiguous? Why or why not? (Either show two accepting computations on some short input, or argue informally why this is impossible.)
3. Consider the following PDA M :



- (a) Describe in English the language L recognized by M . I want a *nonprocedural* description; don't say "do this then if that do something else..." (Hint: it has a very simple description.)

- (b) Prove informally that M recognizes this language. Don't forget to argue that it rejects all strings *not* in L , on all possible computations. You don't need to do a detailed induction proof or the like, but I want an argument that is thorough and convincing.

4. Consider the following PDA M :



- (a) Show that M is nondeterministic by giving a short input on which it has two computations. Show the two computations.
- (b) Describe in English the language L recognized by M . I want a *nonprocedural* description; don't say "do this then if that do something else..." (Hint: it has a very simple description.)
- (c) Prove informally that M recognizes this language. Don't forget to argue that it rejects all strings *not* in L , on all possible computations (and recall it's nondeterministic). You don't need to do a detailed induction proof or the like, but I want an argument that is thorough and convincing.
- (d) Show that M is "ambiguous" by giving a short input on which it has two accepting computations. Show the two computations.