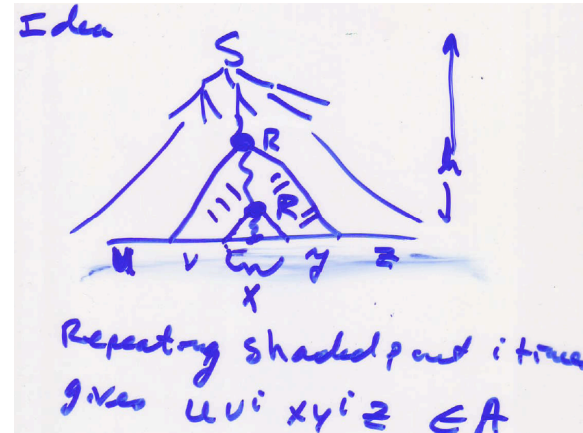


Alt proof: $\{a^n b^n c^n \mid n \geq 0\}$ not CFL:
 v & y can't have
 both a & c

Then case 1 not a
 uv^2xy^2z too few c 's
 Similarly
 case 2 not a
 ... too few a 's



$S \Rightarrow^* uRz$

$R \Rightarrow^* vRy$

$R \Rightarrow^* x$

$R \rightarrow RE$
 $E \rightarrow \epsilon$

$p = |v| + 1$

subtlety 1: true w/ fewest nodes $v^i y^i z$
 subtlety 2: prech rep. nearest lower.

$|vxy| \leq p \leq 2$

And now for something
 completely different

CFG utility beyond compilers

Actually, a Stochastic CFG

Associate probabilities with rules:

$$\begin{array}{l}
 S \rightarrow LS \quad (0.87) \quad | \quad L \quad (0.13) \\
 L \rightarrow S \quad (0.89 * p(s)) \quad | \quad dFd \quad (0.11 * p(dd)) \\
 F \rightarrow LS \quad (0.21) \quad | \quad dFd \quad (0.79 * p(dd))
 \end{array}$$

Where $p(s)$ & $p(dd)$ are the probabilities of the specific single/paired nucleotides, perhaps from empirical data or a model of sequence evolution

What SCFG Gives

“Prior” probabilities for
 frequencies of nucleotides/pairs
 fraction paired vs unpaired
 average lengths of each, etc.

Result: a probability distribution on sequences/structures

E.g., is my sequence more likely to arise under this RNA model or a simple “background” model, say where $A/C/G/T = 1/4$?

Cocke-Kasami-Younger Parser

Suppose all rules of form $A \rightarrow BC$ or $A \rightarrow a$
 (by mechanically transforming grammar, or algorithm below...)

Given $x = x_1 \dots x_n$, want $M_{ij} = \{A \mid A \rightarrow x_{i+1} \dots x_j\}$

For $j=2$ to n

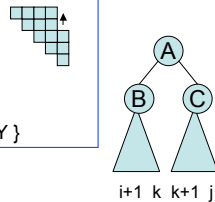
$M[j-1, j] = \{A \mid A \rightarrow x_j \text{ is a rule}\}$

for $i = j-1$ down to 1

$M[i, j] = \bigcup_{i < k < j} M[i, k] \otimes M[k, j]$

Where $X \otimes Y = \{A \mid A \rightarrow BC, B \in X, \text{ and } C \in Y\}$

Time: $O(n^3)$



“Inside” Algorithm for SCFG

Just like CKY, but instead of just recording *possibility* of A in $M[i, j]$, record its *probability*:
 For each A , do sum instead of union, over all possible k and all possible $A \rightarrow BC$ rules, of products of their respective probabilities.

Result: for each i, j, A , have $\Pr(A \Rightarrow^* x_{i+1} \dots x_j)$

The SCFG “Viterbi” algorithm

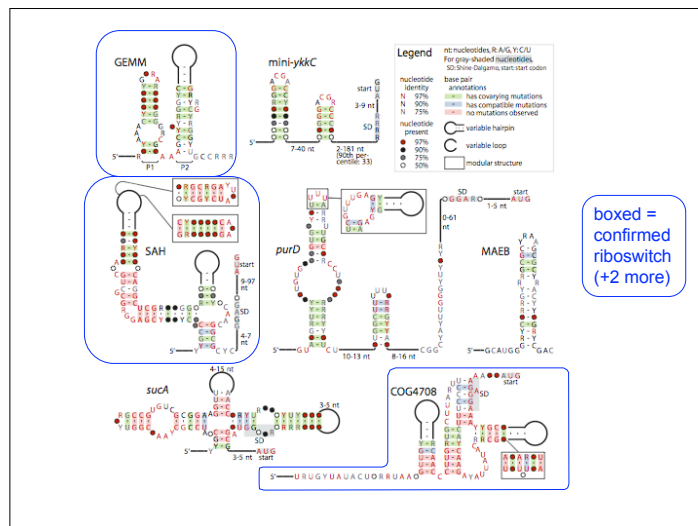
Like inside, but use max instead of sum;
Gives probability of the *single* parse tree
having max probability; (inside sums
probability over *all* legal trees)

ncRNA Discovery in Bacteria

Cmfinder--A Covariance Model Based RNA Motif Finding Algorithm, Yao, Weinberg, Ruzzo,
Bioinformatics, 2006, 22(4): 445-452,

A Computational Pipeline for High Throughput Discovery of cis-Regulatory Noncoding RNA in Prokaryotes. Yao, Barrick, Weinberg, Neph, Breaker, Tompa and Ruzzo.
PLoS Comput Biol. 3(7): e126, July 6, 2007.

Identification of 22 candidate structured RNAs in bacteria using the CMfinder comparative genomics pipeline. Weinberg, Barrick, Yao, Roth, Kim, Gore, Wang, Lee, Block, Sudarsan, Neph, Tompa, Ruzzo and Breaker.
Nucl. Acids Res., July 2007 35: 4809-4819.



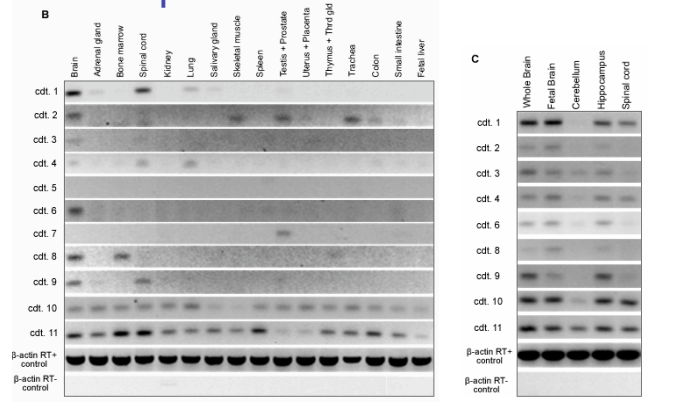
ncRNA Discovery in Vertebrates

Comparative genomics beyond sequence based alignments: RNA structures in the ENCODE regions

Torarinsson, Yao, Wiklund, Bramsen, Hansen, Kjems, Tommerup, Ruzzo and Gorodkin

Genome Research, to appear

Experimental Validation



Bottom Line

CFG technology is a key tool for RNA description, discovery and search
A very active research area. (Some call RNA the “dark matter” of the genome.)
Huge compute hog: results above represent hundreds of CPU-years, and smart algorithms can have a big impact

More?

Check out CSE 427