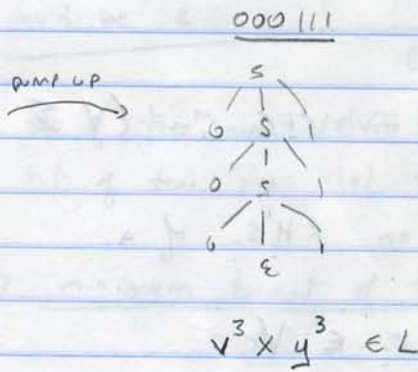
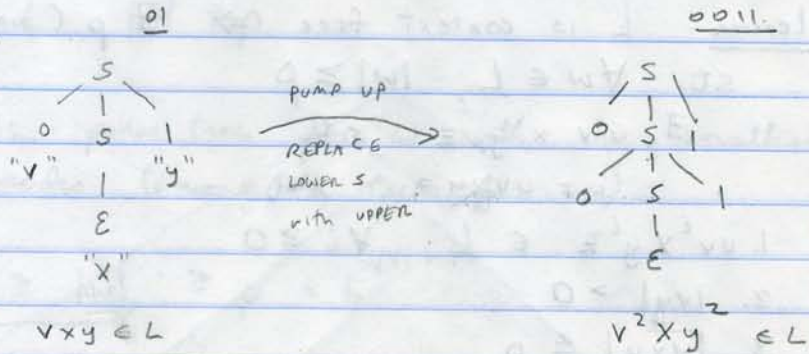


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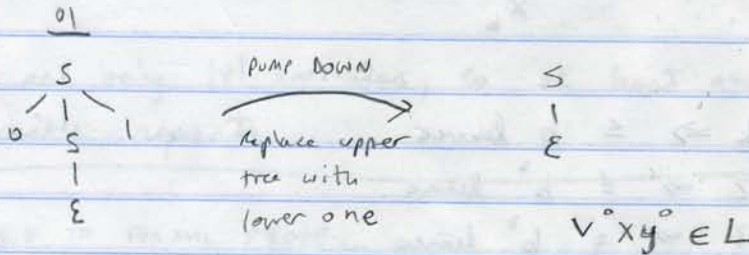
$$L = \{ 0^n 1^n \mid n \geq 0 \}$$

$$S \rightarrow 0S1 \mid \epsilon$$

Pumping Lemma for CFL's



$$v^i x y^i \in L, \forall i \geq 1 \quad \text{What about } i=0?$$



Yes,  $i=0$  also works.

Proof

Pumping Lemma:  $L$  is context free  $\Rightarrow \exists p$  (pumping length)

st.  $\forall w \in L, |w| \geq p$

$\exists u, v, x, y, z$  s.t.

$w = uvxyz$  AND

1.  $uv^i x y^i z \in L \quad \forall i \geq 0$

2.  $|v| > 0$

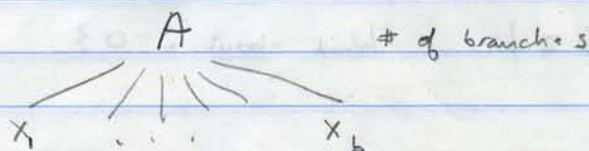
3.  $|vxy| \leq p$

Pf

$L$  is C.F.  $\Rightarrow L = L(G)$  for CFG  $G = (V, \Sigma, R, S)$

① Let  $b = \max \#$  of symbols on R.H.S. of a rule in  $R$ :

$A \rightarrow x_1 x_2 \dots x_b, \quad x_i \in V$



Ht 1  $\Rightarrow \leq b$  leaves

Ht 2  $\Rightarrow \leq b^2$  leaves

Ht 3  $\Rightarrow \leq b^3$  leaves

⋮

$h \Rightarrow \leq b^h$  leaves (can be proved by induction)

$$\text{Let } p = b^{|\nu|+1}$$

$$\text{Let } w \in L, |w| \geq p$$

Let  $T$  be a parse tree for  $w$  with the smallest # of nodes (among all trees for  $w$ )

$$\# \text{ Leaves} \geq |w| \geq p = b^{|\nu|+1}$$

because some leaves could be  $\epsilon$

#### THE REASONING

If ht of tree was  $|\nu|+1 \Rightarrow \leq b^{|\nu|+1}$  leaves  
If ht of tree was  $|\nu| \Rightarrow \leq b^{|\nu|}$  leaves

minimum height of tree is  $|\nu|+1 \Rightarrow$

there is at least one path of length  $|\nu|+1 \Rightarrow$   
this path has  $|\nu|+2$  nodes,  
the last node is a terminal  $\Rightarrow$   
the path contains at least  $|\nu|+1$  variables

There are only  $|\nu|$  variables, so at least one variable repeats.

BACK TO FORMAL PROOF...

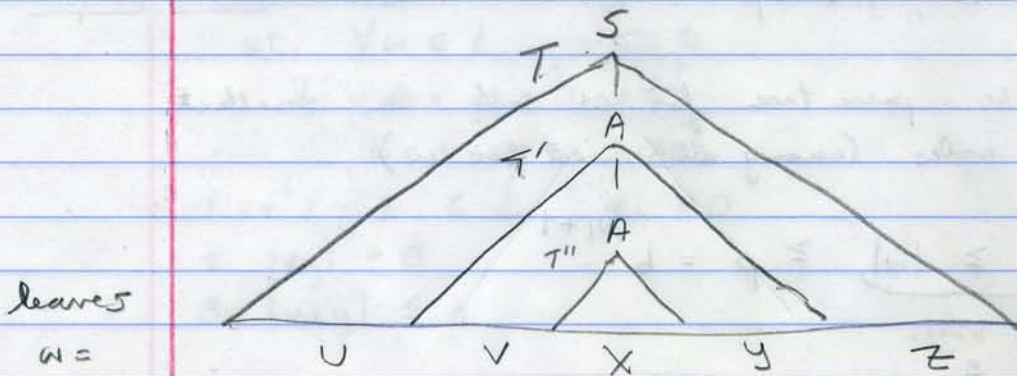
$T$ 's ht  $\geq |\nu|+1$

$\Rightarrow T$  has a path of  $|\nu|+2$  nodes, last is a terminal  
 $|\nu|+1$  nodes are variables.

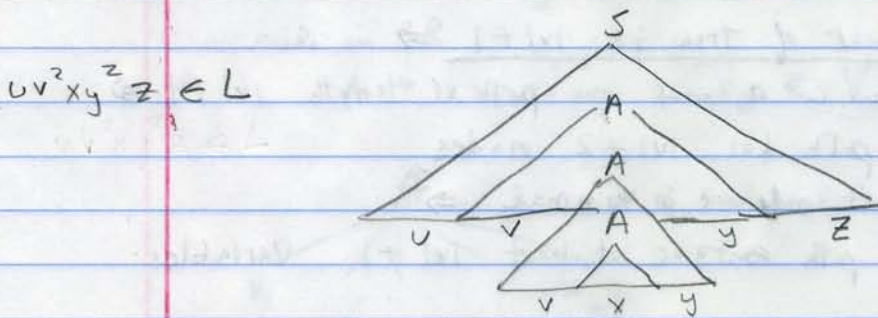
There are only  $|\nu|$  variables

By pigeonhole principle, some var  $A$  repeats

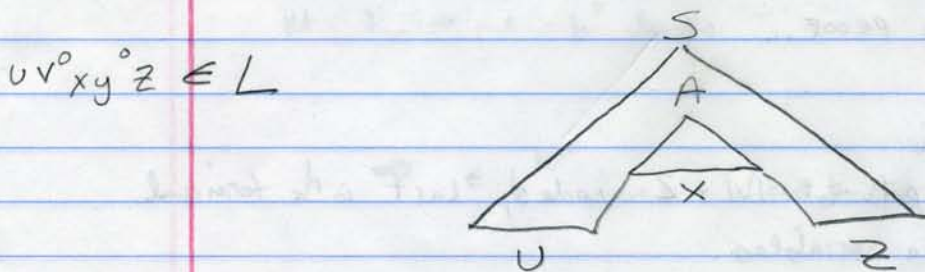
Choose  $A$  in lowermost  $|v| + 1$  vars  
on path



Pump up: Replace lower  $A$  w/ upper

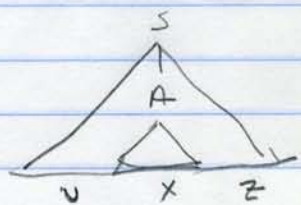


Pump down: Replace upper tree with lower



②  $|vy| > 0$

Suppose  $|vy| = 0 \Rightarrow v = E, y = E$



$w = vxz$

$\Rightarrow \exists$  a smaller tree for  $w$  than  $T$

$\Rightarrow$  contradicts  $T$  is smallest tree for  $w$

③  $|vxy| \leq p$

$A$  is in lowermost  $|v|+1$  vars on path.

$A$  is at height  $\leq |v|+1$  from bottom

$T'$  has  $\leq \begin{matrix} |v|+1 \\ b \\ p \end{matrix}$  leaves