

CSE 326: Proving asymptotic comparisons

Thursday, Jan 20, 2000

1 A little-o definition

1.1 First, Big- Ω

$f(n) \in \Omega(g(n))$ means that there is at least one $c > 0$ and some n_0 such for all $n > n_0$, $f(n) \geq c \cdot g(n)$.

1.2 little-o

$f(n) \in o(g(n))$ means that for all $c > 0$ there exists some n_0 such that for all $n > n_0$, $f(n) < c \cdot g(n)$.

1.3 BTW, another way of looking at little-o

little-o means that there's no c that will satisfy the Big- Ω condition, since *all* c don't.

1.4 Alternate little-o definition (Warning: Calculus)

For any c , and for sufficiently large n , $f(n) < c \cdot g(n)$.

In other words,

$$\begin{aligned} f(n) \in o(g(n)) \\ \iff \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0 \end{aligned}$$

Similarly,

$$\begin{aligned} f(n) \in \omega(g(n)) \\ \iff \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty \end{aligned}$$

And, interestingly,

$$\begin{aligned} f(n) \in \Theta(g(n)) \\ \iff \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = k \end{aligned}$$

for some *finite* constant $k > 0$.

(see p. 43 of the Weiss book)

2 First proof: $n \in o(n^2)$.

We use the limit of the fraction format:

$$\begin{aligned} & \lim_{n \rightarrow \infty} \frac{n}{n^2} \\ &= \lim_{n \rightarrow \infty} \frac{1}{n} \quad \text{cancel out } n \\ &= 0 \end{aligned}$$

So it's true.

3 Generalization - useful theorem

Let $g(n)$ be a monotonically increasing function (mainly $\lim_{n \rightarrow \infty} g(n) = \infty$).

Then $o(f(n) \cdot g(n)) \in o(f(n) \cdot g(n))$ if $\lim_{n \rightarrow \infty} g(n) = \infty$.

Proof: use the fraction format.

$$\begin{aligned} & \lim_{n \rightarrow \infty} \frac{g(n)}{f(n) \cdot g(n)} \\ &= \lim_{n \rightarrow \infty} \frac{1}{f(n)} \quad \text{cancel out } g(n) \\ &= 0 \quad \text{if } f(n) \text{ goes to } \infty \end{aligned}$$

4 An easy one now: $n^k \in o(n^{k+\varepsilon})$ if $k, \varepsilon > 0$

$$n^{k+\varepsilon} = n^k n^\varepsilon.$$

Clearly, $\lim_{n \rightarrow \infty} n^\varepsilon = \infty$ if $\varepsilon > 0$. So, by our above theorem, $n^k \in o(n^\varepsilon n^k)$.

5 Same idea: $\log^k n \in o(\log^{k+\varepsilon} n)$ if $k, \varepsilon > 0$

$\log^{k+\varepsilon} n = (\log^k n)(\log^\varepsilon n)$. And we have the same idea as above.

6 Another easy one: $n \in o(n \cdot \log n)$

Follows from our theorem, and the fact that $\lim_{n \rightarrow \infty} \log n = \infty$.

7 l'Hôpital's Rule

$$\begin{aligned} & \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} \\ \iff & \lim_{n \rightarrow \infty} \frac{f'(n)}{g'(n)} \end{aligned}$$

$f'(n)$ is the first derivative of $f(n)$.

8 Using l'Hôpital's Rule to show $\log n \in o(n)$

Using the fraction format,

$$\begin{aligned} & \lim_{n \rightarrow \infty} \frac{\log n}{n} \\ &= \lim_{n \rightarrow \infty} \frac{1/n}{1} \quad \text{by l'Hôpital's Rule} \\ &= \lim_{n \rightarrow \infty} \frac{1}{n} \\ &= 0 \end{aligned}$$

9 $\log^i \in o(n^j)$ for $i, j > 0$

Look at what happens with the fractional format:

$$\begin{aligned} & \lim_{n \rightarrow \infty} \frac{\log^i n}{n^j} \\ &= \lim_{n \rightarrow \infty} \frac{(\log^{i-1} n)(1/n)i}{jn^{j-1}} \quad \text{the Chain rule} \\ &= \lim_{n \rightarrow \infty} \frac{\log^{i-1} n)i}{jn^j} \\ &= \lim_{n \rightarrow \infty} \frac{\log^{i-1} n}{n^j} \quad i/j \text{ is just a constant} \end{aligned}$$

Note: taking i/j out is kind of sloppy, since we used an $=$ sign. But, it is valid given that we're only concerned about 0, some finite constant or ∞ .

Now, we can prove it inductively (we just proved the induction step). The base case is just $\log^i \in o(n^j)$ for $j > 0$ and $0 < i \leq 1$, which is pretty straightforward from what we've already done.

10 Similar idea: $n^i \in o(j^n)$ for $i, j > 0$

The induction step:

$$\begin{aligned}
 & \lim_{n \rightarrow \infty} \frac{n^i}{j^n} \\
 = & \lim_{n \rightarrow \infty} \frac{n^i}{e^{n \ln j}} \quad \text{by log rules} \\
 = & \lim_{n \rightarrow \infty} \frac{in^{i-1}}{(\ln j)e^{n \ln j}} \\
 = & \lim_{n \rightarrow \infty} \frac{in^{i-1}}{(\ln j)e^{n \ln j}} \\
 = & \lim_{n \rightarrow \infty} \frac{n^{i-1}}{e^{n \ln j}} \quad \text{eliminate constants} \\
 = & \lim_{n \rightarrow \infty} \frac{n^{i-1}}{j^e}
 \end{aligned}$$

10.1 A funky theorem: $\log f(n) \in o(\log g(n)) \implies f(n) \in o(g(n)^c)$ for any $c > 0$

Proof: By the definition of little-o,

$$\begin{aligned}
 & \log f(n) \in o(\log g(n)) \\
 \rightarrow & \forall c > 0 \exists n_0 \forall n > n_0. \log f(n) < c \cdot \log g(n)
 \end{aligned}$$

Now, we add a constant to the right side of the inequality, which preserves the little-o relation. We obtain

$$\forall k \forall c > 0 \exists n_0 \forall n > n_0. \log f(n) < c \cdot \log g(n) + k$$

So, exponentiating both sides of the inequality.

$$\begin{aligned}
 \forall k \forall c > 0 \exists n_0 \forall n > n_0. \log f(n) < c \cdot \log g(n) + k \\
 \rightarrow & \forall k \forall c > 0 \exists n_0 \forall n > n_0. f(n) < 2^k g(n)^c \\
 & \rightarrow f(n) \in o(g(n)^c)
 \end{aligned}$$

We reached the definition for little-o, since 2^k can take on all positive values for some k

A corollary is $\log f(n) \in o(\log g(n)) \implies f(n) \in o(g(n))$, by simply selecting $c = 1$.

Note that the reverse is not necessarily true. i.e. if $f(n) \in o(g(n))$ we don't necessarily know that $\log f(n) \in o(\log g(n))$. Can you think of a counterexample?

11 Interlude: $\log \log n \in o(\log n)$.

We do this using the fractional format, and by substituting $m = 2^n$:

$$\begin{aligned} & \lim_{n \rightarrow \infty} \frac{\log \log n}{\log n} \\ = & \lim_{n \rightarrow \infty} \frac{\log m}{m} \quad \text{do the substitution} \end{aligned}$$

and we know $\log m \in o(m)$ from before.

Note that it's important that $\lim_{n \rightarrow \infty} \log n = \infty$, otherwise the substitution wouldn't necessarily be valid.

12 Using funky theorem: $n^k \in o((\log n)^{\log n})$ for any k

We take log of both sides. Now, it turns out that

$$k \log n \in o((\log \log n)(\log n))$$

because the left side is $\Theta(\log n)$, while the right side has an extra $(\log \log n)$ factor on it.

And fortunately we have

$$\begin{aligned} k \log n &= \log(n^k) \\ \text{and } (\log \log n)(\log n) &= \log((\log n)^{\log n}) \end{aligned}$$

So, we can exponentiate both sides using the funky theorem, and get our answer.

13 Using funky theorem: $(\log n)^{\log n} \in o(2^k)$

Same idea as the previous one.