

5: Trees

CSE326 Spring 2002

April 8, 2002

— Dictionaries —

Remember ADTs?

Dictionary

- MakeEmpty, IsEmpty
- Find
- Insert, Delete



— Using Dictionaries —

Dictionaries are everywhere in computers

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-

Implementations of Dictionaries

Implementation	Time to		
	Find	Insert	Delete
Unordered List			
Unordered Array			
Sorted Array			

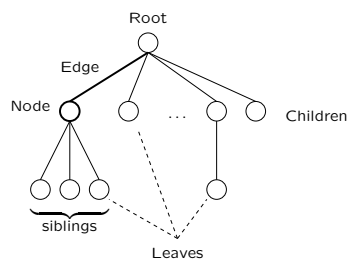
Trees

Trees in the World



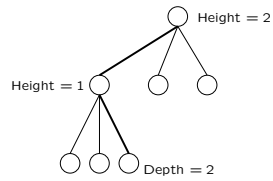
Trees in CS

Definitions



Anatomy of a Tree

Depth and Height



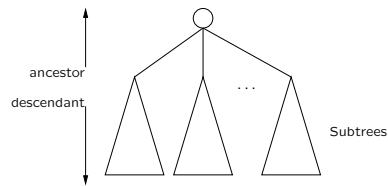
Recursive Definition

A tree is . . .

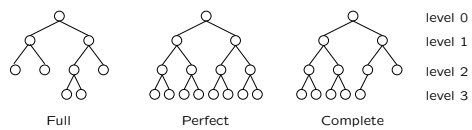
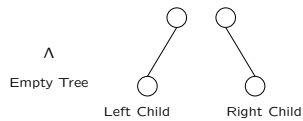
A root



A root with distinct child trees

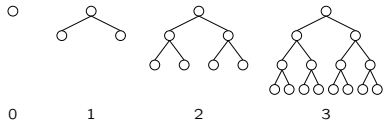


Binary Trees



Counting with Perfect Trees

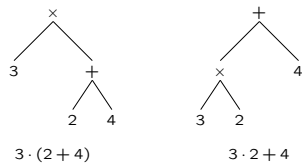
- How many nodes in level i of a perfect tree?



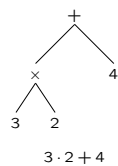
- How many nodes, total?
- What *fraction* of the tree is in the last level?

Binary Tree Traversals

- A *traversal* is an order for visiting all nodes of a tree
- Easy to think of in terms of *expression trees*



Three Kinds of Traversals



- Preorder: $+$ \times 3 2 4
- Inorder: $3 \times 2 + 4$
- Postorder: 3 2 \times 4 $+$
 - For an expression, this is *prefix* order

Code for Traversals

```
void TreeNode::PreOrder() {
    Visit ();
    if ( LeftChild () ) LeftChild()->PreOrder();
    if ( RightChild () ) RightChild()->PreOrder();
}

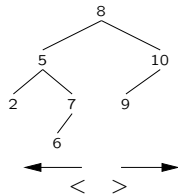
void TreeNode::InOrder() {

}

void TreeNode::PostOrder() {

}
```

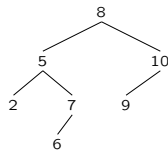
Binary Search Trees



```
Node *Node::Find(Key k)
{
    if (key == k) return this;
    if (k < key)
        && LChild()
        return LChild->Find(k);
    if (k > key)
        && RChild()
        return RChild->Find(k);
    return NULL;
}
```

```
void Node::Insert (Key k)
{
    if (k <= key) {
        if (LChild())
            LChild()->Insert(k);
        else
            SetLChild(new Node(k));
    }
    if (k > key) {
        if (RChild())
            RChild()->Insert(k);
        else
            SetRChild(new Node(k));
    }
}
```

Traversals on BSTs

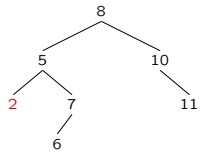


Preorder

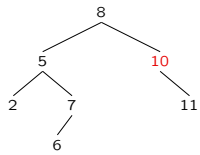
Inorder

Postorder

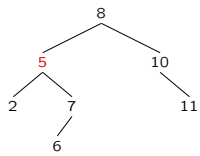
BST Deletion I



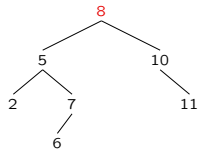
BST Deletion II



BST Deletion III



BST Deletion IV



BST Deletion

- Locate node to delete
- If no children, delete
- If one child, replace with child
- Otherwise,
 - Locate *successor* (or *predecessor*)
 - Replace with successor (or predecessor)

BST Analysis

- Perfect Tree

Complete? Full?

- *Unbalanced* Tree

— Making an Unbalanced Tree —