| CSE 326: Data Structures |
| :---: |
| Topic \#17: |
| Let's get connected... minimally! |
| Ashish Sabharwal |
| Autumn, 2003 |

## Today's Outline

- Discuss Quiz \#5
- Finish Shortest Path Problems
- Minimum Spanning Trees


## Before we move on...

- Dijkstra's algorithm, as we saw, gives the minimum distance between $s$ and $t$.
- Can we modify it to output the shortest path between $s$ and $t$ ?


An Application:
Moving Around Washington


What's the fastest way from Seattle to Pullman?
Answer:

A Different Application:
Communication in Washington


What's the cheapest inter-city network?

Is This Problem Really Different?

- Is knowing Dijkstra's algorithm enough to solve the latter application?

Yes? Then how?
No? Then why?


## Prim's Algorithm for MST

## A node-based greedy algorithm

 Builds MST by greedily adding nodes1. Select a node to be the "root"

- mark it as known
- Update cost of all its neighbors

2. While there are unknown nodes left in the graph
a. Select an unknown node $b$ with the smallest cost from some known node $a$
b. Mark $b$ as known
c. Add $(a, b)$ to MST
d. Update cost of all nodes adjacent to $b$

## Prim's Algorithm: Complexity

- Depends on what?
- How long does each step take?


## Runtime:

Prim's Algorithm: Example


Prim's Algorithm: Correctness

- A proof very similar to that of Dijkstra's algorithm works!
(left as exercise)


## Kruskal's Algorithm for MST

An edge-based greedy algorithm Builds MST by greedily adding edges

1. Initialize with

- empty MST
- all vertices marked unconnected
- all edges unmarked

2. While there are still unmarked edges
a. Pick the lowest cost edge ( $u, v$ ) and mark it
b. If $\mathbf{u}$ and $\mathbf{v}$ are not already connected, add $(\mathbf{u}, \mathbf{v})$ to the MST and mark $\mathbf{u}$ and $\mathbf{v}$ as connected to each other

## Kruskal's Algorithm: Complexity

- Depends, of course, on the data structures/ADT used. What should we use?
- How long does each step take?


## Kruskal's Algorithm: Correctness

It clearly generates a spanning tree. Call it $\mathrm{T}_{\mathrm{K}}$.
Suppose $\mathrm{T}_{\mathrm{K}}$ is not minimum:
Pick another spanning tree $\mathrm{T}_{\text {min }}$ with lower cost than $\mathrm{T}_{\mathrm{K}}$ Pick the smallest edge $e_{1}=(u, v)$ in $\mathrm{T}_{\mathrm{K}}$ that is not in $\mathrm{T}_{\text {min }}$ $\mathrm{T}_{\text {min }}$ already has a path $p$ in $\mathrm{T}_{\text {min }}$ from $u$ to $v$ $\Rightarrow$ Adding $e_{1}$ to $\mathrm{T}_{\text {min }}$ will create a cycle in $\mathrm{T}_{\text {min }}$
Pick an edge $e_{2}$ in $p$ that Kruskal's algorithm considered after adding $e_{1}$ (must exist: u and v unconnected when $\mathrm{e}_{1}$ considered) $\Rightarrow \operatorname{cost}\left(e_{2}\right) \geq \operatorname{cost}\left(e_{1}\right)$
$\Rightarrow$ can replace $e_{2}$ with $e_{1}$ in $\mathrm{T}_{\text {min }}$ without increasing cost!
Keep doing this until $\mathrm{T}_{\text {min }}$ is identical to $\mathrm{T}_{\mathrm{K}}$
$\Rightarrow \mathrm{T}_{\mathrm{K}}$ must also be minimal - contradiction!

## Kruskal's Algorithm: Example




1. Starting at node A, find the MST using Prim's method. (continue on next slide)

Play at Home with Kruskal

2. Now find the MST using Kruskal's method.
3. Under what conditions will these methods give the same result?
4. What data structures should be used for Kruskal's? Running time?


