## CSE 326: Data Structures <br> Dynamic Programming - <br> Floyd/Warhsall Algorithm

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Lectures 26

## Analysis

- Total running time for Dijkstra's:
$\mathrm{O}\left(|\mathrm{V}|^{2}+|\mathrm{E}|\right) \quad$ (linear scan)
$\mathrm{O}(|\mathrm{V}| \log |\mathrm{V}|+|\mathrm{E}| \log |\mathrm{V}|) \quad$ (heaps)

What if we want to find the shortest path from each point to ALL other points?

## Single-Source Shortest Path

- Given a graph $G=(V, E)$ and a single distinguished vertex s, find the shortest weighted path from s to every other vertex in $G$.


## All-Pairs Shortest Path:

- Find the shortest paths between all pairs of vertices in the graph.
- How?


## Dynamic Programming

Algorithmic technique that systematically records the answers to sub-problems in a table and reuses those recorded results (rather than recomputing them).

Simple Example: Calculating the Nth Fibonacci number.

$$
\operatorname{Fib}(\mathrm{N})=\operatorname{Fib}(\mathrm{N}-1)+\operatorname{Fib}(\mathrm{N}-2)
$$

## Floyd-Warshall

```
for (int k = 1; k =< V; k++)
    for (int i = 1; i =< V; i++)
    for (int j = 1; j =< v; j++)
        if ((M[i][k]+ M[k][j] ) < M[i][j] )
        M[i][j] = M[i][k]+ M[k][j]
```

Invariant: After the kth iteration, the matrix includes the shortest paths for all pairs of vertices (i,j) containing only vertices $1 . . \mathrm{k}$ as intermediate vertices

Floyd-Warshall for All-pairs shortest path


|  | $a$ | $b$ | $c$ | $d$ | $e$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $a$ | 0 | 2 | 0 | -4 | 0 |
| $b$ | - | 0 | -2 | 1 | -1 |
| c | - | - | 0 | - | 1 |
| d | - | - | - | 0 | 4 |
| e | - | - | - | - | 0 |

Final Matrix Contents

Initial state of the matrix:

|  | a | b | c | d | e |
| :--- | :--- | :--- | :--- | :--- | :--- |
| a | 0 | 2 | - | -4 | - |
| b | - | 0 | -2 | 1 | 3 |
| c | - | - | 0 | - | 1 |
| d | - | - | - | 0 | 4 |
| e | - | - | - | - | 0 |

$M[i][j]=\min (M[i][j], M[i][k]+M[k][j])$
6

