
CSE 331

Software Design & Implementation

Autumn 2022

Section 6 – HW6 and Midterm Review

Administrivia

- Done with HW5!
- HW6 (ADT implementation) due next week (**Thurs.** 11/10)
 - We will get feedback out for you early next week
- Midterm tomorrow during lecture!
- Any questions?

Agenda

- Walk-through of the test-script driver (to run `.test` files)
- Managing an expensive `checkRep`
- Midterm review

Refresher: Format of script tests

Each script test is expressed as text-based script `foo.test`

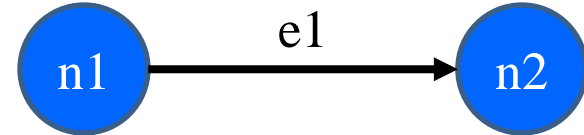
- One command per line, of the form: **Command** *arg₁* *arg₂* ...
- Script's output compared against `foo.expected`
- Precise details specified in the homework
- Match format **exactly**, including whitespace!

Command (in <code>foo.test</code>)	Output (in <code>foo.expected</code>)
<code>CreateGraph</code> <i>name</i>	<code>created graph</code> <i>name</i>
<code>AddNode</code> <i>graph label</i>	<code>added node</code> <i>label</i> <code>to graph</code>
<code>AddEdge</code> <i>graph parent child label</i>	<code>added edge</code> <i>label</i> <code>from parent to child in graph</code>
<code>ListNodes</code> <i>graph</i>	<code>graph contains:</code> <i>label</i> _{<i>node</i>} ...
<code>ListChildren</code> <i>graph parent</i>	<code>the children of parent in graph are:</code> <i>child</i> (<i>label</i> _{<i>edge</i>}) ...
<code>#</code> <i>This is comment text ...</i>	<code>#</code> <i>This is comment text ...</i>

Refresher: example.test

```
# Create a graph  
CreateGraph graph1
```

```
# Add a pair of nodes  
AddNode graph1 n1  
AddNode graph1 n2
```



```
# Add an edge  
AddEdge graph1 n1 n2 e1
```

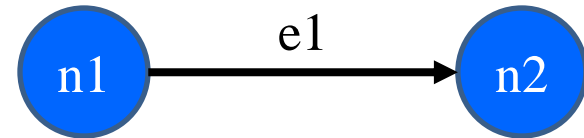
```
# Print all nodes in the graph  
ListNodes graph1
```

```
# Print all child nodes of n1 with outgoing edge  
ListChildren graph1 n1
```

Refresher: example.expected

```
# Create a graph
created graph graph1
```

```
# Add a pair of nodes
added node n1 to graph1
added node n2 to graph1
```



```
# Add an edge
added edge e1 from n1 to n2 in graph1
```

```
# Print all nodes in the graph
graph1 contains: n1 n2
```

```
# Print all child nodes of n1 with outgoing edge
the children of n1 in graph1 are: n2(e1)
```

How the script tests work

- In HW5, you wrote script tests in the form of `.test` files
 - As well as an `.expected` file for each test's expected outcome
- The JUnit class `ScriptFileTests` runs all these tests
 - Looks for all the `.test` files in the `src/test/resources/testScripts` folder
 - Compares test output against corresponding `.expected` file
- `ScriptFileTests` needs a bridge to your graph implementation
 - That's exactly what the `GraphTestDriver` class is for

Graph Test Driver

- **GraphTestDriver** knows how to read these test scripts
- **GraphTestDriver** calls a method to “do” each verb
 - **CreateGraph**, **AddNode**, **AddEdge** ...
 - One method stub per script command for you to fill with calls to your graph code
- Note: Completed test driver should sort lists before printing for **ListNodes** and **ListChildren**
 - Just to ensure predictable, deterministic output
 - Your graph implementation itself should not worry about sorting

Graph Test Driver Output

- The Graph Test Driver is a client of our graph...
 - ...but not the only client.
 - Your graph should not be designed to be exclusively used for the test driver.
- ListChildren in the test driver should print out: “**the children of parent in graph are: $child(label_{edge}) \dots$ ”**”
- This does **not** mean that you should have a method on your graph called ListChildren that returns this String
 - Because that isn’t useful for other clients

Sorting with the driver

- **Use the test driver appropriately!**
 - From before: “Completed test driver should sort lists before printing.”
- Script test output for hw5 needs to be sorted so we can mechanically check it.
- This means sorted output for tests does ***NOT*** mean sorted internal storage in graph.
 - If sorting behavior is needed, Graph ADT clients (including the test driver) can sort those labels.

In other words...

The Graph ADT in general should **NOT** assume that node or edge labels are sorted or even comparable(!).
(of course they can be tested for equality with equals())

Demo

Here's a quick tour of the `GraphTestDriver`!

Expensive checkReps

- A complicated rep. invariant can be expensive to check
 - Especially iterating over internal collection(s)
 - For example, examining every edge in a graph
- A slow `checkRep` could cause our grading scripts to time-out
 - Can be really useful during testing/debugging, but
 - Need to disable the really slow checks before submitting
- We have a tension between two goals:
 - Thorough, possibly slow checking for development
 - Essential, necessarily fast checking for production/grading
- What to do?

Use a debug flag to tune `checkRep`

- Repeatedly (un)commenting sections of code is a poor solution
- Instead, use a class-level constant as a toggle
 - Ex.: `private static final boolean DEBUG = ...;`
 - `false` for only the fast, essential checks
 - `true` for all the slow, thorough checks
 - Real-world code often has several such “debug levels”

```
private void checkRep() {
    assert fast_checks();
    if (DEBUG)
        assert slow_checks();
}
```

Midterm Review

intToString()

Fill in the implementation of a method that converts a **positive integer** to its **string representation in decimal** (invariant given on next slide).

```
{ { P: x > 0 } }
```

```
String intToString(int x)
```

Useful facts to recall:

1. Convert **char** **ch** that is one of `'0'`, `'1'`, ..., `'9'` to a corresponding **int** by doing `ch - '0'`
2. Convert **int** **x** that is one of `0`, `1`, ..., `9` to a corresponding **char** by doing `(char) (x + '0')`

intToString()

```
{ { P: x > 0 } }  
String intToString(int x) {  
    StringBuilder buf =  
    int k = , y = ;  
    { { Inv: P and buf stores the lowest k digits of x  
        in reverse order and y = x / 10^k } }  
    while (y != 0) {  
  
        k = k + 1;  
    }  
  
    return buf.reverse().toString();  
}
```

How do we get the invariant to hold initially?

intToString()

```
{ { P: x > 0 } }  
String intToString(int x) {  
    StringBuilder buf = new StringBuilder();  
    int k = 0, y = x;  
    { { Inv: P and buf stores the lowest k digits of x  
        in reverse order and y = x / 10^k } }  
    while (y != 0) {  
  
        k = k + 1;  
    }  
  
    return buf.reverse().toString();  
}
```

How do we fill out the loop body?

intToString()

```
{ { P: x > 0 } }  
String intToString(int x) {  
    StringBuilder buf = new StringBuilder();  
    int k = 0, y = x;  
    { { Inv: P and buf stores the lowest k digits of x  
        in reverse order and y = x / 10^k } }  
    while (y != 0) {  
  
        k = k + 1;  
    }  
  
    return buf.reverse().toString();  
}
```

Inv changes k to $k+1$, so

- y becomes $x / 10^{k+1}$
- $y_{\text{post}} = x / 10^{k+1} = y_{\text{pre}} / 10$

intToString()

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{ { P: x > 0 } }  
String intToString(int x) {  
    StringBuilder buf = new StringBuilder();  
    int k = 0, y = x;  
    { { Inv: P and buf stores the lowest k digits of x  
        in reverse order and y = x / 10^k } }  
    while (y != 0) {  
  
        y = y / 10;  
        k = k + 1;  
    }  
  
    return buf.reverse().toString();  
}
```

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intToString()

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    while (y != 0) {  
  
        y = y / 10;  
        k = k + 1;  
    }  
  
    return buf.reverse().toString();  
}
```

- Inv changes k to $k+1$, so
- buf stores lowest $k+1$ digits

intToString()

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{ { P: x > 0 } }  
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    StringBuilder buf = new StringBuilder();  
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    while (y != 0) {  
  
        y = y / 10;  
        k = k + 1;  
    }  
  
    return buf.reverse().toString();  
}
```

Inv changes k to $k+1$, so

- buf stores lowest $k+1$ digits

$(k+1)$ -st lowest digit goes at end
since buf stores them reversed

intToString()

```
{ { P: x > 0 } }  
String intToString(int x) {  
    StringBuilder buf = new StringBuilder();  
    int k = 0, y = x;  
    { { Inv: P and buf stores the lowest k digits of x  
        in reverse order and y = x / 10^k } }  
    while (y != 0) {  
        char ch = ?  
        buf.append(ch);  
        y = y / 10;  
        k = k + 1;  
    }  
  
    return buf.reverse().toString();  
}
```

How can we get the (k+1)-st lowest digit of x? And make it a char?

intToString()

```
{ { P: x > 0 } }
String intToString(int x) {
    StringBuilder buf = new StringBuilder();
    int k = 0, y = x;
    { { Inv: P and buf stores the lowest k digits of x
        in reverse order and y = x / 10^k } }
    while (y != 0) {
        char ch = (char) (y % 10 + '0');
        buf.append(ch);
        y = y / 10;
        k = k + 1;
    }

    return buf.reverse().toString();
}
```


intToString() solution

```
{ { P: x > 0 } }
String intToString(int x) {
    StringBuilder buf = new StringBuilder();
    int k = 0, y = x;
    { { Inv: P and buf stores the lowest k digits of x
        in reverse order and y = x / 10^k } }
    while (y != 0) {
        char ch = (char) (y % 10 + '0');
        buf.append(ch);
        y = y / 10;
        k = k + 1;
    }
    { { buf stores the digits of x in reverse order } }
    return buf.reverse().toString();
}
```

Why does this hold?




intToString() solution

```
{ { P: x > 0 } }
String intToString(int x) {
    StringBuilder buf = new StringBuilder();
    int k = 0, y = x;
    { { Inv: P and buf stores the lowest k digits of x
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    while (y != 0) {
        char ch = (char) (y % 10 + '0');
        buf.append(ch);
        y = y / 10;
        k = k + 1;
    }
    { { buf stores the digits of x in reverse order } }
    return buf.reverse().toString();
}
```

Why does this hold?
 $y = 0 \Rightarrow x < 10^k$ so x only
has k digits

Specifications

Suppose we have a **BankAccount** class with instance variable `balance`. Consider the following specifications (ignore `@param`):




- A. `@effects` decreases `balance` by `amount` 
- B. `@requires` `amount >= 0` and `amount <= balance`
`@effects` decreases `balance` by `amount` 
- C. `@throws` `InsufficientFundsException` if `balance < amount`
`@effects` decreases `balance` by `amount` 

Which specifications does this implementation meet?

```
void withdraw(int amount) {  
    balance -= amount;  
}
```

Specifications

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- A. `@effects` decreases `balance` by `amount` 
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`@effects` decreases `balance` by `amount` 
- C. `@throws` `InsufficientFundsException` if `balance < amount`
`@effects` decreases `balance` by `amount` 

Which specifications does this implementation meet?

```
void withdraw(int amount) {  
    if (balance >= amount) balance-=amount;  
}
```

Specifications

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- A. `@effects` decreases `balance` by `amount` 
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`@effects` decreases `balance` by `amount` 
- C. `@throws` `InsufficientFundsException` if `balance < amount`
`@effects` decreases `balance` by `amount` 

Which specifications does this implementation meet?

```
void withdraw(int amount) {  
    if (amount < 0) throw new IllegalArgumentException();  
    balance -= amount;  
}
```

Specifications

Suppose we have a **BankAccount** class with instance variable `balance`. Consider the following specifications (ignore `@param`):

- A. `@effects` decreases `balance` by `amount` 
- B. `@requires` `amount >= 0` and `amount <= balance`
`@effects` decreases `balance` by `amount` 
- C. `@throws` `InsufficientFundsException` if `balance < amount`
`@effects` decreases `balance` by `amount` 

Which specifications does this implementation meet?

```
void withdraw(int amount) throws InsufficientFundsException {  
    if (balance < amount) throw new InsufficientFundsException();  
    balance -= amount;  
}
```

Testing

Consider the `BankAccount` class again. What are some good test cases?

```
public class BankAccount {
    /** @return current balance of account */
    public void balance() { ... }

    /**
     * @param amount to withdraw
     * @requires amount >= 0
     * @throws InsufficientFundsException
     *         if balance < amount
     * @effects decreases balance by amount
     */
    public void withdraw(int amount) { ... }
}
```

Specification test heuristic:

- amount \leq balance
- amount $>$ balance

Boundary test heuristic:

- amount = balance
- amount $>$ balance

Others?

Should we test amount $<$ 0?

More Reasoning

Let's check that this method is correct.

```
/** Return the value of this IntPoly at point x */
public int valueAt(int x) {
    int val = a[0];
    int xk = 1;
    int k = 0;
    int n = a.length - 1;
    {{ inv: xk = x^k && val = a[0] + a[1]*x + ... + a[k]*x^k }}
    while (k != n) {
        {{ ? }}
        xk = xk * x;
        {{ ? }}
        val = val + a[k+1]*xk;
        {{ ? }}

        k = k + 1;
        {{ ? }}
    }
    {{ val = a[0] + a[1]*x + ... + a[n]*x^n }}
    return val;
}
```


More Reasoning

Let's check that this method is correct.

```
/** Return the value of this IntPoly at point x */
public int valueAt(int x) {
    int val = a[0];
    int xk = 1;
    int k = 0;
    int n = a.length - 1;
    {{ inv: xk = x^k && val = a[0] + a[1]*x + ... + a[k]*x^k }}
    while (k != n) {
        {{ ? }}
        xk = xk * x;
        {{ ? }}
        val = val + a[k+1]*xk;
        {{ ? }}

        k = k + 1;
        {{ ? }}
    }
    {{ val = a[0] + a[1]*x + ... + a[n]*x^n }}
    return val;
}
```

Does the invariant hold before the loop?

More Reasoning

Let's check that this method is correct.

```
/** Return the value of this IntPoly at point x */
public int valueAt(int x) {
    int val = a[0];
    int xk = 1;
    int k = 0;
    int n = a.length - 1;
    {{ inv: xk = x^k && val = a[0] + a[1]*x + ... + a[k]*x^k }}
    while (k != n) {
        {{ ? }}
        xk = xk * x;
        {{ ? }}
        val = val + a[k+1]*xk;
        {{ ? }}

        k = k + 1;
        {{ ? }}
    }
    {{ val = a[0] + a[1]*x + ... + a[n]*x^n }}
    return val;
}
```

More Reasoning

Let's check that this method is correct.

```
/** Return the value of this IntPoly at point x */
public int valueAt(int x) {
    int val = a[0];
    int xk = 1;
    int k = 0;
    int n = a.length - 1;
    {{ inv: xk = x^k && val = a[0] + a[1]*x + ... + a[k]*x^k }}
    while (k != n) {
        {{ inv && k != n }}
        xk = xk * x;
        {{ ? }}
        val = val + a[k+1]*xk;
        {{ ? }}

        k = k + 1;
        {{ ? }}
    }
    {{ val = a[0] + a[1]*x + ... + a[n]*x^n }}
    return val;
}
```

More Reasoning

Let's check that this method is correct.

```
/** Return the value of this IntPoly at point x */
public int valueAt(int x) {
    int val = a[0];
    int xk = 1;
    int k = 0;
    int n = a.length - 1;
    {{ inv: xk = x^k && val = a[0] + a[1]*x + ... + a[k]*x^k }}
    while (k != n) {
        {{ inv && k != n }}
        xk = xk * x;
        {{ xk = x^(k+1) && val = a[0] + a[1]*x + ... + a[k]*x^k && k != n }}
        val = val + a[k+1]*xk;
        {{ ? }}

        k = k + 1;
        {{ ? }}
    }
    {{ val = a[0] + a[1]*x + ... + a[n]*x^n }}
    return val;
}
```

More Reasoning

Let's check that this method is correct.

```
/** Return the value of this IntPoly at point x */
public int valueAt(int x) {
    int val = a[0];
    int xk = 1;
    int k = 0;
    int n = a.length - 1;
    {{ inv: xk = x^k && val = a[0] + a[1]*x + ... + a[k]*x^k }}
    while (k != n) {
        {{ inv && k != n }}
        xk = xk * x;
        {{ xk = x^(k+1) && val = a[0] + a[1]*x + ... + a[k]*x^k && k != n }}
        val = val + a[k+1]*xk;
        {{ xk = x^(k+1) && val = a[0] + a[1]*x + ... + a[k+1]*x^(k+1)
            && k != n }}
        k = k + 1;
        {{ ? }}
    }
    {{ val = a[0] + a[1]*x + ... + a[n]*x^n }}
    return val;
}
```

More Reasoning

Let's check that this method is correct.

```
/** Return the value of this IntPoly at point x */
public int valueAt(int x) {
    int val = a[0];
    int xk = 1;
    int k = 0;
    int n = a.length - 1;
    {{ inv: xk = x^k && val = a[0] + a[1]*x + ... + a[k]*x^k }}
    while (k != n) {
        {{ inv && k != n }}
        xk = xk * x;
        {{ xk = x^(k+1) && val = a[0] + a[1]*x + ... + a[k]*x^k && k != n }}
        val = val + a[k+1]*xk;
        {{ xk = x^(k+1) && val = a[0] + a[1]*x + ... + a[k+1]*x^(k+1)
            && k != n }}

        k = k + 1;
        {{ inv && k-1 != n }} -> {{ inv }}
    }
    {{ val = a[0] + a[1]*x + ... + a[n]*x^n }}
    return val;
}
```

Do we reach the postcondition?

Good luck on the midterm!