CSE 332 Autumn 2023
Lecture 8: Dictionaries, BSTs

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Warm Up: Give pseudocode to calculate the height of a Tree

```
treeHeight(root){
        height = 0;
        ???
        return height;
}
```

Tree Height

```
treeHeight(root){
        height = 0;
        if (root.left != Null){
                height = max(height, treeHeight(root.left));
        }
        if (root.right != Null){
                height = max(height, treeHeight(root.right));
        }
        return height;
}
```

Dictionary (Map) ADT

- Contents:
  - Sets of key+value pairs
  - Keys must be comparable
- Operations:
  - insert(key, value)
    - Adds the (key,value) pair into the dictionary
    - If the key already has a value, overwrite the old value
      - Consequence: Keys cannot be repeated
  - find(key)
    - Returns the value associated with the given key
  - delete(key)
    - Remove the key (and its associated value)

Naïve attempts

Data Structure	Time to insert	Time to find	Time to delete
Unsorted Array	$\Theta(n)$	$\Theta(n)$	$\Theta(n)$
Unsorted Linked List	$\Theta(n)$	$\Theta(n)$	$\Theta(n)$
Sorted Array	$\Theta(n)$	$\Theta(\log n)$	$\Theta(n)$
Sorted Linked List	$\Theta(n)$	$\Theta(n)$	$\Theta(n)$

Less Naïve attempts

- Binary Search Trees (today)
- Tries (Project)
- AVL Trees (next week)
- B-Trees (next week)
- HashTables (week after)
- Red-Black Trees (not included in this course)
- Splay Trees (not included in this course)

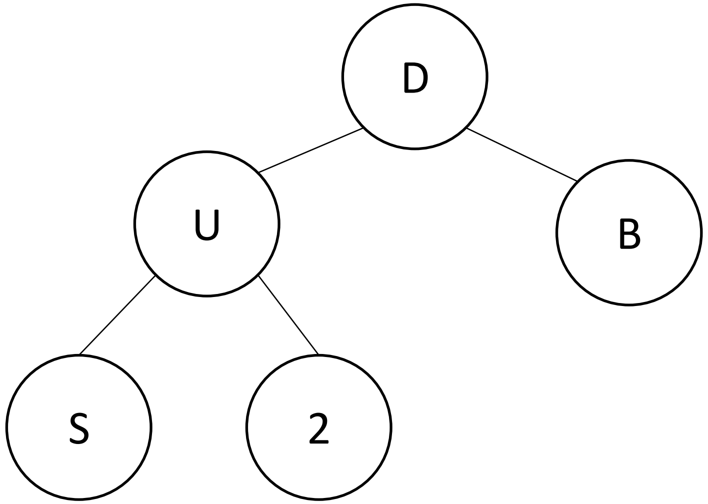
Naïve attempts

Data Structure	Time to insert	Time to find	Time to delete
Unsorted Array	$\Theta(n)$	$\Theta(n)$	$\Theta(n)$
Unsorted Linked List	$\Theta(n)$	$\Theta(n)$	$\Theta(n)$
Sorted Array	$\Theta(n)$	$\Theta(\log n)$	$\Theta(n)$
Sorted Linked List	$\Theta(n)$	$\Theta(n)$	$\Theta(n)$
Binary Search Tree (W.C.)	$\Theta(n)$	$\Theta(n)$	$\Theta(n)$
Binary Search Tree (average)	$\Theta(\log n)$	$\Theta(\log n)$	$\Theta(\log n)$

Tree Height

```
treeHeight(root){
    height = 0;
    if (root.left != Null){
        height = max(height, treeHeight(root.left));
    }
    if (root.right != Null){
        height = max(height, treeHeight(root.right));
    }
    return height;
}
```

More Tree “Vocab”



- Traversal:
  - An algorithm for “visiting” every node in a tree
- Pre-Order Traversal:
  - Root, Left Subtree, Right Subtree

- In-Order Traversal:
  - Left Subtree, Root, Right Subtree

- Post-Order Traversal
  - Left Subtree, Right Subtree, Root

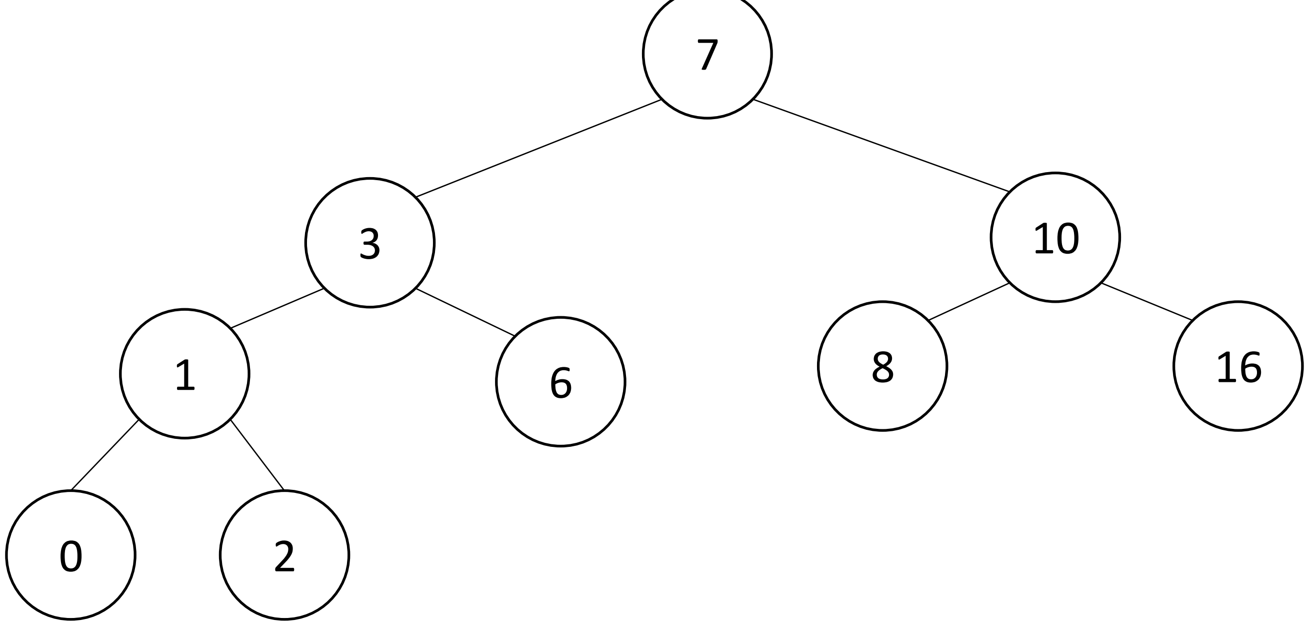
Name that Traversal!

```
AorderTraversal(root){
        if (root.left != Null){
                process(root.left);
        }
        if (root.right != Null){
                process(root.right);
        }
        process(root);
}
```

```
BorderTraversal(root){
        process(root);
        if (root.left != Null){
                process(root.left);
        }
        if (root.right != Null){
                process(root.right);
        }
}
```

```
CorderTraversal(root){
        if (root.left != Null){
                process(root.left);
        }
        process(root)
        if (root.right != Null){
                process(root.right);
        }
}
```

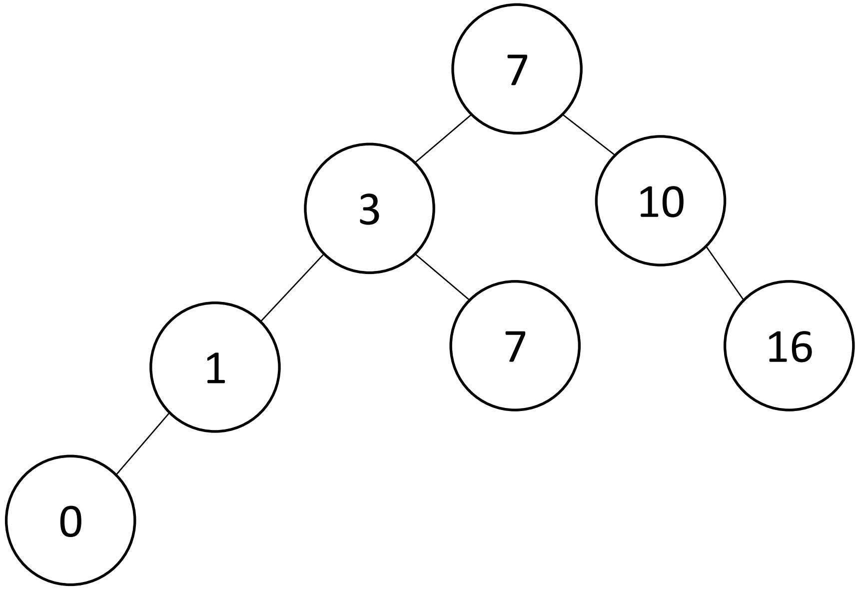
Binary Search Tree

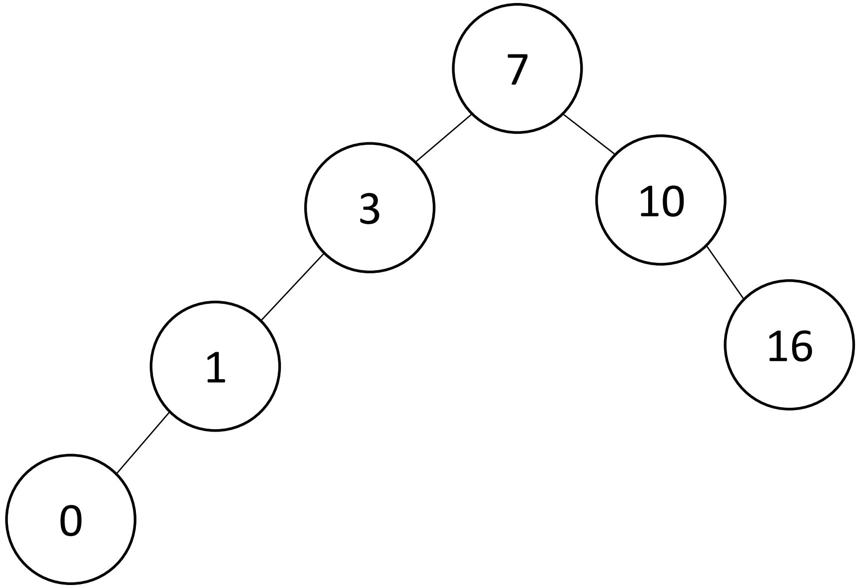


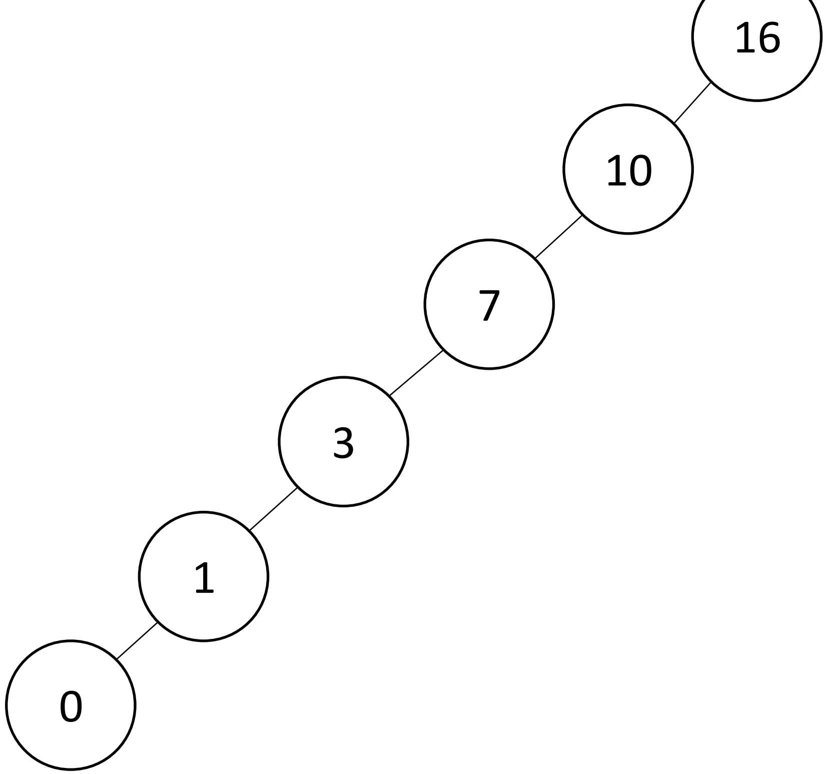
- Binary Tree
  - Definition:
- Order Property
  - All keys in the left subtree are smaller than the root
  - All keys in the right subtree are larger than the root

- Why?

Are these BSTs?

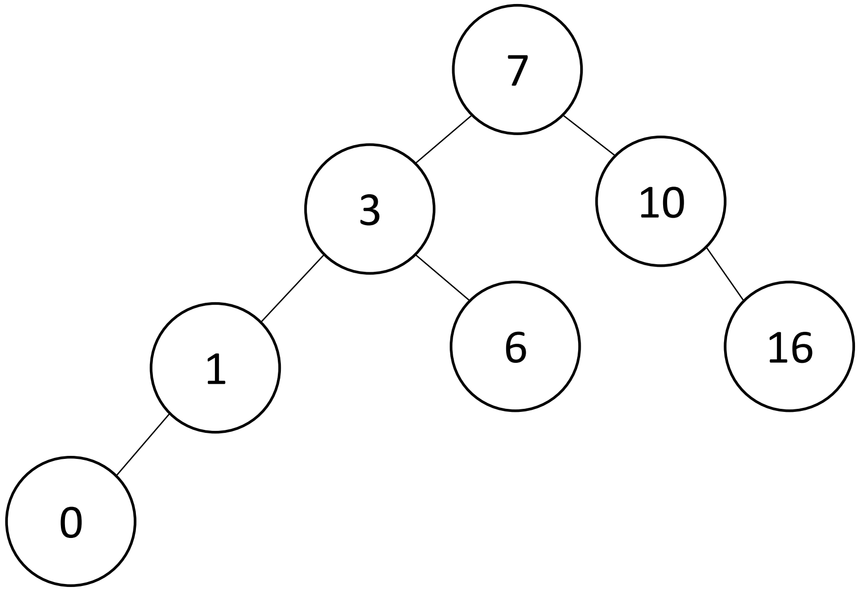






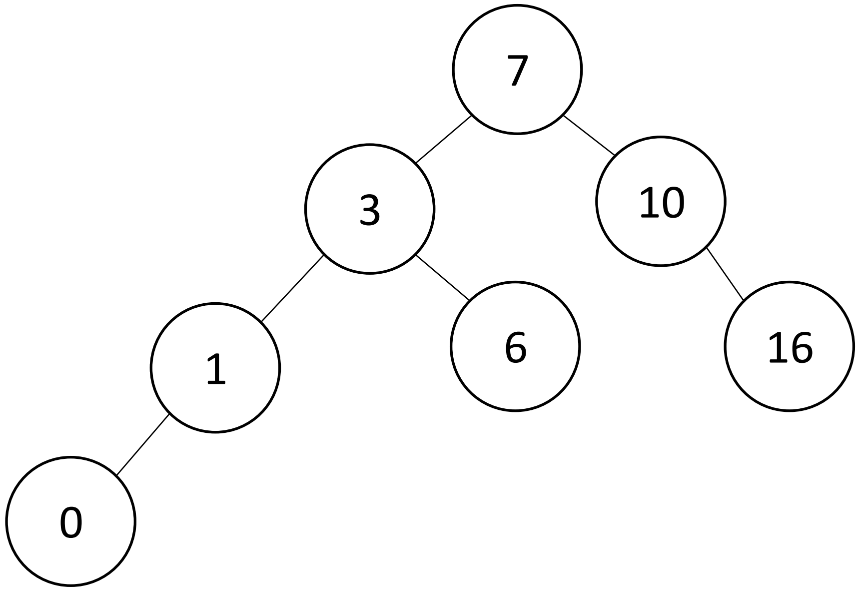
Find Operation (recursive)

```
find(key, root){
        if (root == Null){
                return Null;
        {
        if (key == root.key){
                return root.value;
        }
        if (key < root.key){
                return find(key, root.left);
        }
        if (key > root.key){
                return find(key, root.right);
        }
        return Null;
}
```

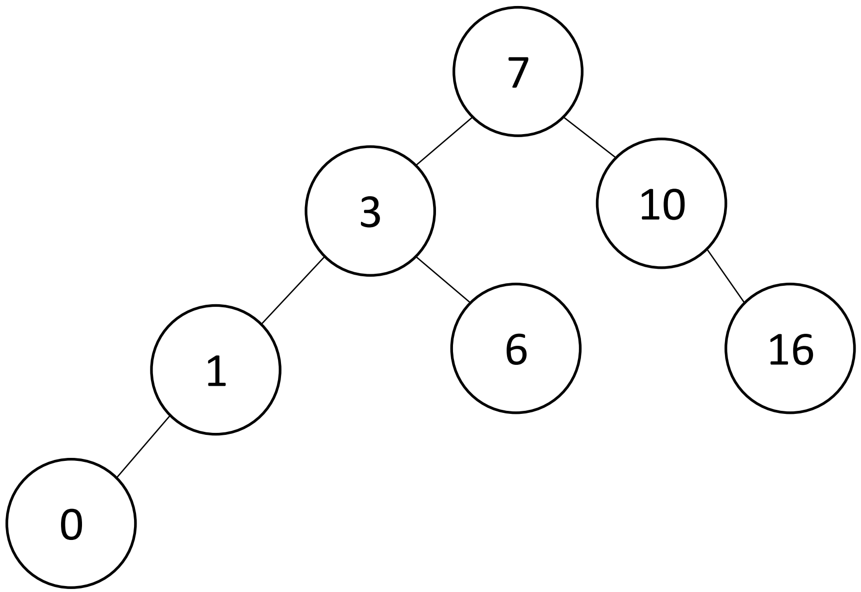


Find Operation (iterative)

```
find(key, root){
        while (root != Null && key != root.key){
                if (key < root.key){
                        root = root.left;
                }
                else if (key > root.key){
                        root = root.right;
                }
        }
        if (root == Null){
                return Null;
        }
        return root.value;
}
```



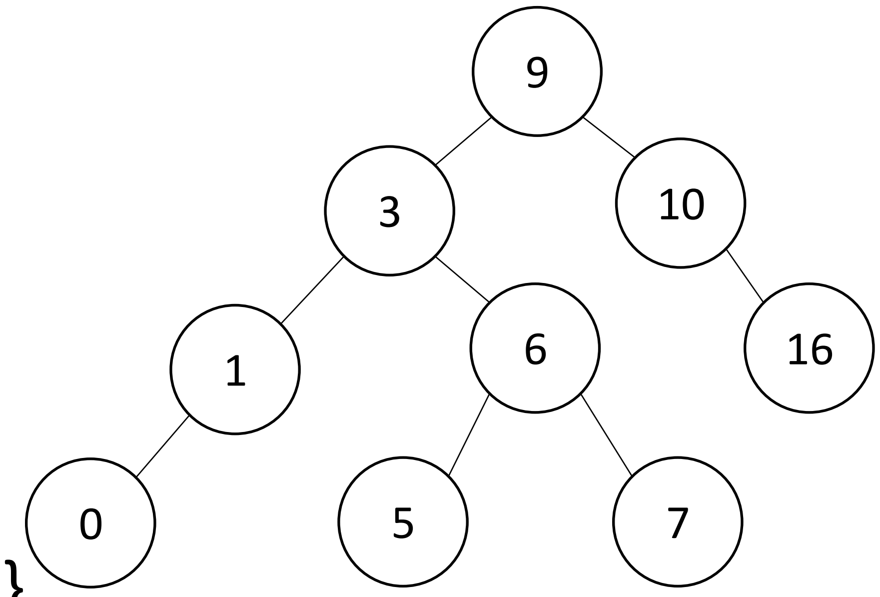
Insert Operation (iterative)



```
insert(key, value, root){
    if (root == Null){ this.root = new Node(key, value); }
    parent = Null;
    while (root != Null && key != root.key){
        parent = root;
        if (key < root.key){ root = root.left; }
        else if (key > root.key){ root = root.right; }
    }
    if (root != Null){ root.value = value; }
    else if (key < parent.key){ parent.left = new Node(key, value); }
    else{ parent.right = new Node (key, value); }
}
```

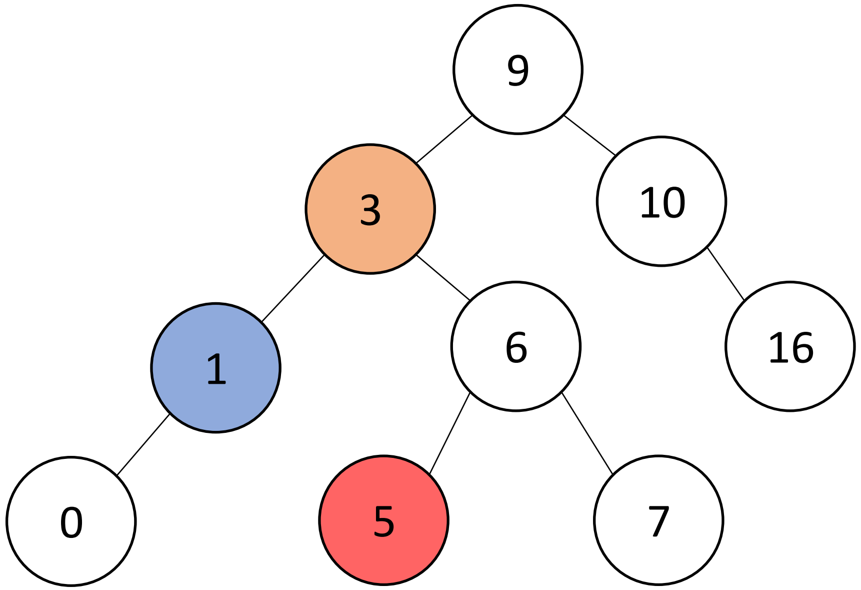
Note: Insert happens only at the leaves!

Delete Operation (iterative)



```
delete(key, root){
        while (root != Null && key != root.key){
                if (key < root.key){ root = root.left;      }
                else if (key > root.key){ root = root.right; }
        }
        if (root == Null){ return; }
        // Now root is the node to delete, what happens next?
}
```

Delete – 3 Cases



- 0 Children (i.e. it's a leaf)
- 1 Child
- 2 Children

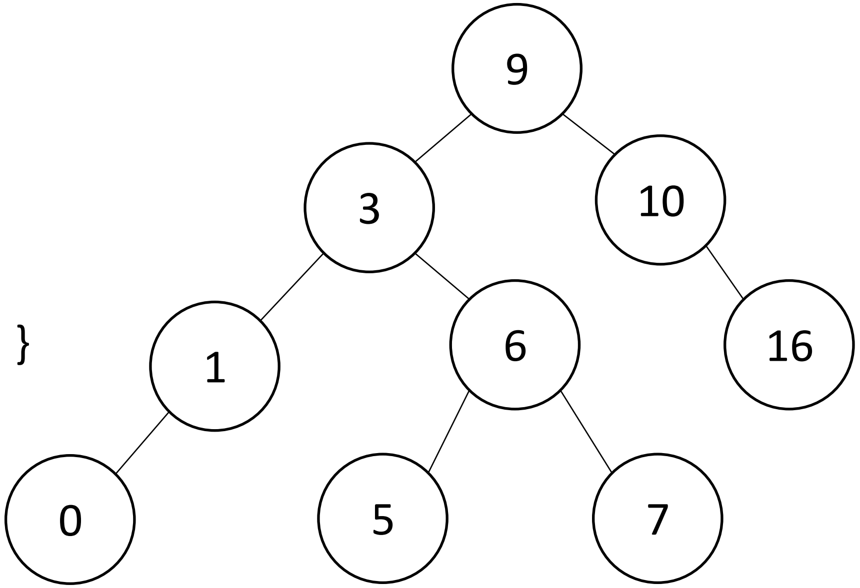
Finding the Max and Min

- Max of a BST:
    - Right-most Thing

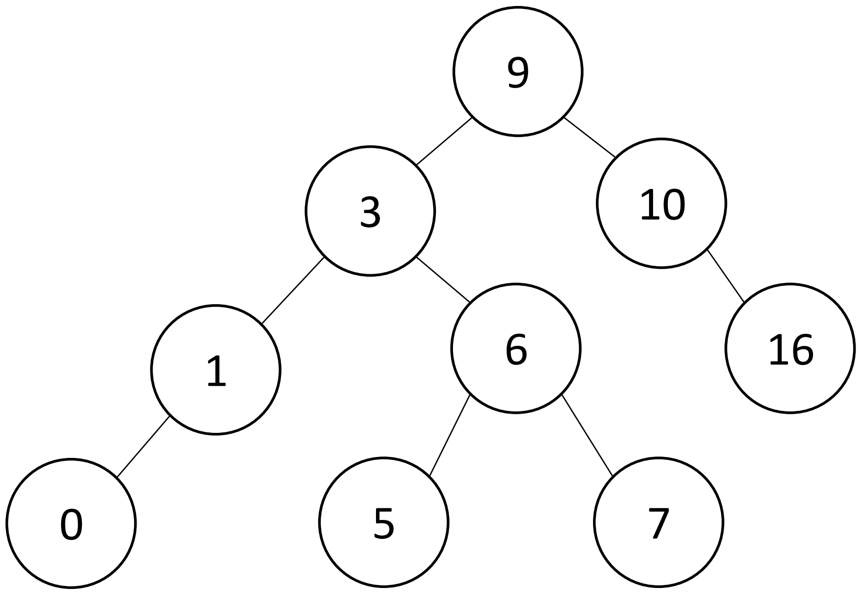
```
maxNode(root){
        if (root == Null){ return Null; }
        while (root.right != Null){
                root = root.right;
        }
        return root;
}
```

- Min of a BST:
    - Left-most Thing

```
minNode(root){
        if (root == Null){ return Null; }
        while (root.left != Null){
                root = root.left;
        }
        return root;
}
```



Delete Operation (iterative)



```
delete(key, root){
        while (root != Null && key != root.key){
                if (key < root.key){ root = root.left;  }
                else if (key > root.key){ root = root.right; }
        }
        if (root == Null){ return; }
        if (root has no children){
                make parent point to Null Instead;
        }
        if (root has one child){
                make parent point to that child instead;
        }
        if (root has two children){
                make parent point to either the max from the left or min from the right
        }
}
```

Worst Case Analysis

- For each of Find, insert, Delete:
  - Worst case running time matches height of the tree
- What is the maximum height of a BST with n nodes?

Improving the worst case

- How can we get a better worst case running time?

“Balanced” Binary Search Trees

- We get better running times by having “shorter” trees
- Trees get tall due to them being “sparse” (many one-child nodes)
- Idea: modify how we insert/delete to keep the tree more “full”

Idea 1: Both Subtrees of Root have same # Nodes

Idea 2: Both Subtrees of Root have same height

Idea 3: Both Subtrees of every Node have same # Nodes

Idea 4: Both Subtrees of every Node have same height

Teaser: AVL Tree

- A Binary Search tree that maintains that the left and right subtrees of every node have heights that differ by at most one.
    - Not too weak (ensures trees are short)
    - Not too strong (works for any number of nodes)