## CSE 341 — Scheme Discussion Questions

1. Suppose we evaluate the following Scheme expressions:

```
(define x '(snail clam))
(define y '(octopus squid scallop))
```

Draw box-and-arrow diagrams of the result of evaluating the following expressions. In particular, note carefully what parts of the list are created fresh, and which are shared with the variables x and y.

```
(a) (define a (cons 'geoduck x))
(b) (define b (cons y y))
(c) (define c (append x y))
(d) (define d (cdr y))
```

- 2. Given the variables defined in Question 1, what is the result of evaluating the following expressions?
  - (a) (eq? (car a) 'geoduck)
    (b) (eq? d '(squid scallop))
    (c) (equal? d '(squid scallop))
    (d) (eq? d (cdr y))
- 3. Continuing with the same variables used in Question 1, show how the box-and-arrow diagram changes after evaluating each of the following expressions.

```
(set-car! x 'tuna)
(set-car! y 'dolphin)
(set-cdr! (cdr y) ())
```

4. Suppose we define some functions and variables:

```
(define (test1 n)
  (set-cdr! n ())
  (display n))
(define (test2 n)
  (set! n ())
  (display n))
(define x '(snail clam))
(define y '(octopus squid scallop))
```

What gets printed when we evaluate

(test1 x) (test2 y)

What are the values of x and y afterward?

5. What is the result of evaluating the following Scheme expressions?

- 6. Define a function mylength to find the length of a list.
- 7. Define a recursive function add1 that takes a list of numbers, and returns a new list of numbers, each being 1 plus the original. For example, (add1 ' (10 20 30)) should evaluate to (11 21 31).
- 8. Define a non-recursive version of add1 that uses map and lambda.
- 9. Using map and lambda, define a function averages that accepts two lists of numbers, and returns a list of the average of each pair. For example:

(averages '(1 2 3) '(11 12 13)) => (6 7 8)

10. Aloysius Q. Hacker, 341 student, is puzzled by the following code.

```
(define incr ())
(define get ())
(let ((n 0))
   (set! incr (lambda (i) (set! n (+ n i))))
   (set! get (lambda () n)))
```

Aloysius is unsure how the functions incr and get can possibly work ... if n is in a stack frame, why do incr and get still work correctly even though we're done with evaluating the let?

Is there something special about let at the top level? So Aloysius tries an experiment:

```
(define newincr ())
(define newget ())
(define (test k)
  (let ((n 0))
      (set! newincr (lambda (i) (set! n (+ n i k))))
      (set! newget (lambda () n))))
```

Then he evaluates (test 100). This time the let is embedded in a function, and so (Aloysius reasons) certainly its stack frame will go away when test returns.

What is the result when Aloysius evaluates each of the following expressions in turn?

(newget)
(newincr 10)
(newget)

Explain (or at least make some reasonable hypotheses).

11. Define a tail-recursive version of "map" for 1-argument functions. (Avoid side effects if possible, but use them if necessary.)