CSE 341: Programming Languages

Spring 2005 Lecture 5 — Type synonyms, more pattern-matching, accumulators

Goals

- Contrast type synonyms with new types
- See pattern-matching for built-in "one of" types (not really a concept, but important for ML programming) and "each of" types
- Investigate why accumulator-style recursion can be more efficient

Type synonyms

You can bind a *type name* to a type. Example:

```
type intpair = int * int
```

(We call something else a *type variable*.)

In ML, this creates a *synonym*, also known as a *transparent* type definition. Recursion not allowed.

So a type name is *equivalent* to its definition.

To contrast, the type a datatype binding introduces is not equivalent to any other type (until possibly a later type binding).

Review: datatypes and pattern-matching

Evaluation rules for datatype bindings and case expressions:

datatype t = C1 of t1 | C2 of t2 | \dots | Cn of tn

Adds constructors Ci where Ci v is a value (and Ci has type $ti \rightarrow t$).

case e of p1 => e1 | p2 => e2 | ... | pn => en

- Evaluate e to v
- If pi is the first pattern to *match* v, then result is evaluation of ei in environment extended by the match.
- If C is a constructor of type t1 * ... * tn -> t, then
 C(x1,...,xn) is a pattern that matches C(v1,...,vn) and the match extends the environment with x1 to v1 ... xn to vn.
- Coming soon: many more pattern forms.

Why patterns?

Even without more pattern forms, this design has advantages over functions for "testing and destructing" (e.g., null, hd, and tl):

- easier to check for missing and redundant cases
- more concise syntax by combining "test, destruct, and bind"
- you can easily define testing and destructing in terms of pattern-matching

In fact, case expressions are the preferred way to test variants and extract values from all ML's "one-of" types, including predefined ones ([] and :: just funny syntax).

So: Do not use functions hd, tl, null, isSome, valOf

Teaser: These functions are useful for *passing as values*

Tuple/record patterns

You can also use patterns to extract fields from tuples and records: pattern $\{f1=x1, \ldots, fn=xn\}$ (or $(x1,\ldots,xn)$) matches $\{f1=v1, \ldots, fn=vn\}$ (or $(v1,\ldots,vn)$).

For record-patterns, field-order does not matter.

This is better style than #1 and #foo, and it means you do not (ever) need to write function-argument types.

Instead of a case with one pattern, better style is a pattern directly in a val binding.

Next time: "deep" (i.e., nested) patterns.

Recursion

You should now have the hang of recursion:

- It's no harder than using a loop (whatever that is)
- It's much easier when you have multiple recursive calls (e.g., with functions over ropes or trees)

But there are idioms you should learn for *elegance*, *efficiency*, and *understandability*.

Today: using an *accumulator*.

Accumulator lessons

- Accumulators can avoid data-structure copying
- Accumulators can reduce the depth of recursive calls that are not *tail calls*
- Key idioms:
 - Non-accumulator: compute recursive results and combine
 - Accumulator: use recursive result as new accumulator
 - The base case becomes the initial accumulator

You will use recursion in non-functional languages—this lesson still applies.

Let's investigate the evaluation of to_list_1 and to_list_2.