

## CSE 341 - Programming Languages Midterm - Autumn 2015 - Answer Key

1. (6 points) Suppose the following Racket program has been read in.

```
(define x '(1 2 3))  
(define y '(10 11 12))  
(define z (append (cdr x) (cdr y)))
```

Draw box-and-arrow diagram of the resulting lists, being careful to show correctly what parts are shared and what parts are separate. (If you run out of room, please draw a fresh copy of the diagram on the back of this page instead.)

See answer on a separate page at the end of this answer key.

2. (12 points) Consider the following Haskell function definitions. (Some of these are from the Prelude; some are written from scratch.)

```
filter p xs = [ x | x <- xs, p x ]  
repeat x = x : repeat x  
memb xs a = or (map (==a) xs)  
vowel = memb "aeiou"
```

Here are some possible types for each of these functions. After each type, write G if the type is correct and the most general, C if it is correct but not the most general, and I if it is incorrect.

```
filter :: [a] -> [a] -> [a] I  
filter :: (a -> Bool) -> [a] -> [a] G  
filter :: (Eq a) => (a -> Bool) -> [a] -> [a] C  
filter :: (Num a) => (a -> Bool) -> [a] -> [a] C  
repeat :: a -> a I  
repeat :: a -> [a] G  
repeat :: [a] -> [[a]] C  
memb :: [a] -> a -> Bool I  
memb :: Eq a => [a] -> a -> Bool G  
memb :: Ord a => [a] -> a -> Bool C  
vowel :: Char -> Bool G  
vowel :: String -> Bool I
```

3. (10 points) Consider a Haskell function `remove` that takes an item and a list, and returns a new list, dropping elements that are equal to the item. For example, `remove 3 [1,2,3,4,3,2,1]` evaluates to `[1,2,4,2,1]`, and `remove 3 []` evaluates to `[]`.

(a) What is the most general type for `remove`?  
`remove :: Eq t => t -> [t] -> [t]`

(b) Write a recursive definition of `remove`.

```
remove x [] = []
remove x (y:ys)
  | x==y      = remove x ys
  | otherwise = y : remove x ys
```

(c) Write a non-recursive definition of `remove` using Haskell's `filter` function. (Hint: there is a definition of `filter` in Question 2.)

```
remove x = filter (/=x)
```

4. (5 points) What is the value of `mystery`? (If it's infinite give the first several elements.)

```
mystery = 0 : map (\x->2*x+1) mystery

[0,1,3,7,15,31,63,127,255,511, .....
```

5. (6 points) Define a **curried plus function** `curried-plus` in Racket. So `(curried-plus 1)` should return a function that adds 1 to numbers, and `((curried-plus 1) 3)` should evaluate to 4.

```
(define (curried-plus i)
  (lambda (j) (+ i j)))
```

6. (6 points) Now define a function `extra-spicy-plus` in Racket that is also curried but that works with 3 numbers instead of 2. So `((extra-spicy-plus 1) 2) 3)` should evaluate to 6.

```
(define (extra-spicy-plus i)
  (lambda (j) (lambda (k) (+ i j k))))
```

7. (6 points) Consider the following Racket expressions. (They are identical except that the first uses `let`, the second uses `let*`, and the third uses `letrec`.) What does each one evaluate to?

```
(let ([y 1]
      [f (lambda (x) (+ x 1))])
  (let ([y 5]
        [f (lambda (x) (if (> x 5) x (f (+ x y)))]))
    (f 3)))
```

=> 5

```
(let ([y 1]
      [f (lambda (x) (+ x 1))])
  (let* ([y 5]
         [f (lambda (x) (if (> x 5) x (f (+ x y)))]))
    (f 3)))
```

=> 9

```
(let ([y 1]
```

```

      [f (lambda (x) (+ x 1))])
(letrec ([y 5]
         [f (lambda (x) (if (> x 5) x (f (+ x y))))])
  (f 3)))

```

=> 8

8. (5 points) Consider the following Racket program.

```

(define x 1)
(define z 2)

(define (add-it y)
  (+ x y))

(define (test)
  (let ([x 10]
        [z 20])
    (add-it z)))

```

(a) What is the result of evaluating `(test)`?

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(b) Suppose Racket used dynamic scoping. What would be the result of evaluating `(test)`?

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9. (8 points) The lecture notes for Racket macros included a cosmetic macro for `my-if` (which just provides a different syntax for `if`).

```

(define-syntax my-if          ; macro name
  (syntax-rules (then else) ; literals it uses, if any
    [(my-if e1 then e2 else e3) ; pattern
     (if e1 e2 e3)])          ; template

```

Show the code to add `my-if` to `OCTOPUS`. For full credit, do this by rewriting the `my-if` expression as a normal `if` and then calling `eval` on the result, in the same way that the Racket `my-if` macro produces new Racket source code that is then interpreted.

Hint: you should write another case for the `eval` function. Here are two example cases from the starter file:

```

-- A quoted expression evaluates to that expression.
eval (OctoList [OctoSymbol "quote", x]) env = x

```

```

{- An expression starting with (lambda ...) evaluates to a closure,
where a closure consists of a list of variable names (OctoSymbols),
the environment of definition, and the body. -}
eval (OctoList [OctoSymbol "lambda", OctoList vars, body]) env =
  OctoClosure vars env body

```

```
eval (OctoList [OctoSymbol "my-if", test, OctoSymbol "then", truebranch,
               OctoSymbol "else", falsebranch]) env =
  eval (OctoList [OctoSymbol "if", test, truebranch, falsebranch]) env
```

10. (8 points) Write a slightly silly Racket macro called `const` that does something like `const` in Haskell: the result of evaluating `(const exp1 exp2)` should be the result of evaluating `exp1` (just throw `exp2` away without evaluating it). Unlike Haskell, though, evaluate `exp1` each time the `const` expression is evaluated. For example, `(const (+ 2 3) (/ 1 0))` should evaluate to 5. (There is no divide-by-zero error since we don't evaluate the second expression.)

```
(define-syntax const
  (syntax-rules ()
    [(const k expr)
     k]))
```

11. (8 points) Convert the following Haskell action into an equivalent one that doesn't use `do`.

```
echo = do
  x <- readLn
  y <- readLn
  putStr "the sum is "
  putStrLn (show (x+y))
```

```
echo = readLn >>= \x -> readLn >>= \y -> putStr "the sum is " >>
  putStrLn (show (x+y))
```

Question 1

