

CSE341: Programming Languages

Lecture 25
Subtyping for OOP;
Comparing/Combining Generics and
Subtyping

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Now...

Use what we learned about subtyping for records and functions to understand subtyping for class-based OOP

- Like in Java/C#

Recall:

- Class names are also types
- Subclasses are also subtypes
- Substitution principle: Instance of subclass should usable in place of instance of superclass

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An object is...

- · Objects: mostly records holding fields and methods
 - Fields are mutable
 - Methods are immutable functions that also have access to self
- So could design a type system using types very much like record types
 - Subtypes could have extra fields and methods
 - Overriding methods could have contravariant arguments and covariant results compared to method overridden
 - Sound only because method "slots" are immutable!

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Actual Java/C#...

Compare/contrast to what our "theory" allows:

- 1. Types are class names and subtyping are explicit subclasses
- 2. A subclass can add fields and methods
- 3. A subclass can override a method with a covariant return type
 - (No contravariant arguments; instead makes it a nonoverriding method of the same name)
- (1) Is a subset of what is sound (so also sound)
- (3) Is a subset of what is sound and a different choice (adding method instead of overriding)

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Classes vs. Types

- · A class defines an object's behavior
 - Subclassing inherits behavior and changes it via extension and overriding
- · A type describes an object's methods' argument/result types
 - A subtype is substitutable in terms of its field/method types
- · These are separate concepts: try to use the terms correctly
 - Java/C# confuse them by requiring subclasses to be subtypes
 - A class name is both a class and a type
 - Confusion is convenient in practice

Optional: More details

Java and C# are sound: They do not allow subtypes to do things that would lead to "method missing" or accessing a field at the wrong type $\frac{1}{2} \left(\frac{1}{2} \right) = \frac{1}{2} \left(\frac{1}{2} \right) \left(\frac{1}{2}$

Confusing (?) Java example:

- Subclass can declare field name already declared by superclass
- Two classes can use any two types for the field name
- Instance of subclass have two fields with same name
- "Which field is in scope" depends on which class defined the method

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self/this is special

- · Recall our Racket encoding of OOP-style
 - "Objects" have a list of fields and a list of functions that take self as an explicit extra argument
- So if self/this is a function argument, is it contravariant?
 - No, it is covariant: a method in a subclass can use fields and methods only available in the subclass: essential for OOP

```
class A {
  int m(){ return 0; }
}
class B extends A {
  int x;
  int m(){ return x; }
}
```

- Sound because calls always use the "whole object" for self
- This is why coding up your own objects manually works much less well in a statically typed languages

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What are generics good for?

Some good uses for parametric polymorphism:

· Types for functions that combine other functions:

```
fun compose (g,h) = fn x => g (h x)
(*compose: ('b->'c) * ('a->'b) -> ('a->'c) *)
```

· Types for functions that operate over generic collections

```
val length : 'a list -> int
val map : ('a -> 'b) -> 'a list -> 'b list
val swap : ('a * 'b) -> ('b * 'a)
```

- · Many other idioms
- General point: When types can "be anything" but multiple things need to be "the same type"

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Generics in Java

- Java generics a bit clumsier syntactically and semantically, but can express the same ideas
 - Without closures, often need to use (one-method) objects
 - See also earlier optional lecture on closures in Java/C
- Simple example without higher-order functions (optional):

```
class Pair<T1,T2> {
   T1 x;
   T2 y;
   Pair(T1 _x, T2 _y){ x = _x; y = _y; }
   Pair<T2,T1> swap() {
      return new Pair<T2,T1>(y,x);
   }
   ...
}
```

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Subtyping is not good for this

- · Using subtyping for containers is much more painful for clients
 - Have to downcast items retrieved from containers
 - Downcasting has run-time cost
 - Downcasting can fail: no static check that container holds the type of data you expect
 - (Only gets more painful with higher-order functions like map)

```
class LamePair {
  Object x;
  Object y;
  LamePair(Object _x, Object _y){ x=_x; y=_y; }
  LamePair swap() { return new LamePair(y,x); }
}
// error caught only at run-time:
String s = (String)(new LamePair("hi",4).y);
```

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What is subtyping good for?

Some good uses for subtype polymorphism:

- Code that "needs a Foo" but fine to have "more than a Foo"
- · Geometry on points works fine for colored points
- GUI widgets specialize the basic idea of "being on the screen" and "responding to user actions"

Awkward in ML

ML does not have subtyping, so this simply does not type-check:

```
(* {x:real, y:real} -> real *)
fun distToOrigin ({x=x,y=y}) =
   Math.sqrt(x*x + y*y)
val five = distToOrigin {x=3.0,y=4.0,color="red"}
```

Cumbersome workaround: have caller pass in getter functions:

- And clients still need different getters for points, color-points

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Wanting both

- · Could a language have generics and subtyping?
 - Sure!
- · More interestingly, want to combine them
 - "Any type T1 that is a subtype of T2"
 - Called bounded polymorphism
 - Lets you do things naturally you cannot do with generics or subtyping separately

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Example

Method that takes a list of points and a circle (center point, radius)

- Return new list of points in argument list that lie within circle

Basic method signature:

Java implementation straightforward assuming Point has a distance method:

```
List<Point> result = new ArrayList<Point>();
for(Point pt : pts)
  if(pt.distance(center) < r)
    result.add(pt);
return result;</pre>
```

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Subtyping?

- Would like to use inCircle by passing a List<ColorPoint> and getting back a List<ColorPoint>
- Java rightly disallows this: While inCircle would "do nothing wrong" its type does not prevent:
 - Returning a list that has a non-color-point in it
 - Modifying pts by adding non-color-points to it

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Generics?

· We could change the method to be

```
<T> List<T> inCircle(List<T> pts,
Point center,
double r) { ... }
```

- Now the type system allows passing in a List<Point> to get a List<Point> returned or a List<ColorPoint> to get a List<ColorPoint> returned
- But cannot implement inCircle properly: method body should have no knowledge of type T

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Bounds

· What we want:

```
<T> List<T> inCircle(List<T> pts,
Point center,
double r) where T <: Point
{ ... }
```

- Caller uses it generically, but must instantiate T with some subtype of Point (including Point)
- Callee can assume T <: Point so it can do its job
- Callee must return a List<T> so output will contain only elements from pts

Real Java

• The actual Java syntax:

- Note: For backward-compatibility and implementation reasons, in Java there is actually always a way to use casts to get around the static checking with generics ®
 - With or without bounded polymorphism

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