

CSE 341:

*Programming Languages*

Section AC with Nate Yazdani

# about me

- CSE BS/MS student (last year of BS)
- I am *really* into programming languages
- I am *really* into research on programming languages
  - program synthesis
  - formal verification
  - crazy theoretical stuff (“homotopy type theory??”)



about you?

# why are we here

- to get a bit more interactive learning
- to supplement the material from lecture
- to take a closer look at important but subtle details
- to ask questions (please)! :-)

# agenda

- ML development workflow
  - emacs
  - using **use**
  - REPL
- some more ML
  - variable shadowing
  - debugging with the REPL
  - boolean operations
  - comparison operations

# emacs

- recommended (not required) editor for this course
- a powerful tool for programming
- learning curve may seem steep, but you get the hang of it more quickly than you'd think
- Dan's emacs guide is super helpful
- if you need help with setup, please let us know

emacs demo time

# using `use`

```
use "foo.sml";
```

- parses the local file `foo.sml` and then evaluates the bindings one after another
- result is the dummy value `()`
  - automatically bound to variable `it`
  - completely safe to ignore



# the REPL

- stands for the “read-eval[uate]-print-loop”
  - it reads, evaluates, prints, and loops!
- works with both expressions and bindings
  - expects semicolons to know when to evaluate
- handy to quickly try stuff out
  - in emacs, start with C-c C-s and end with C-d
- as we will see in a bit, **use**-ing multiple files without restarting your REPL session is dangerous

# shadowing

```
val a = 1; (* a -> 1 *)  
val b = a * 10; (* a -> 1, b -> 10 *)  
val a = 42; (* a -> 1, b -> 10, a -> 42 *)
```

- eager” evaluation of expressions in variable bindings
  - computes the value and *then* binds the name to that value
  - afterwards, the original expression is forgotten
- multiple variable bindings to the same variable name is called “**shadowing**”
  - affects both static and dynamic environments
  - ML will use the most-recently bound value in the current environment
- remember: there is no variable “assignment” in ML
  - you can only shadow it in a later environment
  - once bound, a variable’s value is an immutable constant

# avoid shadowing

- it can confuse yourself and (especially) others
- it's often considered poor style
- why? shadowing variables in a REPL session may
  - make *wrong* code seem *correct*
  - make *correct* code seem *wrong*
  - this can easily happen when you re-**use** a file

# using a shadowed variable

- is it ever possible to use a shadowed variable?
  - yes!
  - and also no...
- when the shadowing binding of a variable name goes out of **scope**, the shadowed binding is available again
  - environments are like a “lookup stack”

```
val x = "Hello World";  
fun plus1 (x : int) = x + 1;  
val y = plus1 2;  
val z = x ^ "!!"; (* ..., z -> "Hello World!!" *)
```

# be careful with **use**

- **warning:** variable shadowing makes it dangerous to call **use** multiple times without restarting the REPL session
- it *might* be safe to call **use** more than once in the same REPL session, but think twice about it
  - at the beginning of a session, loading distinct files with distinct variable names is probably fine
  - while the behavior of **use** is well-defined, even experts can easily get confused
- best to always restart the REPL session

# debugging errors

- your mistake could be
  - syntactic: the source code means nothing (not in the ML grammar) or something unintended

```
val 0 = x
```

- typing: the code fails to typecheck

```
3 + true
```

- semantic (evaluation): the program's behavior is not what you want, *e.g.*, raises an exception, computes the wrong value, or loops infinitely

```
val three = 2 + 2
```

- keep these straight when debugging
- sometimes one kind of mistake will appear to be another

# play around

- best way to learn something: try lots of things and don't be afraid of errors
- work on developing resilience to mistakes
  - slow down
  - don't panic
  - read what you wrote very carefully
  - reconsider what assumptions you're making
- maybe it will help to see me make some mistakes?

let's give it a try



# boolean operations

operation	syntax	type-checking	semantics (evaluation)
conjunction	<b>e1</b> <b>andalso</b> <b>e2</b>	<b>e1</b> and <b>e2</b> must have type <b>bool</b>	same as <b>&amp;&amp;</b> in Java
disjunction	<b>e1</b> <b>orelse</b> <b>e2</b>	<b>e1</b> and <b>e2</b> must have type <b>bool</b>	same as <b>  </b> in Java
negation	<b>not</b> <b>e</b>	<b>e</b> must have type <b>bool</b>	same as <b>!e</b> in Java

- **not** is essentially just a pre-defined function
- **andalso** and **orelse** must be built in, because they “short-circuit” and may not always evaluate **e2**
- be careful to not use **and** instead of **andalso**
  - they mean totally different things

# booleans with *style*

- ML does not “need” **andalso**, **orelse**, or **not**

```
(* e1 andalso e2 *)  
if e1  
then e2  
else false
```

```
(* e1 orelse e2 *)  
if e1  
then true  
else e2
```

- more concise forms are generally better style
- definitely please don't do this

```
(* just say e (!!!)” *)  
if e then true else false
```

# comparisons

- you can compare two **int** values with  
= <> > < >= <=
- you might get weird error messages because these operators work with some other types too
- > < >= <= also work with two real values but *not* with one **int** and one **real**
- = <> work with any two values of the same “equality type” but not with **real**
  - we’ll hear more about equality types later

thanks!