

CUTS

P.50

- A cut prunes or “cuts out” an unexplored part of a Prolog search tree.
- Cuts can make a computation more efficient by eliminating futile search and backtracking.
- Cuts are controversial because they are impure.
- A cut is written as “!”.

When a rule

$B :- C_1, \dots, C_{j-1}, !, C_{j+1}, \dots, C_k$

is applied, the cut tells control to backtrack past

C_{j-1}, \dots, C_1 , and B without considering any more rules for them.

Example with CUTS

P.51

```
age(leah, 48) .
age(natalie, 30) .
age(octavia, 34) .
age(darrell, 59) .
age(michael, 8) .
age(sue, 15) .
age(sylvia, 81) .
age(loren, 29) .
age(lura, 87) .
age(ron, 60) .
```

```
blond(leah) .
blond(natalie) .
blond(octavia) .
brunette(darrell) .
brunette(michael) .
redhair(sylvia) .
redhair(loren) .
redhair(sue) .
grayhair(lura) .
grayhair(ron) .
```

```
cast(X) :- age(X, A), satisfactory(X, A) .
satisfactory(X, A) :- between(0, 10, A), !, blond(X) .
satisfactory(X, A) :- between(11, 20, A), !, redhair(X) .
satisfactory(X, A) :- between(20, 50, A), !, brunette(X) .
satisfactory(X, A) :- between(50, 90, A), !, grayhair(X) .
```

This eliminates some needless search.

Cut + Fail achieve Negation	P.52
<p>not(X) :- X, !, fail not(_) .</p> <p>Fail is a system predicate that fails. _ is a wild-card variable.</p> <p>The first rule attempts to satisfy X . If X fails, then the second rule succeeds, because _ unifies with any term.</p> <p>If X succeeds, then the fail predicate forces failure, and the cut prevents consideration of the second rule.</p> <p>Note that if not (X) succeeds, it merely means that X is not provable according to the database.</p> <p>X may or may not be actually false.</p>	

Another Cut/Fail Combination Example	P.53
<p>allow(elephant) :- !, fail .</p> <p>allow(Animal) :- size(Animal, lessthan50), license(Animal).</p> <p>allow(Animal) :- lives(Animal, cage) .</p> <p>meaning</p> <p>If an animal is not an elephant and either weights less than 50 pounds and has a license or lives in a cage, it is allowed.</p> <p>Elephants, even small ones that live in cages, are not allowed.</p>	

Gathering Answers into Bags or Sets

P.54

The predicates `bagof` and `setof` are used to gather instances of objects.

We specify a goal, a variable in the goal, and a bag or set name.

For each success of the goal, the constant that matched this variable is gathered into the bag or set.

Example

```
parent( jan, bet ) .
parent( jan, cat ) .
parent( joe, ann ) .
parent( joe, cat ) .
```

```
|?- bagof( Child, parent( jan, Child ), B ) .
B = [ bet, cat ]
```

```
|?- bagof( Child, Who^( parent ( Who, Child ) ), B ) .
B = [ bet, cat, ann, cat ]
```

→ read "there exists Who"

```
|?- setof( Child, P^( parent ( P, Child ) ), S ) .
S = [ ann, bet, cat ]
```

Dynamic Knowledge Assertion/Retraction

P.55

Prolog provides built in functions to work with Horn Clauses.

You can

- 1) Construct a structure representing a clause
- 2) add a clause to the database
- 3) remove a clause from the database

* All Prolog structures have the form
functor (arguments)

Facts are already in this form. To convert a rule to this form

$$P(X_1, \dots, X_n) \text{ :- } Q_1(X_1, \dots, X_n), Q_a(\dots), \dots, Q_x(\dots)$$

converts to

$$\text{' :- ' (P(X}_1, \dots, X_n), \text{' , ' (Q}_1(\dots), Q_a(\dots), \dots, Q_x(\dots)))$$

example: `' :- '(cat(X) , ' , '(animal(X), furry(X)))`

Some Utilities for Dynamic Knowledge

P.56

read/write

read(T) reads a term **T** from the input stream.

write(T) writes a term **T** to the output stream.

listing

listing(A) writes out all clauses with atom **A** as their predicate to the output stream.

functor

functor(T, F, N) succeeds if **T** is a structure with functor **F** and arity **N**. (If **T** is a variable, it constructs such a structure.)

arg

arg(Num, T, Argument) puts **Argument** into structure **T** as argument number **Num**.

assert

P.57

assert $\left\{ \begin{array}{l} q \\ z \end{array} \right\}(C)$ adds clause **C** to the database at the
{ beginning
end .

retract

retract(C) removes the first clause that matches **C** from the database.

Example

```
new_fact :- read( A1), read( A2 ), read( A3 ),
            functor( C, A3, 2 ),
            arg( 1, C, A1 ),
            arg( 2, C, A2 ),
            assert( C ) .
```

This rule reads 3 terms; uses functor to set up a structure named **C** with **A3** as its predicate, and room for 2 arguments; uses arg to make **A1** and **A2** the arguments; and asserts it.

P.58

```
|?- new_fact.
```

```
 |: bob.
```

```
 |: mike.
```

```
 |: father.
```

```
yes
```

```
|?- new_fact.
```

```
 |: mauro.
```

```
 |: nick.
```

```
 |: father.
```

```
yes
```

```
|?- listing( father ).
```

```
  father( bob, mike ).
```

```
  father( mauro, nick ).
```

```
yes
```

P.59

Call

A call event occurs when Prolog starts trying to satisfy a goal.

You can also invoke call dynamically, like assert.

Example

```
check_fact :- read( B1 ), read( B2 ), read( B3 ),
               functor( D, B3, 2 ),
               arg( 1, D, B1 ),
               arg( 2, D, B2 ),
               call( D ).
```

```
|?- check_fact.
```

```
 |: bob.
```

```
 |: mike.
```

```
 |: father.
```

```
yes
```

```
|?- check_fact.
```

```
 |: mauro.
```

```
 |: mike.
```

```
 |: father.
```

```
no
```

The Univ Operator =..

This is the easiest and clearest way to construct dynamic assertions and calls.

-- The predicate `f(a, b, c)` corresponds to the list `[f, a, b, c]`.

-- The operator `=..` converts back and forth between the two representations.

```
?- f(a, b, c) =.. X .
   X = [ f, a, b, c ]
   yes
```

```
?- X =.. [ w, x, y, z ] .
   X = w(x, y, z) .
   yes
```

Using =.. To Construct Dynamic Calls

```
mother(linda, sylvia) .
father(linda, aaron) .
```

```
answer_questions :-
    write('mother or father?') ,
    read(X) ,
    write('of whom?') ,
    read(Y) ,
    Q =.. [X, Y, Who] ,
    call(Q) ,
    write(Who) ,
    nl .
```

```
1 ?- answer_question.
mother or father? mother .
of whom? linda .
sylvia
Yes
```

```
2 ?- answer_question.
mother or father? father .
of whom? linda.
aaron
Yes
```

Using =.. To Construct Dynamic Asserts

P.62

```
fact :- F =.. [dog, sierra],
      assert(F),
      write(ok),
      nl.

rule :- R =.. [ '- ', animal(X), dog(X)],
      assert(R),
      write(ok),
      nl.

comprule :- C =.. [ ', ', dog(X), waggingtail(X)],
           S =.. [ '- ', friendly(X), C],
           assert(S),
           write(ok),
           nl.

2 ?- fact .                ok   yes
3 ?- rule .                ok   yes
4 ?- comprule .           ok   yes
5 ?- consult(user).
|: waggingtail(sierra) .

6 ?- dog(Who) .
Who = sierra

8 ?- friendly(Who) .
Who = sierra
```

Clause

P.63

Clause provides another way of selecting Horn clauses.

Clause(X, Y) succeeds if it can match X and Y to the head and body of an existing clause in the database.

X must be instantiated enough so that the main predicate is known.
Only works for dynamically asserted clauses!!

Example

```
list1( X ) :- clause( X, Y ),
             output_clause( X, Y ),
             write( '· ' ), nl, fail.
```

```
list1( X ).
```

```
output_clause( X, true ) :- !, write( X ).
```

```
output_clause( X, Y ) :- write( ( X :- Y ) ).
```

Note that for facts, the tail is true.

```
Ex.   assert( q ( a, b ) ).
```

```
list1( q ( V1, V2 ) ).
```

Parsing Simple English Sentences

article(a). article(the). adjective(giant).
 preposition(on). preposition(from).
 verb(rose). verb(sat). verb(was).
 noun(cat). noun(rocket). noun(mat). noun(pad).

sentence(X) :- np(X, R), vp(R, []).
 np([X,Y|Z], Z) :- article(X), noun(Y).
 vp([X|Y], R) :- verb(X), pp(Y, R).
 pp([X|Y], Z) :- preposition(X), np(Y, Z).

|?- sentence([the, cat, sat, on, the, mat]).
 |?- sentence([the, rocket, was, on, the, pad]).
 |?- sentence([the, mat, was, on, the, cat]).
 |?- sentence([the, rocket, rose]).
 |?- sentence([the, giant, cat, rose, from, the, mat]).