

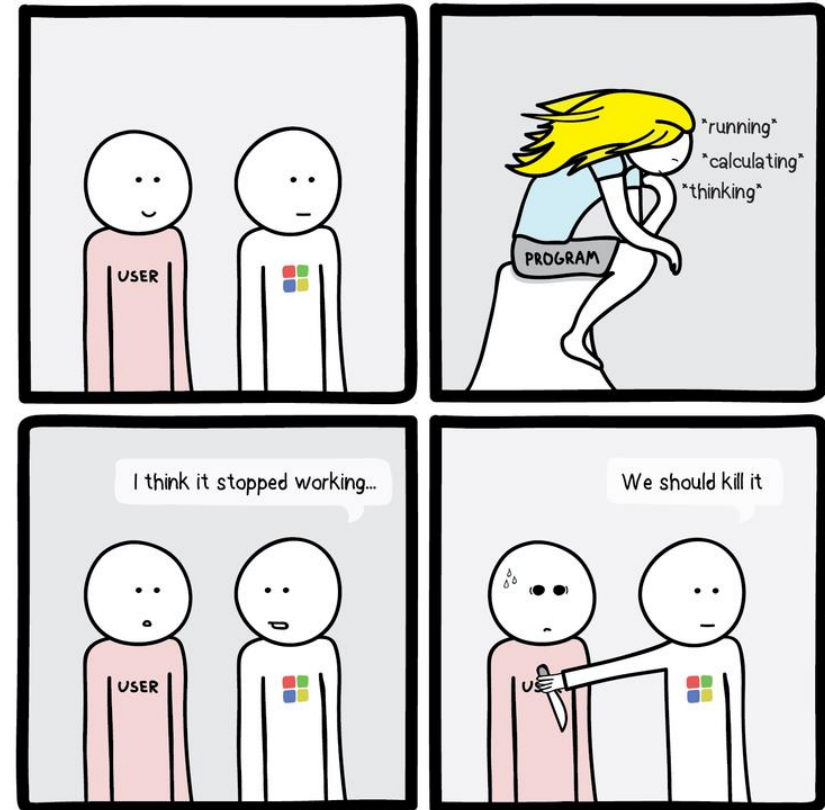
# Processes

## CSE 351 Winter 2024

**Instructor:**  
Justin Hsia

**Teaching Assistants:**

Adithi Raghavan  
Aman Mohammed  
Connie Chen  
Eyoel Gebre  
Jiawei Huang  
Malak Zaki  
Naama Amiel  
Nathan Khuat  
Nikolas McNamee  
Pedro Amarante  
Will Robertson



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# Relevant Course Information

- ❖ HW20 due tonight, HW21 due Friday, HW22 due Monday
- ❖ Lab 4 due Friday
- ❖ Lab 5 due next Friday, 3/8
  - Section this week is to get your started with Lab 5
  - Can use one late day; must be submitted by Sunday, 3/10
- ❖ Final March 11-13, regrade requests only Monday, March 18

# Winter 2024 Crunch

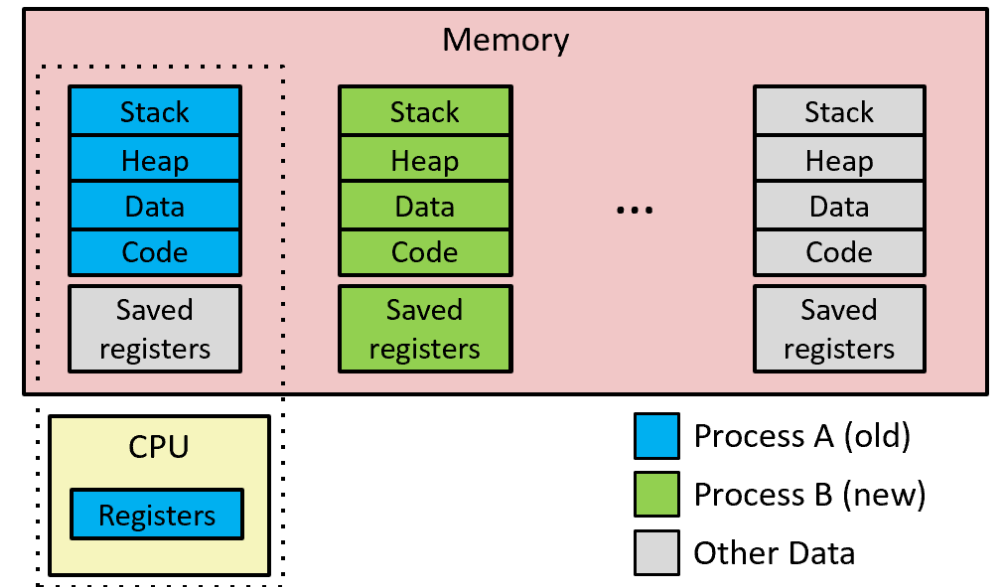
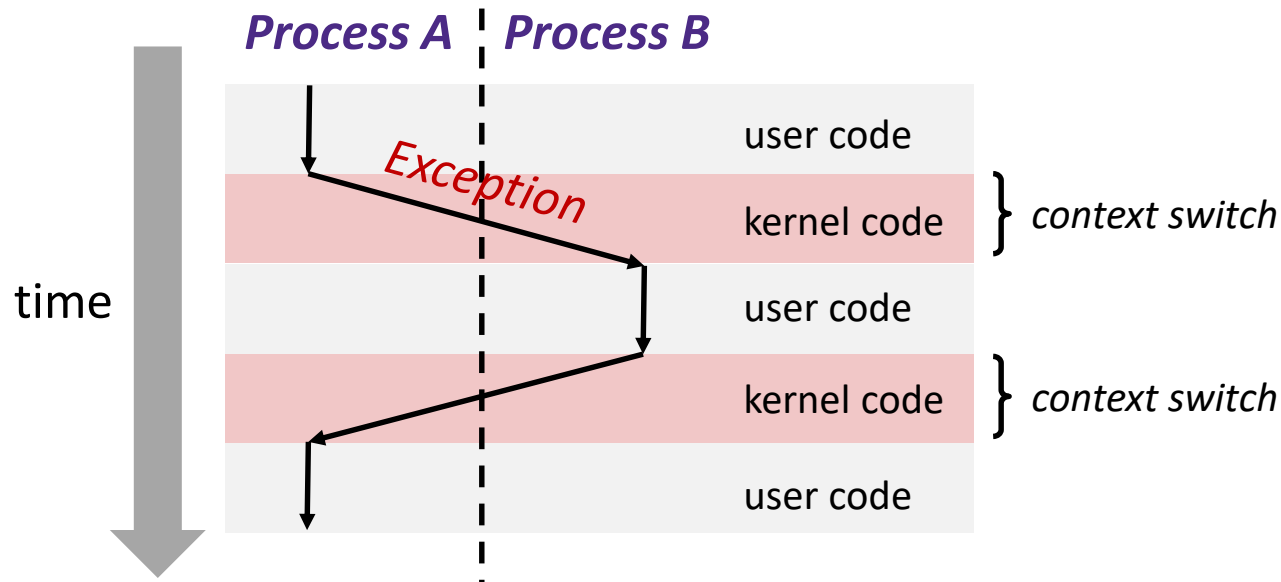
- ❖ This quarter is unusually short
  - Winter is always the shortest of the year for MWF classes due to Monday holidays
  - This year, we also lost the first Monday due to New Year's
- ❖ 29 → 26 lectures compared against 23au
  - Condensed “x86-64 Programming” from 4 lessons to 3
  - Cut “Exceptional Control Flow” (related to Processes)
  - Cutting “C and Java”
- ❖ Assignments compressed, too
  - Cut a number of homework questions throughout the quarter
  - Less time than usual to work on Lab 4 and 5
  - End topics (Processes, VM) will be stressed less than usual on Final

A detailed, colorful micrograph of a microchip die, showing a complex grid of circuitry and various colored regions. The word "Processes" is overlaid in large white text on the left side of the image.

# Processes

# Lesson Summary (1/4)

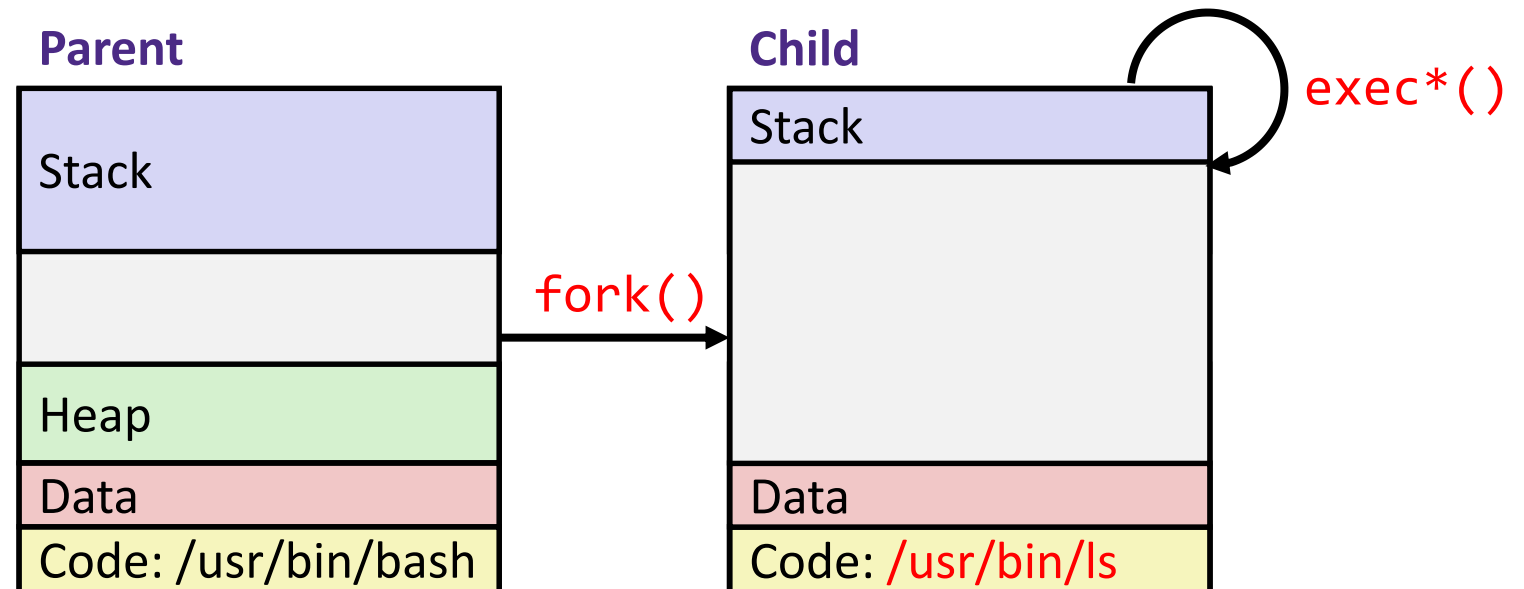
- ❖ A **process** is an instance of an running program and provides two key abstractions: logical control flow and private address space
- ❖ Multiple running processes can be run *concurrently* via **context switching**
  - Parallelism only possible with multiple CPUs/cores



# Lesson Summary (2/4)

## ❖ The *fork-exec model*

- Every process is assigned a unique **process ID** (pid)
- Every process has a parent process except for `init/system` (pid 1)
- `fork()` returns 0 to child, child's PID to parent
- `exec()` replaces the current process' code and address space with the code for a different program



# Lesson Summary (3/4)

## ❖ Terminating a process

- Return from `main()` or explicit call to `exit(status)`
- Passes a ***status code*** (`main`'s return value or `exit`'s argument) to parent process
  - 0 for normal exit, nonzero for abnormal exit

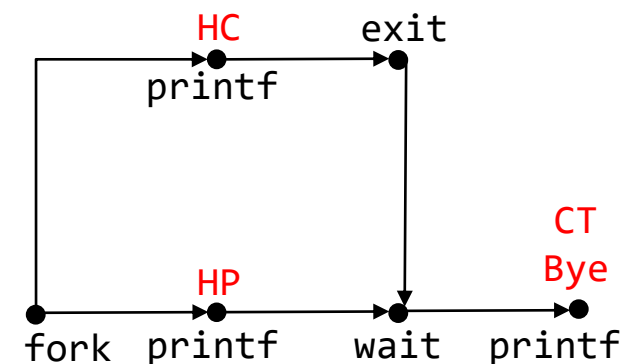
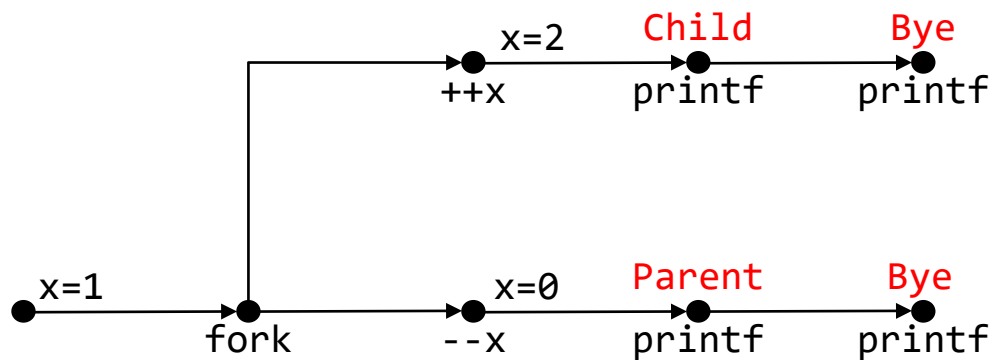
## ❖ Processes and resources

- A terminated (***zombie***) process still consumes system resources until ***reaped***
- Child is reaped when parent process terminates or explicitly calls `wait/waitpid`
- Orphaned children reaped by `init/systemd`

# Lesson Summary (4/4)

## ❖ Concurrency and *process diagrams*

- Concurrently executing processes are scheduled non-deterministically by the operating system
- A process graph is a useful tool for capturing the partial ordering of statements in a concurrent program
  - Vertices are program statements, directed edges capture sequencing *within a process*
  - Flexible visualization tool:





# Lesson Q&A

- ❖ Learning Objectives:
  - Define exceptional control flow and explain its importance in enabling concurrency and error handling.
  - Design process graphs to determine potential orderings of concurrent execution.
  
- ❖ What lingering questions do you have from the lesson?
  - Chat with your neighbors about the lesson for a few minutes to come up with questions

A detailed, colorful micrograph of a microchip die, showing a complex grid of circuitry and various colored regions (purple, blue, yellow, green, red) representing different functional blocks and interconnects.

# Processes – Practice

# Polling Questions (1/2)

❖ Are the following sequences of outputs possible?

```
void nestedfork() {
    printf("L0\n");
    if (fork() == 0) {
        printf("L1\n");
        if (fork() == 0) {
            printf("L2\n");
        }
    }
    printf("Bye\n");
}
```

**Seq 1:**

L0

L1

Bye

Bye

Bye

L2

**Seq 2:**

L0

Bye

L1

L2

Bye

Bye

- A. **No**      **No**
- B. **No**      **Yes**
- C. **Yes**     **No**
- D. **Yes**     **Yes**
- E. **We're lost...**

## Polling Questions (2/2)

- ❖ For the following scenarios, what will the outcome be for a child process that executes `exit(0)`:

Scenario	Outcome for child		
Parent is still executing:	Alive	Reaped	Zombie
Parent has called <code>wait()</code> :	Alive	Reaped	Zombie
Parent has terminated:	Alive	Reaped	Zombie

A detailed, colorful micrograph of a microchip die, showing a complex grid of circuitry and various colored regions (purple, blue, yellow, green, red) representing different functional blocks.

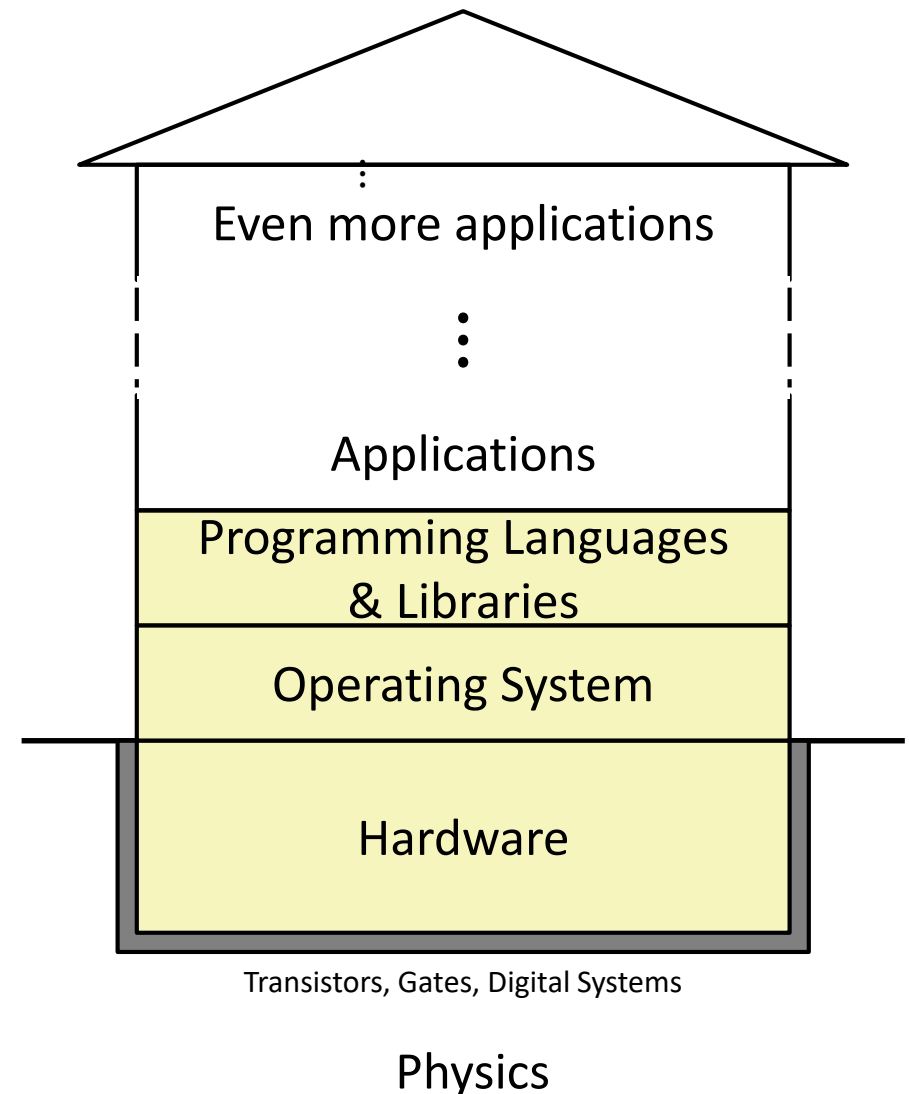
# Processes – Context

# Processes Demos

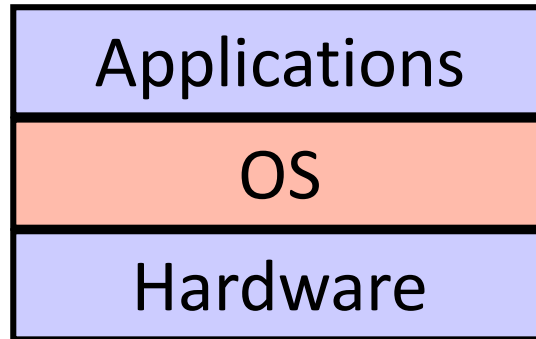
- ❖ How many processes are running on my computer right now?
- ❖ In Linux, the `ps` utility gives a snapshot of currently-running processes and `pstree` formats these as a tree
  - Can run `man ps` and `man pstree` for more info
  - Let's see a simple `pstree`
  - Let's check `attw` for some 351 zombie processes

# The Hardware/Software Interface

- ❖ Topic Group 3: **Scale & Coherence**
  - Caches, Memory Allocation, **Processes**, Virtual Memory
- ❖ How do we maintain logical consistency in the face of more data and more processes?
  - How do we support control flow both within many processes and things external to the computer?
  - How do we support data access, including dynamic requests, **across multiple processes**?



# The Operating System

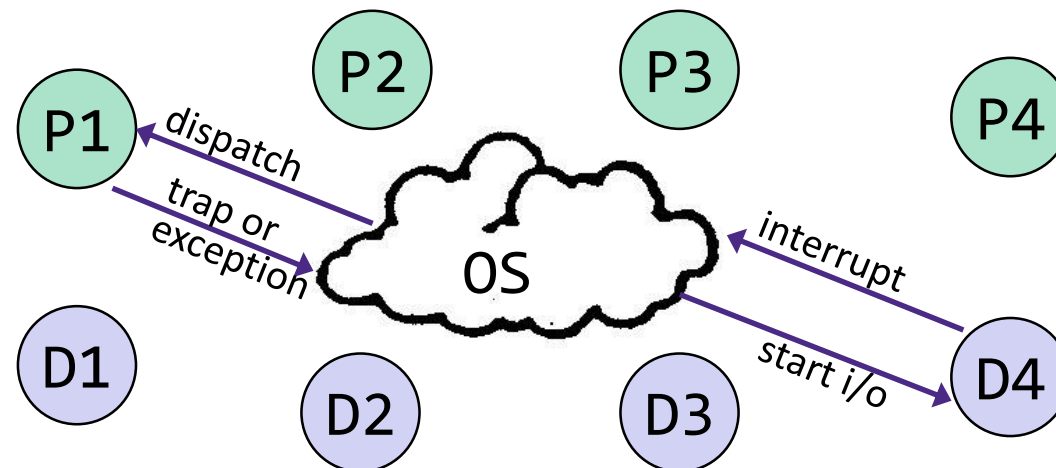


- ❖ “The OS is everything you don’t need to write in order to run your application”
- ❖ This depiction invites you to think of the OS as a library
  - In some ways, it is:
    - All operations on I/O devices require OS calls (syscalls – traps)
  - In other ways, it isn't:
    - You use the CPU/memory without OS calls
    - It intervenes without having been explicitly called



# Operating System Structure

- ❖ The OS sits between application programs (P for processes) and the hardware (D for devices)
  - It mediates access (**sharing** and **protection**)
    - Programs request services via *traps* or *exceptions*; devices request attention via *interrupts*
  - It abstracts away hardware into *logical resources* and well-defined *interfaces* to those resources (**ease of use**)
    - *e.g.*, **processes** (CPU, memory), files (disk), programs (sequences of instructions), sockets (network)



# OS Relevance in 351

- ❖ From programmer's perspective, the application benefits include:
  - Programming **simplicity**
    - Can deal with high-level abstractions instead of low-level hardware details
    - Abstractions are *reusable* across many programs
  - **Portability** (across machine configurations or architectures)
    - Device independence: 3com card or Intel card?
  
- ❖ Want to learn more?
  - CSE 333 will cover the application interface with the OS via system calls
  - CSE 451 will have you implementing the complex details of an operating system

# Group Work Time

- ❖ During this time, you are encouraged to work on the following:
  - 1) If desired, continue your discussion
  - 2) Work on the homework problems
  - 3) Work on the lab (if applicable)
  
- ❖ Resources:
  - You can revisit the lesson material
  - Work together in groups and help each other out
  - Course staff will circle around to provide support