

Overview

- ◆ Last lecture
 - Incompletely specified functions
 - Design examples
 - k-maps for POS minimization
- ◆ Today
 - Verilog
 - ☞ Structural constructs
 - ☞ Describing combinational circuits

Ways of specifying circuits

- ◆ Schematics
 - Structural description
 - Describe circuit as interconnected elements
 - ☞ Build complex circuits using hierarchy
 - ☞ Large circuits are unreadable
- ◆ HDLs
 - Hardware description languages
 - ☞ Not programming languages
 - ☞ Parallel languages tailored to digital design
 - Synthesize code to produce a circuit

Hardware description languages (HDLs)

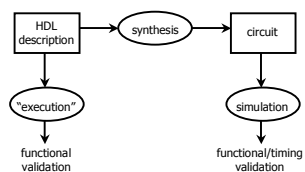
- ◆ Abel (~1983)
 - Developed by Data-I/O
 - Targeted to PLDs
 - Limited capabilities (can do state machines)
- ◆ Verilog (~1985)
 - Developed by Gateway (now part of Cadence)
 - Similar to C
 - Moved to public domain in 1990
- ◆ VHDL (~1987)
 - DoD sponsored
 - Similar to Ada

Verilog versus VHDL

- ◆ Both "IEEE standard" languages
- ◆ Most tools support both
- ◆ Verilog is "simpler"
 - Less syntax, fewer constructs
- ◆ VHDL is more structured
 - Can be better for large, complex systems
 - Better modularization

Simulation versus synthesis

- ◆ Simulation
 - "Execute" a design to verify correctness
- ◆ Synthesis
 - Generate a netlist from HDL code



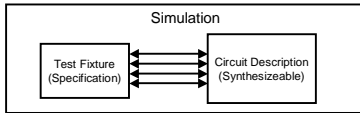
Simulation versus synthesis (con't)

- ◆ Simulation
 - Models what a circuit does
 - ☞ Multiply is "**", ignoring implementation options
 - Can include static timing
 - Allows you to test design options
- ◆ Synthesis
 - Converts your code to a netlist
 - ☞ Can simulate synthesized design
 - Tools map your netlist to hardware
- ◆ Verilog and VHDL simulate and synthesize
 - CSE370: Learn simulation
 - CSE467: Learn synthesis

Simulation

- You provide an environment
 - Using non-circuit constructs
 - Read files, print, control simulation
 - Using Verilog simulation code
 - A "test fixture"

Note: We will ignore timing and test benches until next Verilog lecture



Levels of abstraction

- Verilog supports 4 description levels
 - Switch
 - Gate
 - Dataflow
 - Algorithmic

}

}

structural
behavioral
- Can mix & match levels in a design
- Designs that combine dataflow and algorithmic constructs and synthesis are called RTL
 - Register Transfer Level

Structural versus behavioral Verilog

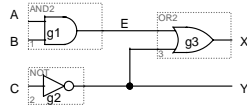
- Structural
 - Describe explicit circuit elements
 - Describe explicit connections between elements
 - Connections between logic gates
 - Just like schematics, but using text
- Behavioral
 - Describe circuit as algorithms/programs
 - What a component does
 - Input/output behavior
 - Many possible circuits could have same behavior
 - Different implementations of a Boolean function

Verilog tips

- Do not write C-code
 - Think hardware, not algorithms
 - Verilog is **inherently parallel**
 - Compilers don't map algorithms to circuits well
- Do describe hardware circuits
 - First draw a dataflow diagram
 - Then start coding
- References
 - Tutorial and reference manual are here:
 - <http://www.cs.washington.edu/education/courses/370/03sp/html/compinfo.html>

Basic building blocks: Modules

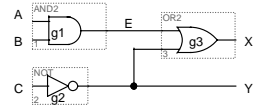
- Instantiated into a design
 - Never called
- Illegal to nest module defs.
- Modules execute in parallel
- Names are case sensitive
- // for comments
- Name can't begin with a number
- Use wires for connections
- and, or, not are keywords
- All keywords are lower case
- Gate declarations (and, or, etc)
 - List outputs first
 - Inputs second



```
// first simple example
module smpl (X,Y,A,B,C);
  input A,B,C;
  output X,Y;
  wire E;
  and g1 (E,A,B);
  not g2 (Y,C);
  or g3 (X,E,Y);
endmodule
```

Modules are circuit components

- Module has ports
 - External connections
 - A,B,C,X,Y in example
- Port types
 - input
 - output
 - inout (tristate)
- Use assign statements for Boolean expressions
 - and $\Leftrightarrow \&$
 - or $\Leftrightarrow |$
 - not $\Leftrightarrow \sim$



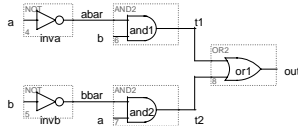
```
// previous example as a
// Boolean expression
module smpl2 (X,Y,A,B,C);
  input A,B,C;
  output X,Y;
  assign X = (A&B) | ~C;
  assign Y = ~C;
endmodule
```

Structural Verilog

```

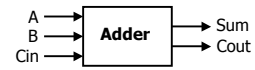
module xor_gate (out,a,b);
input  a,b;
output out;
wire  abar, bbar, t1, t2;
not    inva (abar,a);
not    invb (bbar,b);
and    and1 (t1,abar,b);
and    and2 (t2,bbar,a);
or     or1 (out,t1,t2);
endmodule
    
```

8 basic gates (keywords):
and, or, nand, nor
buf, not, xor, xnor



Behavioral Verilog

- Describe circuit behavior
 - Not implementation



```

module full_addr (Sum,Cout,A,B,Cin);
input  A, B, Cin;
output Sum, Cout;
assign {Cout, Sum} = A + B + Cin;
endmodule
    
```

{Cout, Sum} is a concatenation

Behavioral 4-bit adder

```

module add4 (SUM, OVER, A, B);
input [3:0] A;
input [3:0] B;
output [3:0] SUM;
output OVER;
assign {OVER, SUM[3:0]} = A[3:0] + B[3:0];
endmodule
    
```

"[3:0] A" is a 4-wire bus labeled "A"
Bit 3 is the MSB
Bit 0 is the LSB

Can also write "[0:3] A" Buses are implicitly connected
Bit 0 is the MSB If you write BUS[3:2], BUS[1:0]
Bit 3 is the LSB They become part of BUS[3:0]

Data types

- Values on a wire
 - 0, 1, x (don't care), z (tristate or unconnected)
- Vectors
 - A[3:0] vector of 4 bits: A[3], A[2], A[1], A[0]
 - Unsigned integer value
 - Indices must be constants
 - Concatenating bits/vectors
 - e.g. sign extend
 - $\leftarrow B[7:0] = \{A[3], A[3], A[3], A[3], A[3:0]\};$
 - $\leftarrow B[7:0] = \{4(A[3]), A[3:0]\};$
 - Style: Use $a[7:0] = b[7:0] + c;$
Not $a = b + c;$
 - Legal syntax: $C = \&A[6:7];$ // logical and of bits 6 and 7 of A

Numbers

- Format: <sign><size><base format><number>
- 14
 - Decimal number
- 4'b11
 - 4-bit 2's complement binary of 0011 (is 1101)
- 12'b0000_0100_0110
 - 12 bit binary number (_ is ignored)
- 3'h046
 - 3-digit (12-bit) hexadecimal number
- Verilog values are unsigned
 - $C[4:0] = A[3:0] + B[3:0];$
 - if $A = 0110$ (6) and $B = 1010$ (-6), then $C = 10000$ (not 00000)
 - B is zero-padded, not sign-extended

Operators

Verilog Operator	Name	Functional Group
[]	bit-select or part-select	
()	parenthesis	
!	logical negation	Logical
-	negation	Arithmetic
&	reduction AND	Reduction
	reduction OR	Reduction
&&	reduction NAND	Reduction
	reduction NOR	Reduction
^	reduction XOR	Reduction
~^ or ^~	reduction XNOR	Reduction
+	unary (sign) plus	Arithmetic
-	unary (sign) minus	Arithmetic
{}	concatenation	Concatenation
{N}	replication	Replication
*	multiply	Arithmetic
/	divide	Arithmetic
%	modulus	Arithmetic
+	binary plus	Arithmetic
-	binary minus	Arithmetic
<<	shift left	Shift
>>	shift right	Shift

>	greater than	Relational
>=	greater than or equal to	Relational
<	less than	Relational
<=	less than or equal to	Relational
==	logical equality	Equality
!=	logical inequality	Equality
===	case equality	Equality
!==	case inequality	Equality
&	bit-wise AND	Bit-wise
^	bit-wise XOR	Bit-wise
~^ or ^~	bit-wise XNOR	Bit-wise
	bit-wise OR	Bit-wise
&&	logical AND	Logical
	logical OR	Logical
?	conditional	Conditional

Similar to C operators

Continuous assignment

- ◆ Assignment is continuously evaluated
 - Corresponds to a logic gate
 - Assignments execute in parallel

```

assign A = X | (Y & ~Z);
assign B[3:0] = 4'b01XX;
assign C[15:0] = 4'h00ff;
assign #3 {Cout, Sum[3:0]} = A[3:0] + B[3:0] + Cin;
    
```

Boolean operators (~ for bit-wise negation)
 bits can assume four values (0, 1, X, Z)
 variables can be n-bits wide (MSB:LSB)
 arithmetic operator
 multiple assignment (concatenation)
 Gate delay (used by simulator)

Example: A comparator

```

module Compare1 (Equal, Alarger, Blarger, A, B);
    input A, B;
    output Equal, Alarger, Blarger;
    assign Equal = (A & B) | (~A & ~B);
    assign Alarger = (A & ~B);
    assign Blarger = (~A & B);
endmodule
    
```

Top-down design and bottom-up design are both okay
 ⇒ module ordering doesn't matter
 ⇒ because modules execute in parallel

Comparator example (con't)

```

// Make a 4-bit comparator from 4 1-bit comparators
module Compare4 (Equal, Alarger, Blarger, A4, B4);
    input [3:0] A4, B4;
    output Equal, Alarger, Blarger;
    wire e0, e1, e2, e3, A10, A11, A12, A13, B10, B11, B12, B13;

    Compare1 cp0 (e0, A10, B10, A4[0], B4[0]);
    Compare1 cp1 (e1, A11, B11, A4[1], B4[1]);
    Compare1 cp2 (e2, A12, B12, A4[2], B4[2]);
    Compare1 cp3 (e3, A13, B13, A4[3], B4[3], );

    assign Equal = (e0 & e1 & e2 & e3);
    assign Alarger = (A13 | (A12 & e3) |
                    (A11 & e3 & e2) |
                    (A10 & e3 & e2 & e1));
    assign Blarger = (~Alarger & ~Equal);
endmodule
    
```

Functions

- ◆ Use functions for complex combinational logic

```

module and_gate (out, in1, in2);
    input in1, in2;
    output out;

    assign out = myfunction(in1, in2);
    function myfunction;
        input in1, in2;
        begin
            myfunction = in1 & in2;
        end
    endfunction
endmodule
    
```

Benefit:

Functions force a result
 ⇒ Compiler will fail if function does not generate a result

Overview

- ◆ Last lecture
 - Verilog
 - Structural constructs
 - Describing combinational circuits
- ◆ Today
 - Summary of 2-level combinational logic
 - Class example: A 4-bit prime-number detector
 - Quiz

Summary of two-level combinational-logic

- ◆ Logic functions and truth tables
 - AND, OR, Buf, NOT, NAND, NOR, XOR, XNOR
 - Minimal set
- ◆ Axioms and theorems of Boolean algebra
 - Proofs by re-writing
 - Proofs by perfect induction (fill in truth table)
- ◆ Gate logic
 - Networks of Boolean functions
 - NAND/NOR conversion and de Morgan's theorem
- ◆ Canonical forms
 - Two-level forms
 - Incompletely specified functions (don't cares)
- ◆ Simplification
 - Two-level simplification (K-maps)

Solving combinational design problems

- ◆ Step 1: Understand the problem
 - Identify the inputs and outputs
 - Draw a truth table
- ◆ Step 2: Simplify the logic
 - Draw a K-map
 - Write a simplified Boolean expression
 - ▣ SOP or POS
 - ▣ Use don't cares
- ◆ Step 3: Implement the design
 - Logic gates and/or Verilog