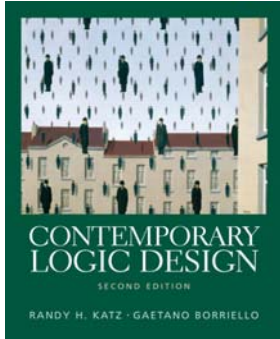


CSE 370 Spring 2006

Introduction to Digital Design

Lecture 7: Karnaugh and Beyond



Last Lecture

- Quiz
- Karnaugh Maps
- K-maps & Minimization

Today

- Design Examples & K-maps
- Minimization Algorithm

Administrivia

- Pick up Quiz 1
Average: 9.2/10, Median 10/10
- Lab 3 this week (Verilog!)
- Homework 3 on the web

Quiz Review

Problem 1: -5_{10} as a four bit expression using

- a) sign and magnitude
- b) ones-complement
- c) twos-complement

Quiz Review

$$f=AB+B'C+AC'$$

- a) Truth table
- b) Sum of products

Quiz Review

$f=AB+B'C+AC'$

b) Product of Sums

c) Circuit using AND, OR, NOT

Karnaugh Maps

■ Last Time

		A		
	0	4	12	8
	1	5	13	9
	3	7	15	11
C	2	6	14	10
		B		

		A		
	0	4	12	8
	1	5	13	9
	3	7	15	11
C	2	6	14	10
		B		

Karnaugh Map Don't Cares

■ $f(A,B,C,D) = \sum m(1,3,5,7,9) + d(6,12,13)$

■ without don't cares

■ $f = A'D + B'C'D$

		A		
	0	0	X	0
	1	1	X	1
	3	1	0	0
C	2	0	X	0
		B		

Karnaugh Map Don't Cares

■ $f(A,B,C,D) = \sum m(1,3,5,7,9) + d(6,12,13)$

■ $f = A'D + B'C'D$

without don't cares

■ $f = A'D + C'D$

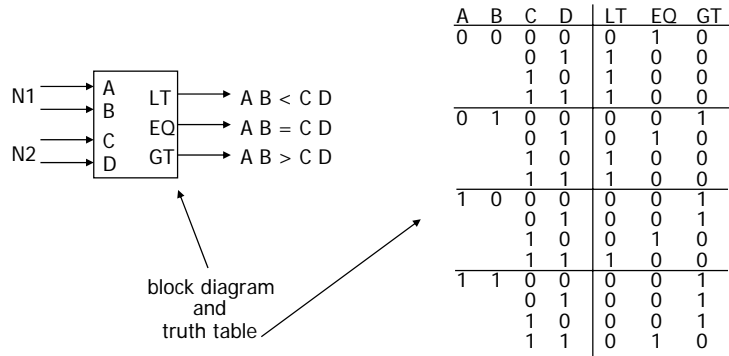
with don't cares

		A		
	0	0	X	0
	1	1	X	1
	3	1	0	0
C	2	0	X	0
		B		

by using don't care as a "1" a 2-cube can be formed rather than a 1-cube to cover this node

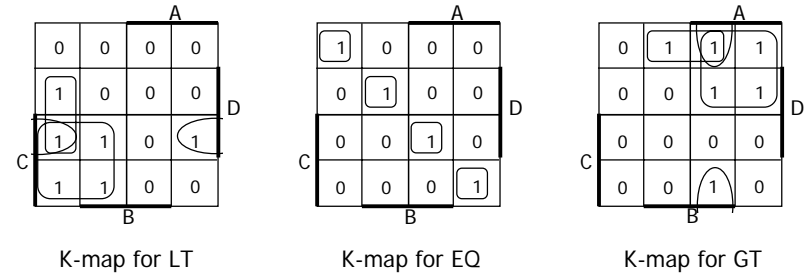
don't cares can be treated as 1s or 0s depending on which is more advantageous

Design example: two-bit comparator



we'll need a 4-variable Karnaugh map for each of the 3 output functions

Design example: two-bit comparator (cont'd)



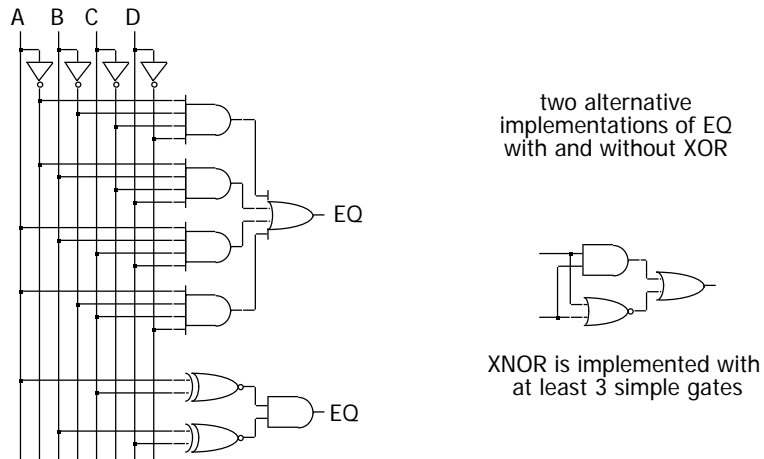
$$LT = A' B' D + A' C + B' C D$$

$$EQ = A' B' C' D' + A' B C' D + A B C D + A B' C D' = (A \text{ xnor } C) \cdot (B \text{ xnor } D)$$

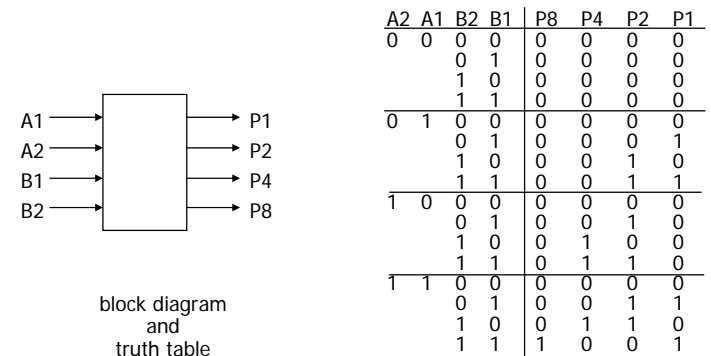
$$GT = B C' D' + A C' + A B D'$$

LT and GT are similar (flip A/C and B/D)

Design example: two-bit comparator (cont'd)

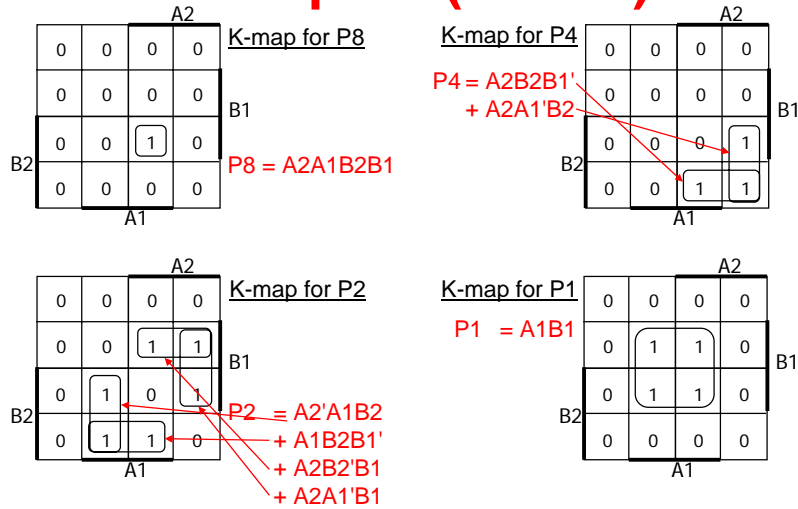


Design example: 2x2-bit multiplier

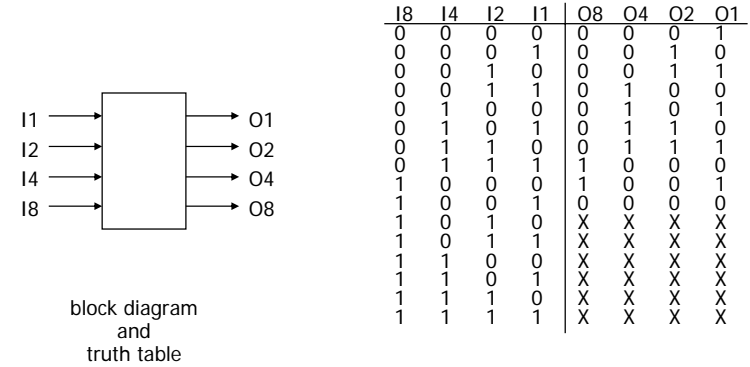


4-variable K-map for each of the 4 output functions

Design example: 2x2-bit multiplier (cont'd)

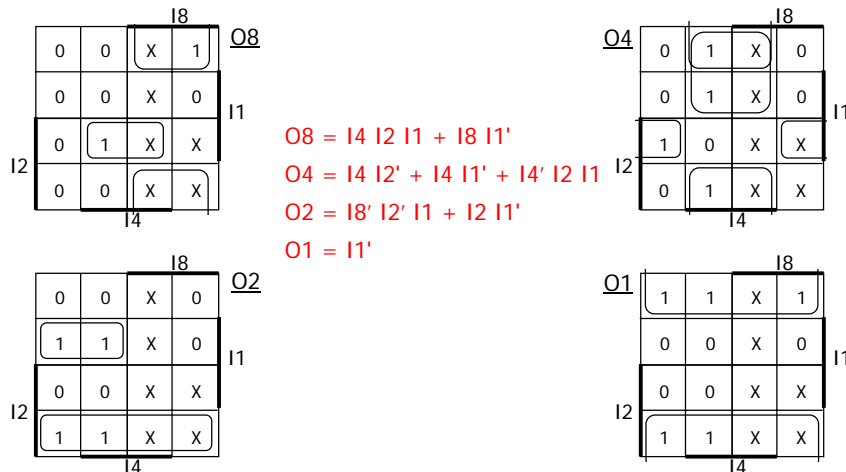


Design example: BCD increment by 1



4-variable K-map for each of the 4 output functions

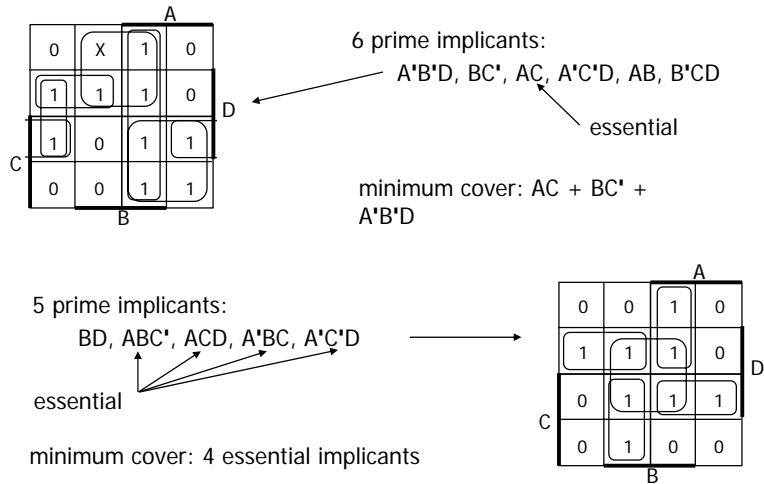
Design example: BCD increment by 1 (cont'd)



Definition of terms for two-level simplification

- **Implicant**
 - single element of ON-set or DC-set or any group of these elements that can be combined to form a subcube
- **Prime implicant**
 - implicant that can't be combined with another to form a larger subcube
- **Essential prime implicant**
 - prime implicant is essential if it alone covers an element of ON-set
 - will participate in ALL possible covers of the ON-set
 - DC-set used to form prime implicants but not to make implicant essential
- **Objective:**
 - grow implicant into prime implicants (minimize literals per term)
 - cover the ON-set with as few prime implicants as possible (minimize number of product terms)

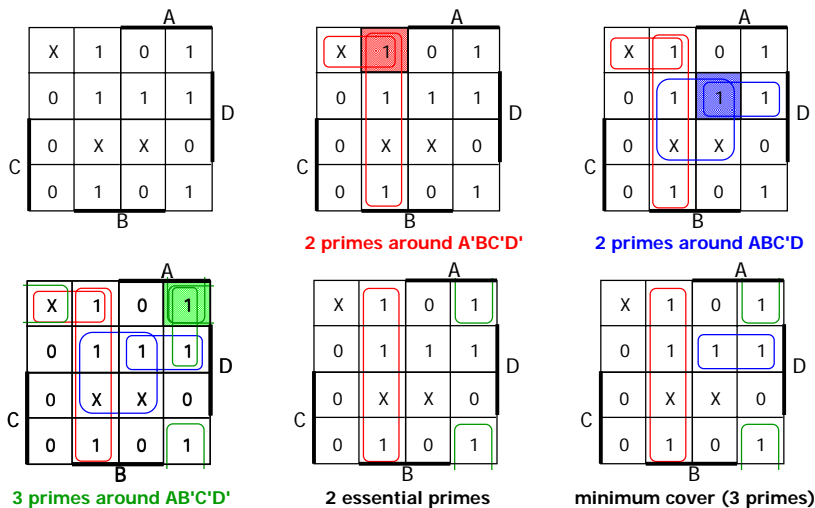
Examples to illustrate terms



Algorithm for two-level simplification

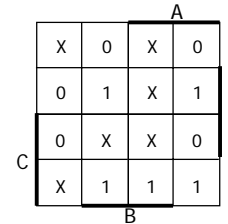
- Algorithm: minimum sum-of-products expression from a Karnaugh map
- Step 1: choose an element of the ON-set
- Step 2: find "maximal" groupings of 1s and Xs adjacent to that element
 - consider top/bottom row, left/right column, and corner adjacencies
 - this forms prime implicants (number of elements always a power of 2)
- Repeat Steps 1 and 2 to find all prime implicants
- Step 3: revisit the 1s in the K-map
 - if covered by single prime implicant, it is essential, and participates in final cover
 - 1s covered by essential prime implicant do not need to be revisited
- Step 4: if there remain 1s not covered by essential prime implicants
 - select the smallest number of prime implicants that cover the remaining 1s

Algorithm for two-level simplification (example)



Activity

- List all prime implicants for the following K-map:



- Which are essential prime implicants?
- What is the minimum cover?

Loose end: POS minimization using k-maps

- Using k-maps for POS minimization
 - Encircle the zeros in the map
 - Interpret indices complementary to SOP form

		AB		A		
		00	01	11	10	
CD	C	00	1	0	0	1
		01	0	1	0	0
	11	1	1	1	1	
	10	1	1	1	1	
		B		D		

$$F = (B'+C+D)(B+C+D')(A'+B'+C)$$

Check using de Morgan's on SOP

$$F' = BC'D' + B'C'D + ABC'$$

$$(F')' = (BC'D' + B'C'D + ABC')'$$

$$(F')' = (BC'D')' + (B'C'D)' + (ABC')'$$

$$F = (B'+C+D)(B+C+D')(A'+B'+C)$$

Implementations of two-level logic

- Sum-of-products
 - AND gates to form product terms (minterms)
 - OR gate to form sum
- Product-of-sums
 - OR gates to form sum terms (maxterms)
 - AND gates to form product

