

Overview

- ◆ Last lecture
 - Sequential Logic Examples
- ◆ Today
 - State encoding
 - ↳ One-hot encoding
 - ↳ Output encoding

One-hot encoding

- ◆ One-hot: Encode n states using n flip-flops
 - Assign a single "1" for each state
 - ↳ Example: 0001, 0010, 0100, 1000
 - Propagate a single "1" from one flip-flop to the next
 - ↳ All other flip-flop outputs are "0"
- ◆ The inverse: One-cold encoding
 - Assign a single "0" for each state
 - ↳ Example: 1110, 1101, 1011, 0111
 - Propagate a single "0" from one flip-flop to the next
 - ↳ All other flip-flop outputs are "1"
- ◆ "almost one-hot" encoding
 - Use no-hot (000...0) for the initial (reset state)
 - Assumes you never revisit the reset state

State encoding

- ◆ Assume n state bits and m states
 - $2^n / (2^n - m)!$ possible encodings [$m \geq n \geq \log_2(m)$]
 - ↳ From binomial expansion
 - ↳ Example: 3 state bits, 4 states, 1680 possible state assignments
- ◆ Hard problem, with no known algorithmic solution
 - Can try heuristic approaches
 - Can try to optimize some metric
 - ↳ FSM size (amount of logic and number of FFs)
 - ↳ FSM speed (depth of logic and fanout)
 - ↳ FSM dependencies (decomposition)
- ◆ Need to consider startup
 - Self-starting FSM or explicit reset input

One-hot encoding (con't)

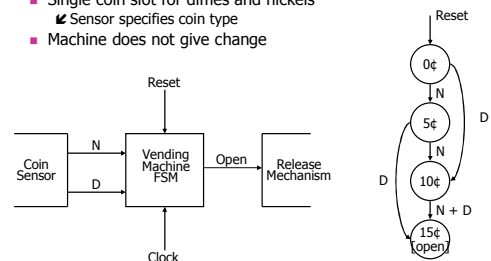
- ◆ Often the best approach for FPGAs
 - FPGAs have many flip-flops
 - One-hot machines use the least next-state logic
- ◆ Draw FSM directly from the state diagram
 - One product term per incoming arc
 - But complex state diagram \Rightarrow complex design
- ◆ One-hot designs have many possible failure modes
 - All states that aren't one-hot
 - Can create logic to reset the FSM if it enters illegal state
- ◆ Large machines require many flip-flops
 - Decompose design into smaller one-hot encoded sub-designs
 - ↳ n+m states for two machines versus n*m states for one

State-encoding strategies

- ◆ No guarantee of optimality
 - An intractable problem
- ◆ Most common strategies
 - Binary (sequential) – number states as in the state table
 - Random – computer tries random encodings
 - Heuristic – rules of thumb that seem to work well
 - ↳ e.g. Gray-code – try to give adjacent states (states with an arc between them) codes that differ in only one bit position
 - One-hot – use as many state bits as there are states
 - Output – use outputs to help encode states

Vending machine again...

- ◆ Release item after receiving 15 cents
 - Single coin slot for dimes and nickels
 - ↳ Sensor specifies coin type
 - Machine does not give change



One-hot encoded transition table

present state inputs	next state	output
$Q_3 Q_2 Q_1 Q_0$ D N	$D_3 D_2 D_1 D_0$	open
0 0 0 1	0 0	0
	0 1	0
	1 0	0
	1 1	-
0 0 1 0	0 0	0
	0 1	0
	1 0	0
	1 1	-
0 1 0 0	0 0	0
	0 1	0
	1 0	0
	1 1	-
1 0 0 0	- -	1

$$D_0 = Q_0 D' N'$$

$$D_1 = Q_0 N + Q_1 D' N'$$

$$D_2 = Q_0 D + Q_1 N + Q_2 D' N'$$

$$D_3 = Q_1 D + Q_2 D + Q_2 N + Q_3$$

$$OPEN = Q_3$$

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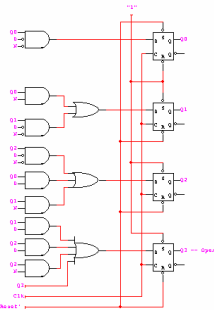
Output encoding

- ◆ Reuse outputs as state bits
 - Why create new functions when you can use outputs?
 - Bits from state assignments are the outputs for that state
 - ▣ Take outputs directly from the flip-flops
- ◆ ad hoc - no tools
 - Yields small circuits for most FSMs
 - Fits nicely with synchronous Mealy machines

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One-hot encoded vending machine



$$D_0 = Q_0 D' N'$$

$$D_1 = Q_0 N + Q_1 D' N'$$

$$D_2 = Q_0 D + Q_1 N + Q_2 D' N'$$

$$D_3 = Q_1 D + Q_2 D + Q_2 N + Q_3$$

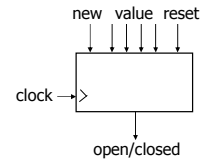
$$OPEN = Q_3$$

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Digital combination lock again...

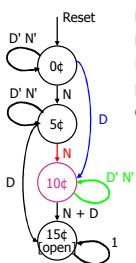
- ◆ An output-encoded FSM
 - Punch in 3 values in sequence and the door opens
 - If there is an error the lock must be reset
 - After the door opens the lock must be reset
 - Inputs: sequence of number values, reset
 - Outputs: door open/close



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Designing from the state diagram



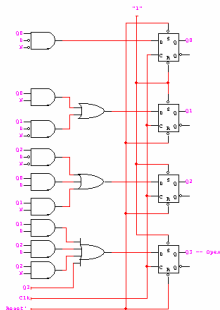
$$D_0 = Q_0 D' N'$$

$$D_1 = Q_0 N + Q_1 D' N'$$

$$D_2 = Q_0 D + Q_1 N + Q_2 D' N'$$

$$D_3 = Q_1 D + Q_2 D + Q_2 N + Q_3$$

$$OPEN = Q_3$$

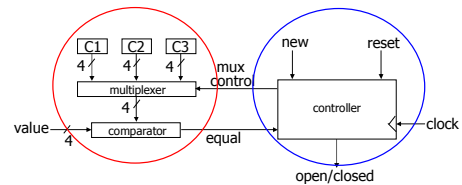


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Separate data path and control

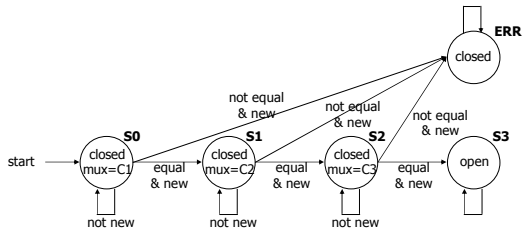
- ◆ Design datapath first
 - After the state diagram
 - Before the state encoding
- ◆ Control has 2 outputs
 - Mux control to datapath
 - Lock open/closed



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Draw the state diagram



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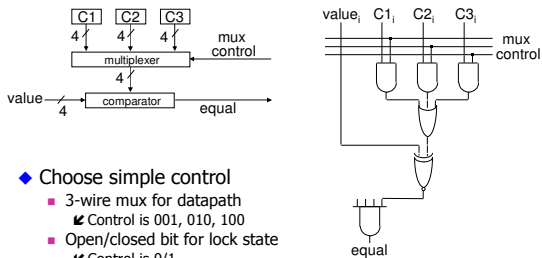
FSM has 4 state bits and 2 inputs...

- ◆ Output encoded!
 - Outputs and state bits are the same
- ◆ How do we minimize the logic?
 - FSM has 4 state bits and 2 inputs (equal, new)
 - 6-variable kmap?
- ◆ Notice the state assignment is close to one-hot
 - ERR state (0000) is only deviation
 - Is there a clever design we can use?

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Design the datapath

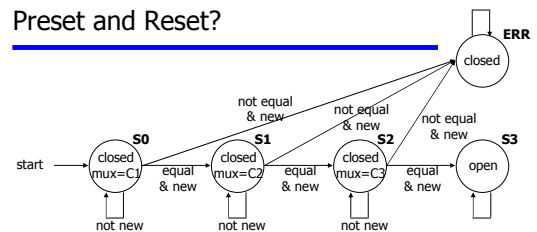


- ◆ Choose simple control
 - 3-wire mux for datapath
 - ✦ Control is 001, 010, 100
 - Open/closed bit for lock state
 - ✦ Control is 0/1

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Preset and Reset?



Assume flip-flops have preset & reset inputs
Can we encode the ERR state as reset?

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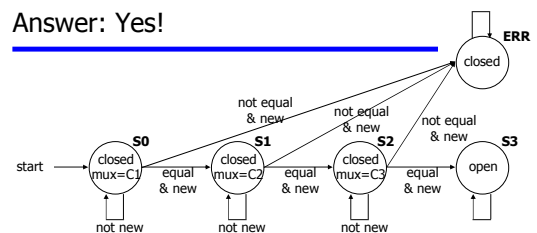
Output encode the FSM

- ◆ FSM outputs
 - Mux control is 100, 010, 001
 - Lock control is 0/1
- ◆ State are: S0, S1, S2, S3, or ERR
 - Can use 3, 4, or 5 bits to encode
 - Have 4 outputs, so choose 4 bits
 - ✦ Encode mux control and lock control in state bits
 - ✦ Lock control is first bit, mux control is last 3 bits
 - S0 = 0001 (lock closed, mux first code)
 - S1 = 0010 (lock closed, mux second code)
 - S2 = 0100 (lock closed, mux third code)
 - S3 = 1000 (lock open)
 - ERR = 0000 (error, lock closed)

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Answer: Yes!

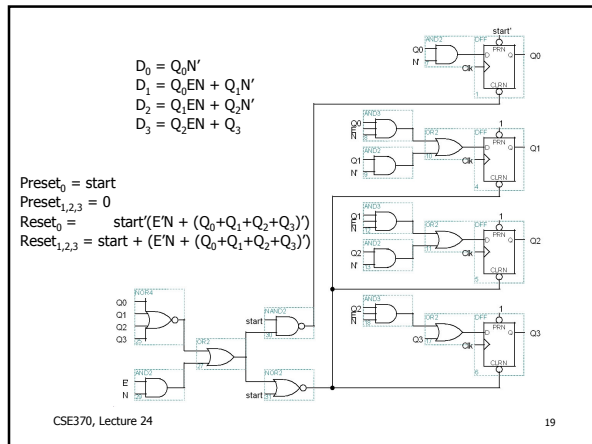


$$\begin{aligned}
 D_0 &= Q_0 N' \\
 D_1 &= Q_0 E N + Q_1 N' \\
 D_2 &= Q_1 E N + Q_2 N' \\
 D_3 &= Q_2 E N + Q_3
 \end{aligned}$$

$$\begin{aligned}
 \text{Preset}_0 &= \text{start} \\
 \text{Preset}_{1,2,3} &= 0 \\
 \text{Reset}_0 &= \text{start}'(E'N + (Q_0 + Q_1 + Q_2 + Q_3)') \\
 \text{Reset}_{1,2,3} &= \text{start} + (E'N + (Q_0 + Q_1 + Q_2 + Q_3)')
 \end{aligned}$$

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FSM design: A 5-step process

1. Understand the problem
 - State diagram and state-transition table
2. Determine the machine's states
 - Consider missing transitions: Will the machine start?
 - Minimize the state diagram: Reuse states where possible
3. Encode the states
 - Encode states, outputs with a reasonable encoding choice
 - Consider the implementation target
4. Design the next-state logic
 - Minimize the combinational logic
 - Choices made in steps 2 & 3 affect the logic complexity
5. Implement the FSM