

Lecture 4

- Logic gates and truth tables
- Implementing logic functions
- Canonical forms
 - Sum-of-products
 - Product-of-sums

Logic gates and truth tables

- AND $X \cdot Y \quad XY$

X	Y	Z
0	0	0
0	1	0
1	0	0
1	1	1
- OR $X + Y$

X	Y	Z
0	0	0
0	1	1
1	0	1
1	1	1
- NOT $X' \quad \bar{X}$

X	Y
0	1
1	0

Logic gates and truth tables

- NAND $\overline{X \cdot Y} \quad \overline{XY}$

X	Y	Z
0	0	1
0	1	1
1	0	1
1	1	0
- NOR $\overline{X + Y}$

X	Y	Z
0	0	1
0	1	0
1	0	0
1	1	0

Logic gates and truth tables

- XOR $X \oplus Y$

X	Y	Z
0	0	0
0	1	1
1	0	1
1	1	0
- XNOR $\overline{X \oplus Y}$

X	Y	Z
0	0	1
0	1	0
1	0	0
1	1	1

Realizing Boolean formulas

- $F = (A \cdot B)' + C \cdot D$
- $F = C \cdot (A + B)'$

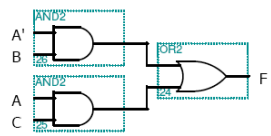
Realizing truth tables

- Given a truth table
 1. Write the Boolean expression
 2. Minimize the Boolean expression
 3. Draw as gates

Example

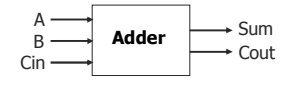
A	B	C	F
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1

$$\begin{aligned}
 F &= A'BC' + A'BC + AB'C + ABC \\
 &= A'B(C'+C) + AC(B'+B) \\
 &= A'B + AC
 \end{aligned}$$



Example: Binary full adder

- 1-bit binary adder
 - Inputs: A, B, Carry-in
 - Outputs: Sum, Carry-out



A	B	Cin	Cout	Sum
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

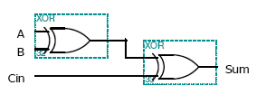
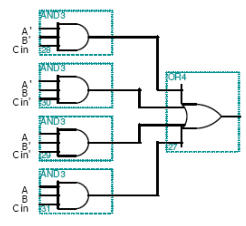
$$\begin{aligned}
 \text{Sum} &= A'B'Cin + A'BCin' + AB'Cin' + ABCin \\
 \text{Cout} &= A'BCin + AB'Cin + ABCin' + ABCin
 \end{aligned}$$

Both Sum and Cout can be minimized.

Full adder: Sum

Before Boolean minimization
 $\text{Sum} = A'B'Cin + A'BCin' + AB'Cin' + ABCin$

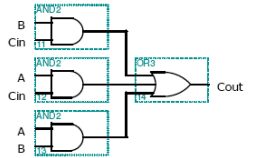
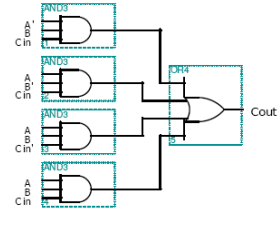
After Boolean minimization
 $\text{Sum} = (A \oplus B) \oplus Cin$



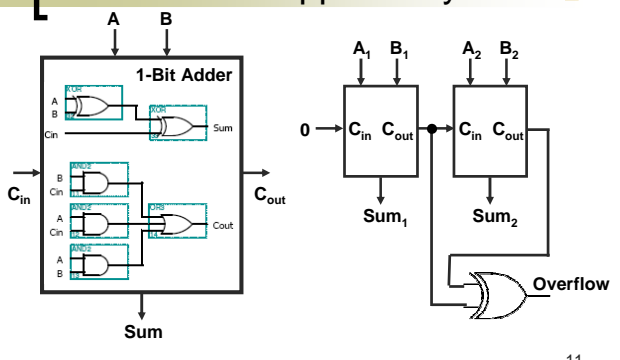
Full adder: Carry-out

Before Boolean minimization
 $\text{Cout} = A'BCin + AB'Cin + ABCin' + ABCin$

After Boolean minimization
 $\text{Cout} = BCin + ACin + AB$

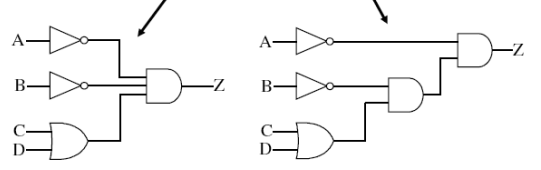


Preview: 2-bit ripple-carry adder



Many possible mappings

- Many ways to map expressions to gates
- Example: $Z = \bar{A} \cdot \bar{B} \cdot (C + D) = \bar{A} \cdot (\bar{B} \cdot (C + D))$



What is the optimal realization?

- We use the axioms and theorems of Boolean algebra to “optimize” our designs
- Design goals vary
 - Reduce the number of gates?
 - Reduce the number of gate inputs?
 - Reduce the number of cascaded levels of gates?

What is the optimal realization?

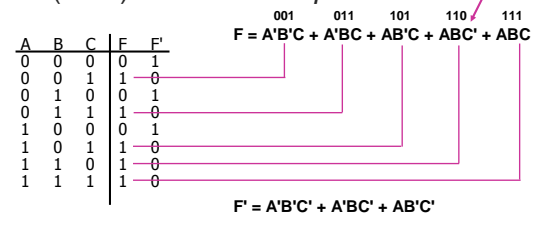
- How do we explore the tradeoffs?
 - Logic minimization: Reduce number of gates and complexity
 - Logic optimization: Maximize speed and/or minimize power
 - CAD tools

Canonical forms

- Canonical forms
 - Standard forms for Boolean expressions
 - Derived from truth table
 - Generally not the simplest forms (can be minimized)
- Two canonical forms
 - Sum-of-products (minterms)
 - Product-of-sums (maxterms)

Sum-of-products (SOP)

- Also called *disjunctive normal form* (DNF) or *minterm expansion*



Minterms

- Variables appear exactly once in each minterm in true or inverted form (but not both)

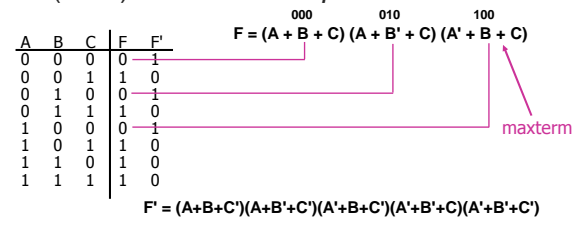
A	B	C	minterms
0	0	0	A'B'C' m0
0	0	1	A'B'C m1
0	1	0	A'BC' m2
0	1	1	A'BC m3
1	0	0	AB'C' m4
1	0	1	AB'C m5
1	1	0	ABC' m6
1	1	1	ABC m7

F in canonical form:
 $F(A,B,C) = \Sigma m(1,3,5,6,7)$
 $= m1 + m3 + m5 + m6 + m7$
 $= A'B'C + A'BC + AB'C + ABC' + ABC$

short-hand notation

Product-of-sums (POS)

- Also called *conjunctive normal form* (CNF) or *maxterm expansion*



Maxterms

- Variables appear exactly once in each maxterm in true or inverted form (but not both)

A	B	C	maxterms
0	0	0	A+B+C M0
0	0	1	A+B+C' M1
0	1	0	A+B'+C M2
0	1	1	A+B'+C' M3
1	0	0	A'+B+C M4
1	0	1	A'+B+C' M5
1	1	0	A'+B'+C M6
1	1	1	A'+B'+C' M7

F in canonical form:

$$F(A,B,C) = \prod M(0,2,4)$$

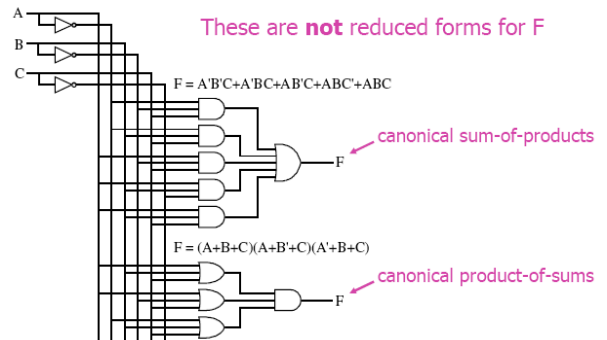
$$= M0 \cdot M2 \cdot M4$$

$$= (A+B+C)(A+B'+C)(A'+B+C)$$

short-hand notation

19

Example: F = AB+C



From SOP to POS and back

- Minterm to maxterm
 - Use maxterms that aren't in minterm expansion
 - $F(A,B,C) = \sum m(1,3,5,6,7) = \prod M(0,2,4)$
- Maxterm to minterm
 - Use minterms that aren't in maxterm expansion
 - $F(A,B,C) = \prod M(0,2,4) = \sum m(1,3,5,6,7)$

21

From SOP to POS and back

- Minterm of F to minterm of F'
 - Use minterms that don't appear
 - $F(A,B,C) = \sum m(1,3,5,6,7) \quad F' = \sum m(0,2,4)$
- Maxterm of F to maxterm of F'
 - Use maxterms that don't appear
 - $F(A,B,C) = \prod M(0,2,4) \quad F' = \prod M(1,3,5,6,7)$

22

SOP, POS, and DeMorgan's

- Sum-of-products
 - $F' = A'B'C' + A'BC' + AB'C'$
- Apply DeMorgan's to get POS
 - $(F')' = (A'B'C' + A'BC' + AB'C)'$
 - $F = (A+B+C)(A+B'+C)(A'+B+C)$

23

SOP, POS, and DeMorgan's

- Product-of-sums
 - $F' = (A+B+C)(A+B'+C)(A'+B+C)(A'+B'+C)$
- Apply DeMorgan's to get SOP
 - $(F')' = ((A+B+C)(A+B'+C)(A'+B+C)(A'+B'+C))'$
 - $F = A'B'C + A'BC + AB'C + ABC$

24