

## Lecture 2: Number Systems

- ◆ Logistics
  - <http://www.cs.washington.edu/370>
  - HW1 is posted on the web in the calendar --- due 1/14 11:30am
  - Email list: please sign up on the web.
- ◆ Last lecture
  - Class introduction and overview
- ◆ Today
  - Binary numbers
  - Base conversion
  - Number systems
    - ◇ Twos-complement
  - A/D and D/A conversion

CSE370, Lecture 2

## The "WHY" slide

- ◆ Binary numbers
  - All computers work with 0's and 1's so it is like learning alphabets before learning English
- ◆ Base conversion
  - For convenience, people use other bases (like decimal, hexadecimal) and we need to know how to convert from one to another.
- ◆ Number systems
  - There are more than one way to express a number in binary. So 1010 could be 10, -2, -5 or -6 and need to know which one.
- ◆ A/D and D/A conversion
  - Real world signals come in continuous/analog format and it is good to know generally how they become 0's and 1's (and vice versa).

CSE370, Lecture 2

## Digital

- ◆ Digital = discrete
  - Binary codes (example: BCD)
  - Decimal digits 0-9
- ◆ Binary codes
  - Represent symbols using binary digits (bits)
- ◆ Digital computers:
  - I/O is digital
    - ◇ ASCII, decimal, etc.
  - Internal representation is binary
    - ◇ Process information in bits

Decimal Symbols	BCD Code
0	0000
1	0001
2	0010
3	0011
4	0100
5	0101
6	0110
7	0111
8	1000
9	1001

CSE370, Lecture 2

## The basics: Binary numbers

- ◆ Bases we will use
  - Binary: Base 2
  - Octal: Base 8
  - Decimal: Base 10
  - Hexadecimal: Base 16
- ◆ Positional number system
  - $101_2 = 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0$
  - $63_8 = 6 \times 8^1 + 3 \times 8^0$
  - $A_{16} = 10 \times 16^1 + 1 \times 16^0$
- ◆ Addition and subtraction

$$\begin{array}{r} 1011 \\ + 1010 \\ \hline 10101 \end{array}$$

$$\begin{array}{r} 1011 \\ - 0110 \\ \hline 0101 \end{array}$$

CSE370, Lecture 2

## Binary → hex/decimal/octal conversion

- ◆ Conversion from binary to octal/hex
  - Binary: 10011110001
  - Octal: 10 | 011 | 110 | 001 =  $2361_8$
  - Hex: 100 | 1111 | 0001 =  $4F1_{16}$
- ◆ Conversion from binary to decimal
  - $101_2 = 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 = 5_{10}$
  - $63.4_8 = 6 \times 8^1 + 3 \times 8^0 + 4 \times 8^{-1} = 51.5_{10}$
  - $A_{16} = 10 \times 16^1 + 1 \times 16^0 = 161_{10}$

CSE370, Lecture 2

## Decimal → binary/octal/hex conversion

Binary			Octal		
Quotient	Remainder		Quotient	Remainder	
56 ÷ 2 =	28	0	56 ÷ 8 =	7	0
28 ÷ 2 =	14	0	7 ÷ 8 =	0	7
14 ÷ 2 =	7	0			
7 ÷ 2 =	3	1			
3 ÷ 2 =	1	1			$56_{10} = 111000_2$
1 ÷ 2 =	0	1			$56_{10} = 70_8$

- ◆ Why does this work?
  - $N = 56_{10} = 111000_2$
  - $Q = N/2 = 56/2 = 111000/2 = 11100$  remainder 0
- ◆ Each successive divide liberates an LSB (least significant bit)

CSE370, Lecture 2

## Number systems

- ◆ How do we write negative binary numbers?
- ◆ Historically: 3 approaches
  - Sign-and-magnitude
  - Ones-complement
  - Twos-complement
- ◆ For all 3, the most-significant bit (MSB) is the sign digit
  - 0 ≡ positive
  - 1 ≡ negative
- ◆ twos-complement is the important one
  - Simplifies arithmetic
  - Used almost universally

CSE370, Lecture 2

## Sign-and-magnitude

- ◆ The most-significant bit (MSB) is the sign digit
  - 0 ≡ positive
  - 1 ≡ negative
- ◆ The remaining bits are the number's magnitude
- ◆ Problem 1: Two representations for zero
  - 0 = 0000 and also -0 = 1000
- ◆ Problem 2: Arithmetic is cumbersome

Add		Subtract		Compare and subtract			
4	0100	4	0100	0100	- 4	1100	1100
+ 3	+ 0011	- 3	+ 1011	- 0011	+ 3	+ 0011	- 0011
= 7	= 0111	= 1	≠ 1111	= 0001	- 1	≠ 1111	= 1001

CSE370, Lecture 2

## Ones-complement

- ◆ Negative number: Bitwise complement positive number
  - $0111 \equiv 7_{10}$
  - $1000 \equiv -7_{10}$
- ◆ Solves the arithmetic problem

Add		Invert, add, add carry		Invert and add	
4	0100	4	0100	- 4	1011
+ 3	+ 0011	- 3	+ 1100	+ 3	+ 0011
= 7	= 0111	= 1	1 0000	- 1	1110
		add carry:	+ 1		
			= 0001		

- ◆ Remaining problem: Two representations for zero
  - 0 = 0000 and also -0 = 1111

CSE370, Lecture 2

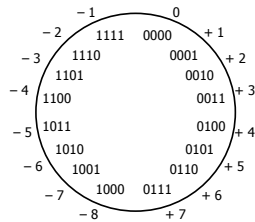
## Why ones-complement works

- ◆ The ones-complement of an 8-bit positive  $y$  is  $11111111_2 - y$
- ◆ What is  $11111111_2$ ?
  - 1 less than  $1\ 00000000_2 \equiv 2^8 \equiv 256_{10}$
  - So in ones-complement  $-y$  is represented by  $(2^8 - 1) - y$
- ◆ Adding representations of  $x$  and  $-y$  where  $x, y$  are positive we get  $(2^8 - 1) + x - y$ 
  - If  $x < y$  then  $x - y < 0$  there is no carry and get  $-ve$  number
    - ◇ Just add the representations if no carry
  - If  $x > y$  then  $x - y > 0$  there is a carry and get  $+ve$  number
    - ◇ Need to add 1 and ignore the  $2^8$  i.e. "add the carry"
  - If  $x = y$  then answer should be 0, get  $2^8 - 1 = 11111111_2$

CSE370, Lecture 2

## Twos-complement

- ◆ Negative number: Bitwise complement **plus one**
  - $0111 \equiv 7_{10}$
  - $1001 \equiv -7_{10}$
- ◆ Number wheel
- ◆ Only one zero!
- ◆ MSB is the sign digit
  - 0 ≡ positive
  - 1 ≡ negative



CSE370, Lecture 2

## Twos-complement (con't)

- ◆ Complementing a complement → the original number
- ◆ Arithmetic is easy
  - Subtraction = negation and addition
    - ◇ Easy to implement in hardware

Add		Invert and add		Invert and add	
4	0100	4	0100	- 4	1100
+ 3	+ 0011	- 3	+ 1101	+ 3	+ 0011
= 7	= 0111	= 1	1 0001	- 1	1111
		drop carry	= 0001		

CSE370, Lecture 2

## Why twos-complement works better

- ◆ Recall: The ones-complement of a b-bit positive  $y$  is  $(2^b - 1) - y$
- ◆ Adding 1 to get the twos-complement represents  $-y$  by  $2^b - y$ 
  - So  $-y$  and  $2^b - y$  are equal mod  $2^b$  (leave the same remainder when divided by  $2^b$ )
  - Ignoring carries is equivalent to doing arithmetic mod  $2^b$
- ◆ Adding representations of  $x$  and  $-y$  yields  $2^b + x - y$ 
  - If there is a carry then that means  $x \geq y$  and dropping the carry yields  $x - y$
  - If there is no carry then  $x < y$  and then we can think of it as  $2^b - (y - x)$

CSE370, Lecture 2

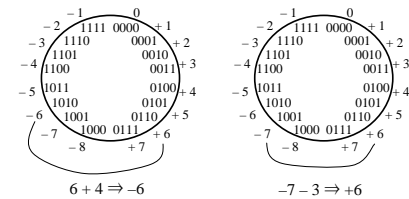
## Miscellaneous

- ◆ Twos-complement of non-integers
  - $1.6875_{10} = 01.1011_2$
  - $-1.6875_{10} = 10.0101_2$
- ◆ Sign extension
  - Write +6 and -6 as twos complement
    - ◇ 0110 and 1010
  - Sign extend to 8-bit bytes
    - ◇ 00000110 and 11111010
- ◆ Can't infer a representation from a number
  - 11001 is 25 (unsigned)
  - 11001 is -9 (sign magnitude)
  - 11001 is -6 (ones complement)
  - 11001 is -7 (twos complement)

CSE370, Lecture 2

## Twos-complement overflow

- ◆ Answers only correct mod  $2^b$ 
  - Summing two positive numbers can give negative result
  - Summing two negative numbers can give a positive result



- ◆ Make sure to have enough bits to handle overflow

CSE370, Lecture 2

## BCD (Binary-Coded Decimal) and Gray codes

Decimal Symbols	BCD Code	Decimal Symbols	Gray Code
0	0000	0	0000
1	0001	1	0001
2	0010	2	0011
3	0011	3	0010
4	0100	4	0110
5	0101	5	0111
6	0110	6	0101
7	0111	7	0100
8	1000	8	1100
9	1001	9	1101

Only one bit changes per step

CSE370, Lecture 2

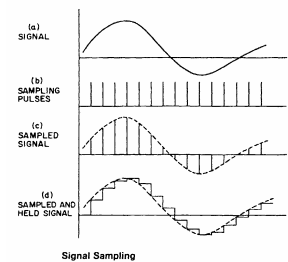
## The physical world is analog

- ◆ Digital systems need to
  - Measure analog quantities
    - ◇ Speech waveforms, etc
  - Control analog systems
    - ◇ Drive motors, etc
- ◆ How do we connect the analog and digital domains?
  - Analog-to-digital converter (A/D)
    - ◇ Example: CD recording
  - Digital-to-analog converter (D/A)
    - ◇ Example: CD playback

CSE370, Lecture 2

## Sampling

- ◆ **Quantization**
  - Conversion from analog to discrete values
- ◆ Quantizing a signal
  - We sample it



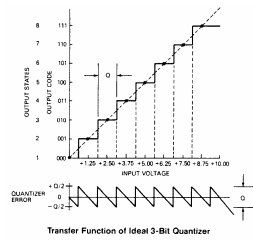
Signal Sampling

Datel Data Acquisition and Conversion Handbook

CSE370, Lecture 2

## Conversion

- ◆ **Encoding**
  - Assigning a digital word to each discrete value
- ◆ Encoding a quantized signal
  - Encode the samples
  - Typically Gray or binary codes



Transfer Function of Ideal 3-Bit Quantizer  
Dattel Data Acquisition and Conversion Handbook