

Implementation Technologies

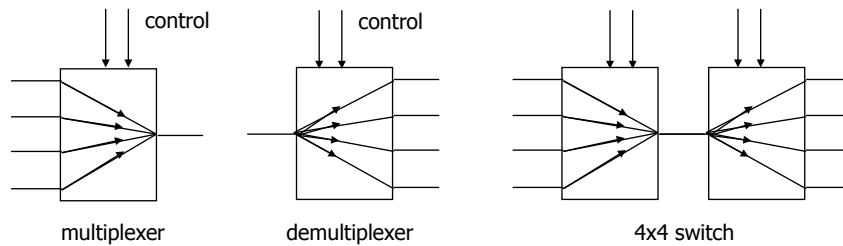
- Standard gates (pretty much done)
 - gate packages
 - cell libraries
- Regular logic (we are here)
 - multiplexers
 - decoders
- Two-level programmable logic (a little later)
 - PALs, PLAs, PLDs
 - ROMs
 - FPGAs

Regular logic

- Need to make design faster
- Need to make engineering changes easier to make
- Simpler for designers to understand and map to functionality
 - harder to think in terms of specific gates
 - easier to think in terms of larger multi-purpose blocks

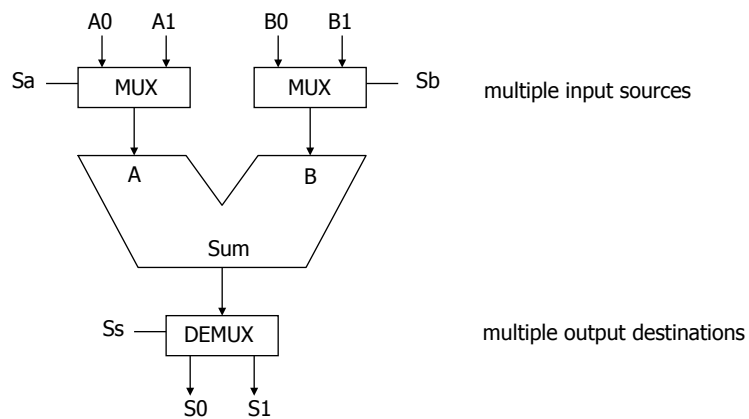
Making connections

- Direct point-to-point connections using wires
- Route one of many inputs to a single output --- multiplexer
- Route a single input to one of many outputs --- demultiplexer



Mux and demux (cont'd)

- Uses of multiplexers/demultiplexers in multi-point connections



Multiplexers/selectors

- Multiplexers/selectors: general concept
 - 2^n data inputs, n control inputs (called "selects"), 1 output
 - used to connect 2^n points to a single point
 - control signal pattern forms binary index of input connected to output

$Z = A' I_0 + A I_1$

A	Z
0	I_0
1	I_1

functional form

logical form

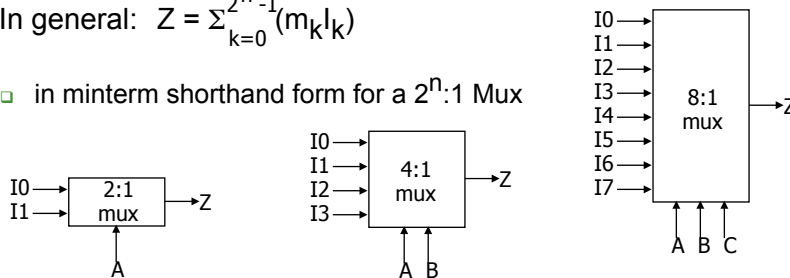
I_1	I_0	A	Z
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

two alternative forms for a 2:1 Mux truth table

Multiplexers/selectors (cont'd)

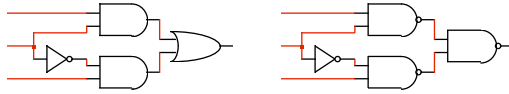
- 2:1 mux: $Z = A'I_0 + AI_1$
- 4:1 mux: $Z = A'B'I_0 + A'BI_1 + AB'I_2 + ABI_3$
- 8:1 mux: $Z = A'B'C'I_0 + A'B'CI_1 + A'BC'I_2 + A'BCI_3 + AB'C'I_4 + AB'CI_5 + ABC'I_6 + ABCI_7$
- In general: $Z = \sum_{k=0}^{2^n-1} (m_k I_k)$

- in minterm shorthand form for a $2^n:1$ Mux

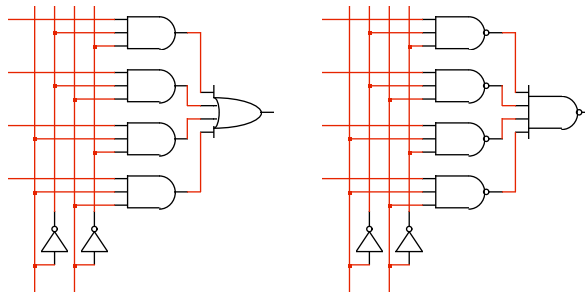


Gate level implementation of muxes

- 2:1 mux



- 4:1 mux



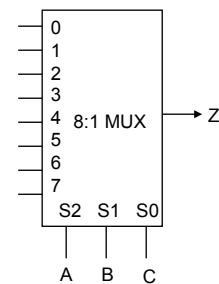
Multiplexers as general-purpose logic

- A $2^n:1$ multiplexer can implement any function of n variables
 - with the variables used as control inputs and
 - the data inputs tied to 0 or 1
 - in essence, a **lookup table (LUT)**

- Example:

$$F(A,B,C) = m_0 + m_2 + m_6 + m_7$$

$$= A'B'C' + A'BC' + ABC' + ABC$$

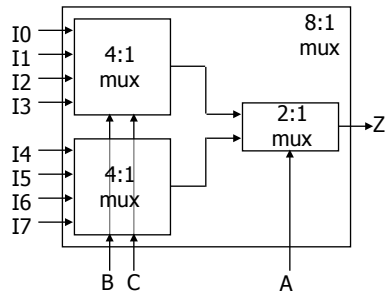


$$Z = A'B'C'I_0 + A'B'CI_1 + A'BC'I_2 + A'BCI_3 +$$

$$AB'C'I_4 + AB'CI_5 + ABC'I_6 + ABCI_7$$

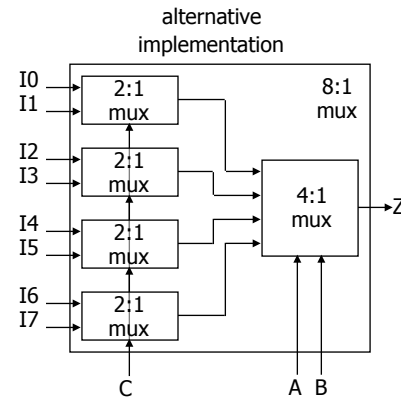
Cascading multiplexers

- Large multiplexers can be made by cascading smaller ones



control signals B and C simultaneously choose one of I0, I1, I2, I3 and one of I4, I5, I6, I7

control signal A chooses which of the upper or lower mux's output to gate to Z

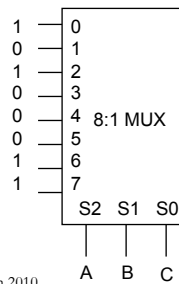


Multiplexers as general-purpose logic (cont'd)

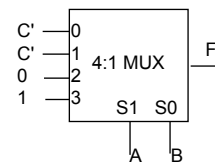
- A $2^{n-1}:1$ multiplexer can implement any function of n variables
 - with $n-1$ variables used as control inputs and
 - the data inputs tied to the last variable or its complement

- Example:

$$\begin{aligned}
 F(A,B,C) &= m_0 + m_2 + m_6 + m_7 \\
 &= A'B'C' + A'BC' + ABC' + ABC \\
 &= A'B'(C') + A'B(C') + AB'(0) + AB(1)
 \end{aligned}$$



A	B	C	F
0	0	0	1
0	0	1	0
0	1	0	1
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1



Multiplexers as general-purpose logic (cont'd)

- Generalization

n-1 mux control variables

single mux data variable

I_0	I_1	...	I_{n-1}	I_n
.	.	.	.	0
.	.	.	.	1

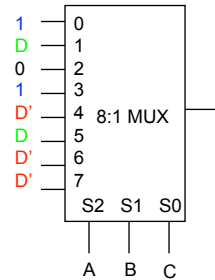
	F			
0	0	1	1	1
0	1	0	1	1
↓	↓	↓	↓	↓
0	I_n	I_n'	1	

four possible configurations of truth table rows can be expressed as a function of I_n

- Example:
 $G(A,B,C,D)$
can be realized by an 8:1 MUX

choose A,B,C as control variables

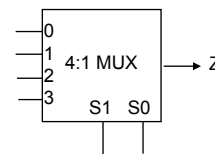
A	B	C	D	G
0	0	0	0	1
0	0	0	1	1
0	0	1	0	0
0	0	1	1	1
0	1	0	0	0
0	1	0	1	0
0	1	1	0	1
0	1	1	1	1
1	0	0	0	1
1	0	0	1	0
1	0	1	0	0
1	0	1	1	1
1	1	0	0	1
1	1	0	1	0
1	1	1	0	1
1	1	1	1	0



Activity

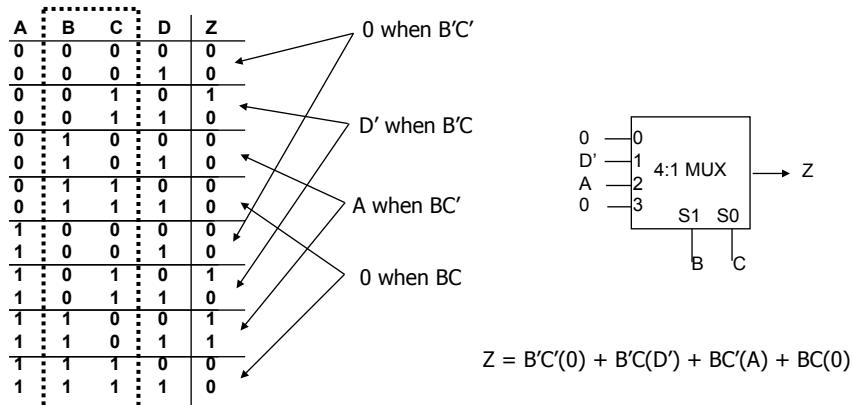
- Realize $F = B'CD' + ABC'$ with a 4:1 multiplexer and a minimum of other gates:

A	B	C	D	Z
0	0	0	0	0
0	0	0	1	0
0	0	1	0	1
0	0	1	1	0
0	1	0	0	0
0	1	0	1	0
0	1	1	0	0
0	1	1	1	0
1	0	0	0	0
1	0	0	1	0
1	0	1	0	1
1	0	1	1	0
1	1	0	0	1
1	1	0	1	1
1	1	1	0	0
1	1	1	1	0



Activity

- Realize $F = B'CD' + ABC'$ with a 4:1 multiplexer and a minimum of other gates:



Demultiplexers/decoders

- Decoders/demultiplexers: general concept
 - single data input, n control inputs, 2^n outputs
 - control inputs (called "selects" (S)) represent binary index of output to which the input is connected
 - data input usually called "enable" (G)

1:2 Decoder:

$$\begin{aligned} O0 &= G \cdot S' \\ O1 &= G \cdot S \end{aligned}$$

2:4 Decoder:

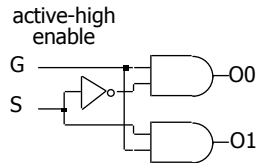
$$\begin{aligned} O0 &= G \cdot S1' \cdot S0' \\ O1 &= G \cdot S1' \cdot S0 \\ O2 &= G \cdot S1 \cdot S0' \\ O3 &= G \cdot S1 \cdot S0 \end{aligned}$$

3:8 Decoder:

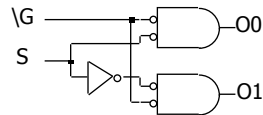
$$\begin{aligned} O0 &= G \cdot S2' \cdot S1' \cdot S0' \\ O1 &= G \cdot S2' \cdot S1' \cdot S0 \\ O2 &= G \cdot S2' \cdot S1 \cdot S0' \\ O3 &= G \cdot S2' \cdot S1 \cdot S0 \\ O4 &= G \cdot S2 \cdot S1' \cdot S0' \\ O5 &= G \cdot S2 \cdot S1' \cdot S0 \\ O6 &= G \cdot S2 \cdot S1 \cdot S0' \\ O7 &= G \cdot S2 \cdot S1 \cdot S0 \end{aligned}$$

Gate level implementation of demultiplexers

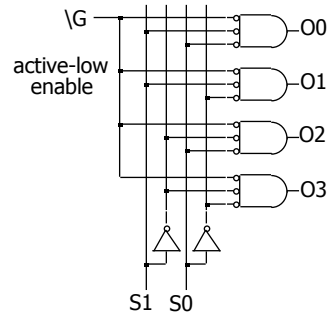
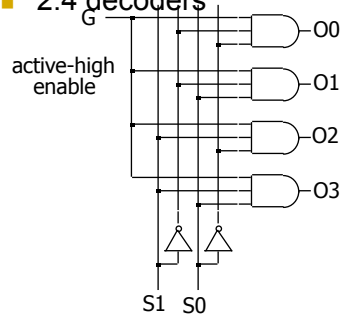
1:2 decoders



active-low enable



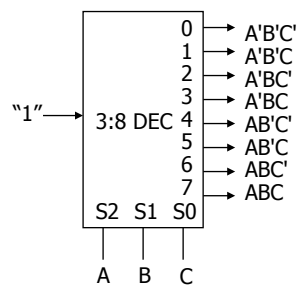
2:4 decoders



Demultiplexers as general-purpose logic

A $n:2^n$ decoder can implement any function of n variables

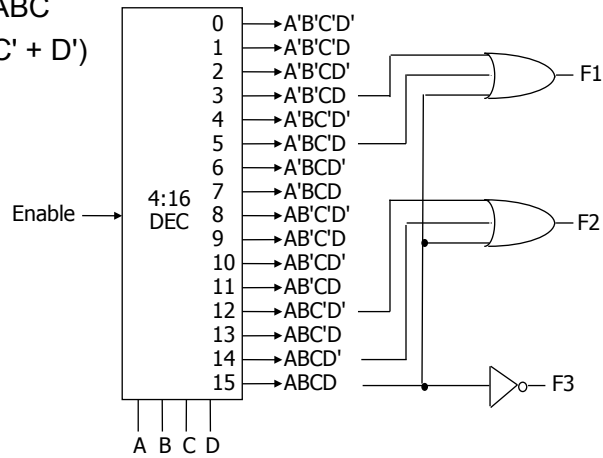
- with the variables used as control inputs
- the enable inputs tied to 1 and
- the appropriate minterms summed to form the function



demultiplexer generates appropriate minterm based on control signals (it "decodes" control signals)

Demultiplexers as general-purpose logic (cont'd)

- $F1 = A'BC'D + A'B'CD + ABCD$
- $F2 = ABC'D' + ABC$
- $F3 = (A' + B' + C' + D')$



Cascading decoders

- 5:32 decoder
 - 1x2:4 decoder
 - 4x3:8 decoders

