Graph Matching

CSE 373

Data Structures

DS.GR.14

Graph Matching

Input: 2 digraphs G1 = (V1,E1), G2 = (V2,E2)

Questions to ask:

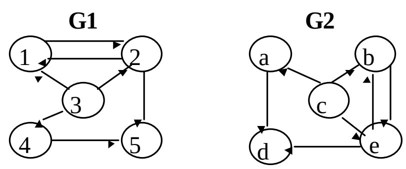
1. Are G1 and G2 isomorphic?
2. Is G1 isomorphic to a subgraph of G2?
3. How similar is G1 to G2?
4. How similar is G1 to the most similar subgraph of G2?

DS.GR.15

Isomorphism for Digraphs

G1 is isomorphic to G2 if there is a 1-1, onto mapping h: V1 → V2 such that

(vi,vj) ∈ E1 iff (h(vi), h(vj)) ∈ E2



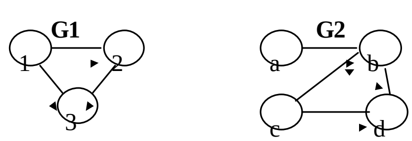
Find an isomorphism h: {1,2,3,4,5} → {a,b,c,d,e}.
Check that the condition holds for every edge.

DS.GR.16

Subgraph Isomorphism for Digraphs

G1 is isomorphic to a subgraph of G2 if there is a 1-1 mapping h: V1 → V2 such that

(vi,vj) ∈ E1⇒ (h(vi), h(vj)) ∈ E2



Isomorphism and subgraph isomorphism are defined similarly for undirected graphs.

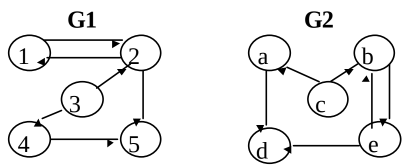
In this case, when (vi,vj) ∈ E1, either (vi,vj) or (vj,vi) can be listed in E2, since they are equivalent and both mean {vi,vj}.

DS.GR.17

Similar Digraphs

Sometimes two graphs are close to isomorphic, but have a few "errors."

Let h(1)=b, h(2)=e, h(3)=c, h(4)=a, h(5) = d.



(1,2)	(b,e)
(2,1)	(e,b)
X	(c,b)
(4,5)	(a,d)
(2,5)	(e,d)
(3,2)	X
(3,4)	(c,a)

The mapping **h** has **2 errors** .

(c,b) ∈ G2, but (3,1) ∉ G1

(3,2) ∈ G1, but (c,e) ∉ G2

DS.GR.18

Error of a Mapping

Intuitively, the error of mapping h tells us
- how many edges of G1 have no corresponding edge in G2 and
- how many edges of G2 have no corresponding edge in G1.

Let G1=(V1,E1) and G2=(V2,E2), and let h:V1→ V2 be a 1-1, onto mapping.

forward error $EF(h) = |\{(vi,vj) \in E1 \mid (h(vi),h(vj)) \notin E2\}|$
edge in E1 corresponding edge not in E2

backward error $EB(h) = |\{(vi,vj) \in E2 \mid (h^{-1}(vi),h^{-1}(vj)) \notin E1\}|$
edge in E2 corresponding edge not in E1

total error $Error(h) = EF(h) + EB(h)$

relational distance $GD(G1,G2) = \min Error(h)$
for all 1-1, onto h:V1→ V2

DS.GR.19

Variations of Relational Distance

1. normalized relational distance:
   Divide by the sum of the number of edges in E1 and those in E2.

2. undirected graphs:
   Just modify the definitions of EF and EB to accommodate.

3. one way mappings:
   h is 1-1, but need not be onto
   Only the forward error EF is used.

4. labeled graphs:
   When nodes and edges can have labels, each node should be mapped to a node with the same label, and each edge should be mapped to an edge with the same label.

DS.GR.20

Graph Matching Algorithms

1. graph isomorphism
* 2. subgraph isomorphism
* 3. relational distance
4. attributed relational distance (uses labels)

Subgraph Isomorphism

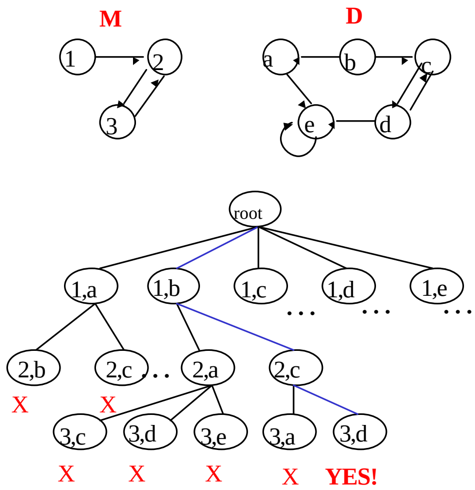
Given model graph M = (VM,EM)
data graph D = (VD,ED)

Find 1-1 mapping $h: VM \rightarrow VD$

satisfying $(vi,vj) \in EM \Rightarrow ((h(vi),h(vj)) \in ED$.

DS.GR.21

Method: Backtracking Tree Search



DS.GR.22

Treesearch for Subgraph Isomorphism in Digraphs

```
procedure Treesearch(VM, VD, EM, ED, h)
  {
  v = first(VM);
  for each w ∈ VD
    {
    h' = h ∪ {(v,w)};   //add to set
    OK = true;
    for each edge (vi,vj) in EM   (with vi < vj for undirected graphs)
      if one of vi or vj is v and the other
           has been assigned a value in h'
            if ( (h'(vi),h'(vj)) is NOT in ED )
              {OK = false; break;};

    if OK  {
            VM' = VM – v;   //remove from set
            VD'  = VD – w'
             if isempty(VM') output(h');
             else Treesearch(VM',VD',EM,ED,h')
             } } }
```

Branch-and-Bound Tree Search

Keep track of the least-error mapping.