

Lists

CSE 373
Data Structures

Readings

- Reading
 - › Goodrich and Tamassia, Chapter 5

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List ADT

- What is a List?
 - › Ordered sequence of elements A_1, A_2, \dots, A_N
- Elements may be of arbitrary type, but all are of the same type
- Common List operations are:
 - › Insert, Find, Delete, IsEmpty, IsLast, FindPrevious, First, Kth, Last, Print, etc.

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Simple Examples of List Use

- Polynomials
 - › $25 + 4x^2 + 75x^{85}$
- Unbounded Integers
 - › 4576809099383658390187457649494578
- Text
 - › "This is an example of text"

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List Implementations

- Two types of implementation:
 - › Array-Based
 - › Pointer-Based

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List: Array Implementation

- Basic Idea:
 - › Pre-allocate a big array of size MAX_SIZE
 - › Keep track of current size using a variable `count`
 - › Shift elements when you have to insert or delete

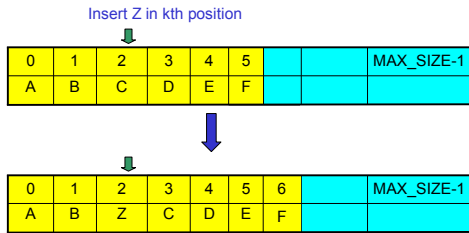
0	1	2	3	...	count-1	MAX_SIZE-1
A_1	A_2	A_3	A_4	...	A_N	

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List: Array Implementation



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Array List Insert Running Time

- Running time for N elements?
- On average, must move half the elements to make room – assuming insertions at positions are equally likely
- Worst case is insert at position 0. Must move all N items one position before the insert
- This is $O(N)$ running time. Probably too slow

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Review Big Oh Notation

- $T(N) = O(f(N))$ if there are positive constants c and n_0 such that:

$$T(N) \leq c f(N) \text{ when } N \geq n_0$$
- $T(N) = O(N)$ linear

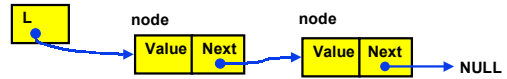
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List: Pointer Implementation

- Basic Idea:
 - › Allocate little blocks of memory (nodes) as elements are added to the list
 - › Keep track of list by linking the nodes together
 - › Change links when you want to insert or delete

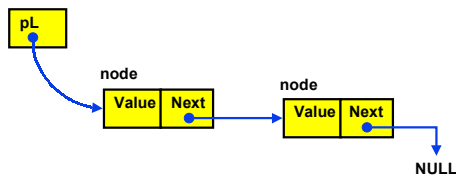


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Pointer-Based Linked List

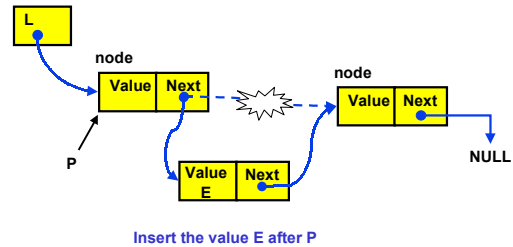


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Pointer-based Insert (after p)



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Insertion After

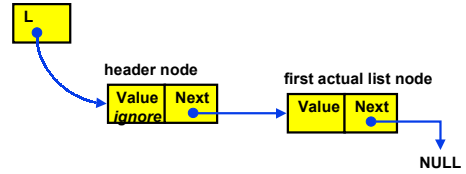
```
InsertAfter(p : node pointer, v : thing): {  
  x : node pointer;  
  x := new node;  
  x.value := v;  
  x.next := p.next;  
  p.next := x;  
}
```

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Linked List with Header Node



Advantage: "insert after" and "delete after" can be done at the beginning of the list.

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Pointer Implementation Issues

- Whenever you break a list, your code should fix the list up as soon as possible
 - › Draw pictures of the list to visualize what needs to be done
- Pay special attention to boundary conditions:
 - › Empty list
 - › Single item – same item is both first and last
 - › Two items – first, last, but no middle items
 - › Three or more items – first, last, and middle items

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Pointer List Insert Running Time

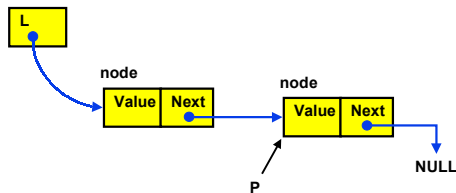
- Running time for N elements?
- Insert takes constant time ($O(1)$)
- Does not depend on input size
- Compare to array based list which is $O(N)$

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Linked List Delete



To delete the node pointed to by P, need a pointer to the previous node; See book for findPrevious method

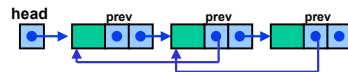
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Doubly Linked Lists

- findPrevious (and hence Delete) is slow [$O(N)$] because we cannot go directly to previous node
- Solution: Keep a "previous" pointer at each node



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Double Link Pros and Cons

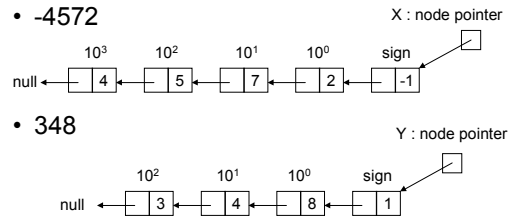
- Advantage
 - › Delete (not DeleteAfter) and FindPrev are faster
- Disadvantages:
 - › More space used up (double the number of pointers at each node)
 - › More book-keeping for updating the two pointers at each node (pretty negligible overhead)

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Unbounded Integers Base 10

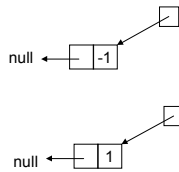


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Zero



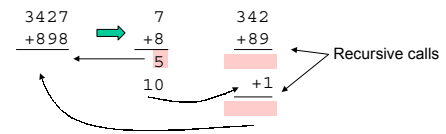
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Recursive Addition

- Positive numbers (or negative numbers)



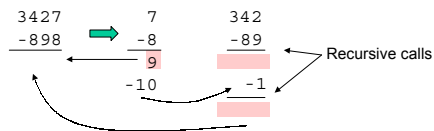
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Recursive Addition

- Mixed numbers



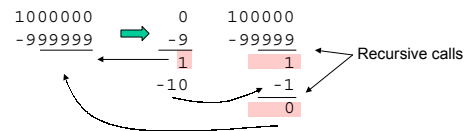
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Example

- Mixed numbers



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