

CSE 373, Autumn 2008, Assignment 2 Solutions

October 18, 2008

1. (8 points)

(a) $n = ((795 - (-10))/7) + 1 = 116$
Sum = $(((-10) + 795)116)/2 = 45530$

(b) Sum = $(256 * (1 - (1/2)^9))/(1 - (1/2)) = 511$

(c) Sum = $(1 * (3^9 - 1))/(3 - 1) = 9841$

(d) Sum = $144/(1 - (1/4)) = 192$

2. (6 points)

(a) 10^{x+y+z}

(b) xy

(c) $1 + 2 \log_2 x + 3 \log_2 y$

3. (7 points)

Basis Step:

$$n = 1, (1 + 1) = 1 * (1 + 3)/2 = 2.$$

Induction hypothesis:

$$\sum_{i=1}^k (i + 1) = \frac{k(k+3)}{2}, \text{ for some } k.$$

Induction step:

$$\sum_{i=1}^{k+1} (i + 1) = \sum_{i=1}^k (i + 1) + ((k + 1) + 1) = \frac{k(k+3)}{2} + ((k + 1) + 1) = \frac{(k+1)((k+1)+3)}{2}$$

This represents the proposition to be proved for the case $n = k + 1$, and completes the proof.

4. (6 points)

(a) $\{\}, \{0\}, \{1\}, \{0, 1\}$

(b) $(0, 0), (0, 1), (1, 0), (1, 1)$

(c) $(0, 0, 0), (0, 0, 1), (0, 1, 0), (0, 1, 1), (1, 0, 0), (1, 0, 1), (1, 1, 0), (1, 1, 1)$

5. (18 points)

	<i>R1</i>	<i>R2</i>	<i>R3</i>	<i>R4</i>	<i>R5</i>	<i>R6</i>
Reflexive	<i>N</i>	<i>N</i>	<i>N</i>	<i>Y</i>	<i>Y</i>	<i>Y</i>
Symmetric	<i>Y</i>	<i>Y</i>	<i>Y</i>	<i>Y</i>	<i>N</i>	<i>Y</i>
Transitive	<i>Y</i>	<i>Y</i>	<i>N</i>	<i>Y</i>	<i>Y</i>	<i>Y</i>
Antisymmetric	<i>Y</i>	<i>Y</i>	<i>N</i>	<i>Y</i>	<i>Y</i>	<i>N</i>
Equivalence Relation	<i>N</i>	<i>N</i>	<i>N</i>	<i>Y</i>	<i>N</i>	<i>Y</i>
Partial Order	<i>N</i>	<i>N</i>	<i>N</i>	<i>Y</i>	<i>Y</i>	<i>N</i>

6. (20 points, 15 for table entries and 5 for explanations)

	100	$2n + 5$	$\log_2 n$	$5n^2$	$n \log_2 n$
$3n + 1$	Ω	Θ	Ω	O	O
$0.001 * 2^{n-10}$	Ω	Ω	Ω	Ω	Ω
$\log_{10} n^n$	Ω	Ω	Ω	O	Θ

$0.001 * 2^{n-10} \geq 5n^2$ for $n \geq 33$ as can be verified by taking base 2 logs on both sides.

$$\log_{10} n^n = n \log_{10} n = n \log_2 n / \log_2 10 = \Theta(n \log_2 n)$$

7. (20 points)

(a) (12 points) We will use stack *Sa* for enqueueing, *Sb* for dequeueing, and a boolean variable *enQmode* for storing the current operating mode. The methods are shown below.

```

boolean isEmpty(){
    if(Sa.isEmpty() && Sb.isEmpty())
        return true;
    else
        return false; }

void enqueue(Object obj){
    if(!enQmode){
        while(!Sb.isEmpty())
            Sa.push(Sb.pop()); }
    Sa.push(obj); }

Object dequeue(){
    if(enQmode){
        while(!Sa.isEmpty())
            Sb.push(Sa.pop()); }
    return Sb.pop(); }

```

- (b) (4 points) The isEmpty method is constant time. The enqueue and dequeue operations take $O(m)$ time in the worst case where m is the current size of the queue. This is because we may need to move all m objects from one stack to another.
- (c) (4 points) The total time complexity is $O(n^2)$. There are $2n$ operations each of which takes $O(n)$ time.

8. (15 points)

- (a) (5 points) The algorithm goes thru the following steps.
 - poly = 4
 - poly = $4 * 3 + 8 = 20$
 - poly = $20 * 3 + 0 = 60$
 - poly = $60 * 3 + 1 = 181$
 - poly = $181 * 3 + 2 = 545$
- (b) (5 points) Observe that the polynomial $a_0 + a_1x + a_2x^2 + \dots$ can be rewritten as $a_0 + x(a_1 + x(a_2 + \dots))$ by repeatedly factoring out x . The algorithm computes the polynomial using this equivalent form.
- (c) (5 points) The running time is $\Theta(n)$.