

# CSE 373

# Data Structures & Algorithms

## Lecture 02

### Lists, Stacks, and Queues

# Announcements

Guest Lecturers next week:

- Monday: Vibhor Rastogi
- Wednesday: Sean Shih-Yen Liu

(I will be giving a keynote lecture at a conference in Brazil, about my research on probabilistic databases)

# Today's Outline

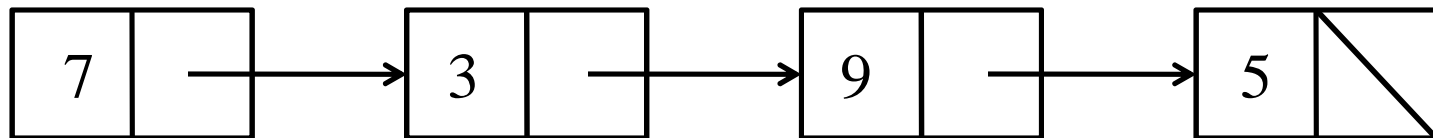
- Lists
- Stacks and queues
- Begin trees

Remember:

- Reading assignment for Monday: Ch. 1- 3
- For today: 3.1, 3.2, 3.3 (did you read them ?)

# The List ADT

- List of items, e.g. 7, 3, 9, 5
- Insert, delete, find, replace, ...
- Implementations: linked list or array list



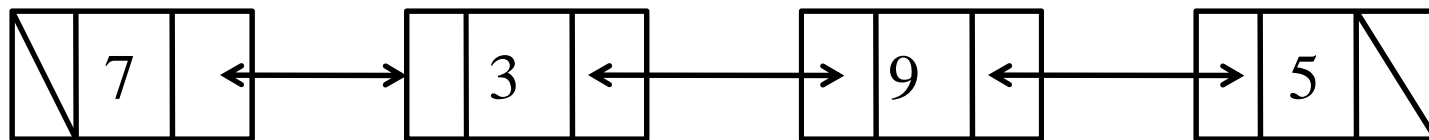
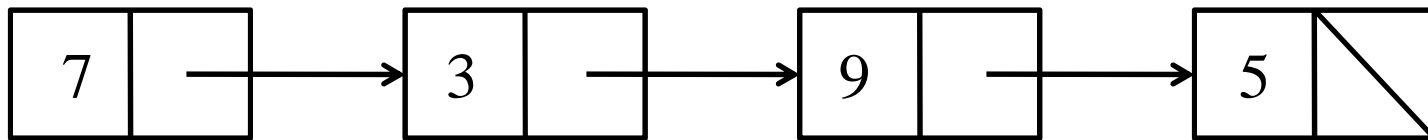
# Linked List Implementation

First, need to define a single node:

```
private static class Node<AnyType>
{
    public Node( AnyType d, Node<AnyType> n )
    {
        data = d;  next = n;
    }

    public AnyType data;
    public Node<AnyType>  next;
}
```

# Simple Linked v.s. Double Linked



What are the tradeoffs ?

# Double Linked List Implementation

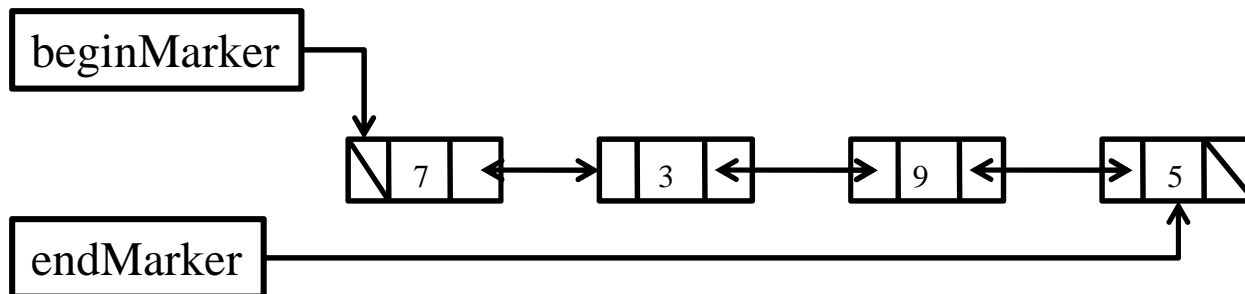
```
private static class Node<AnyType>
{
    public Node( AnyType d, Node<AnyType> p, Node<AnyType> n )
    {
        data = d; prev = p; next = n;
    }

    public AnyType data;
    public Node<AnyType> prev; /* now we can go back */
    public Node<AnyType> next;
}
```

# Linked List Implementation

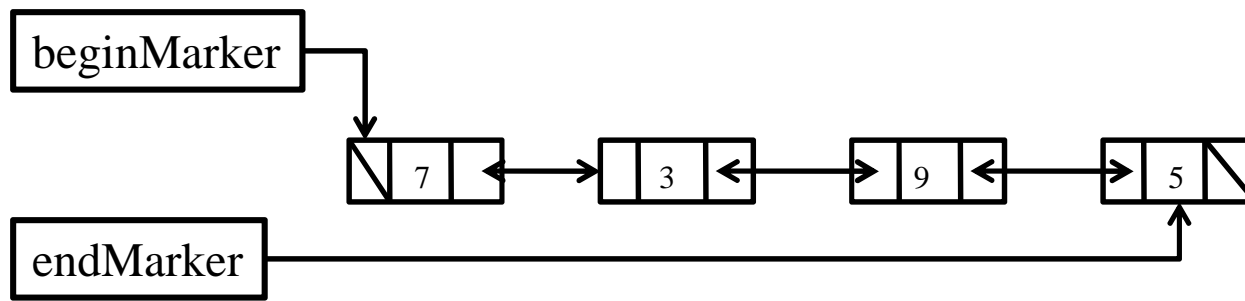
```
public class MyLinkedList<AnyType> implements Iterable<AnyType>
{
    /* all methods go here (see book, code examples) */

    private int theSize;
    private Node<AnyType> beginMarker;
    private Node<AnyType> endMarker;
}
```





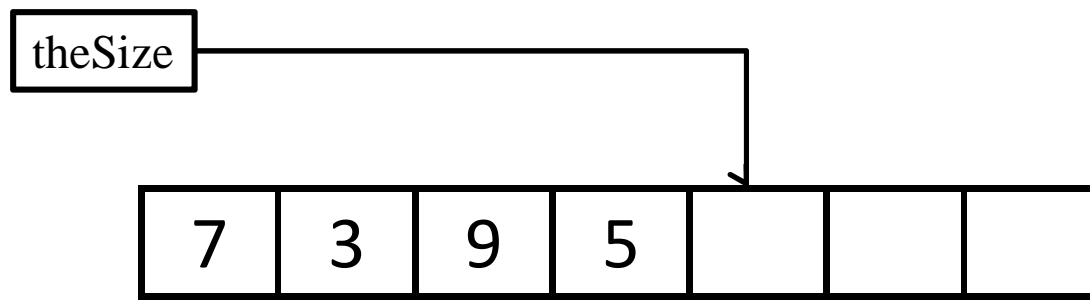
# Operations on the Linked List



Discuss in class the following operations:

1. Find(9); Find(2);
2. Insert(11); at the beginning, at the end, in the middle;
3. Remove(9);

# Array Implementation



Discuss in class the following operations:

1. Find(9); Find(2);
2. Insert(11); at the beginning, at the end, in the middle;
3. Remove(9);

# Review: Big O Notation

**DEFINITION:** The Big-O notation

$T(n) = O(f(n))$  if there exist constants  $c$  and  $n'$  such that:  $T(n) \leq c f(n)$  for all  $n \geq n'$

# Summary: Some Complexities

	Linked list	Array
AnyType get( int idx)	?	?
int find(AnyType x)	?	?
void add(x,idx) /* before/after */	?	?
void remove(int idx)	?	?
InsertAnywhere(x)	?	?

# Summary: Some Complexities

	Linked list	Array
<code>AnyType get( int idx)</code>	$O(k)$ where $k=idx$	$O(1)$
<code>int find(AnyType x)</code>	$O(n)$	$O(n)$
<code>void add(x,idx)</code> <code>/* before/after */</code>	$O(1)$	$O(n)$
<code>void remove(int idx)</code>	$O(1)$	$O(n)$
<code>InsertAnywhere(x)</code>	$O(1)$	$O(1)$

# Iterators

- Review questions from Weiss, Chapt. 3.3
- What is an *Iterator* ?
- What is an *Iterable* ?
- What is a *Collection* ?

# Answers...

```
public interface Iterator<AnyType>
{
    boolean hasNext( );
    AnyType next( );
    void remove( );
}
```

**An** Iterable provides a method named `iterator` that returns an object of type `Iterator`. It can be used in a for loop (see book...).

```
public interface Collection<AnyType>
    extends Iterable<AnyType>
{
    . . .
    boolean add( AnyType x );
    boolean remove( AnyType x );
    . . .
}
```

# Application: Polynomial ADT

Attempt 1:

- linked list implementation:
- $A_i$  is the coefficient of the  $x^{i-1}$  term:

$$5 + 2x + 3x^2 \quad ( \quad 5 \quad 2 \quad 3 \quad )$$

$$7 + 8x \quad ( \quad 7 \quad 8 \quad )$$

$$3 + x^2 \quad ( \quad 3 \quad 0 \quad 1 \quad )$$

**Problem?**





# Sparse Vector Data Structure:

$$4 + 3x^{2001}$$

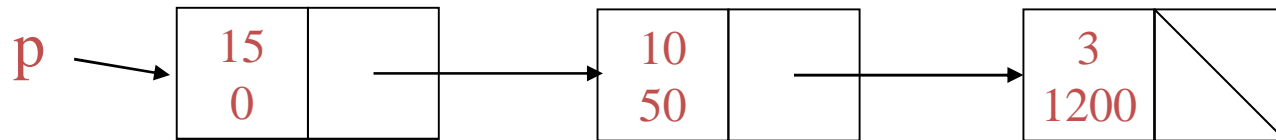
( <4 0>    <3 2001> )



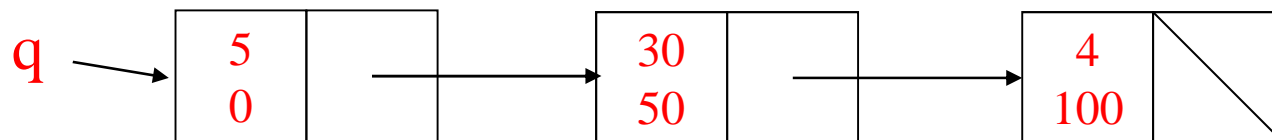
# Addition of Two Polynomials?

*Complexity?*

$$15 + 10x^{50} + 3x^{1200}$$



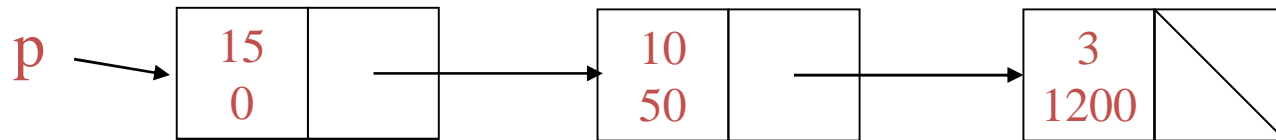
$$5 + 30x^{50} + 4x^{100}$$



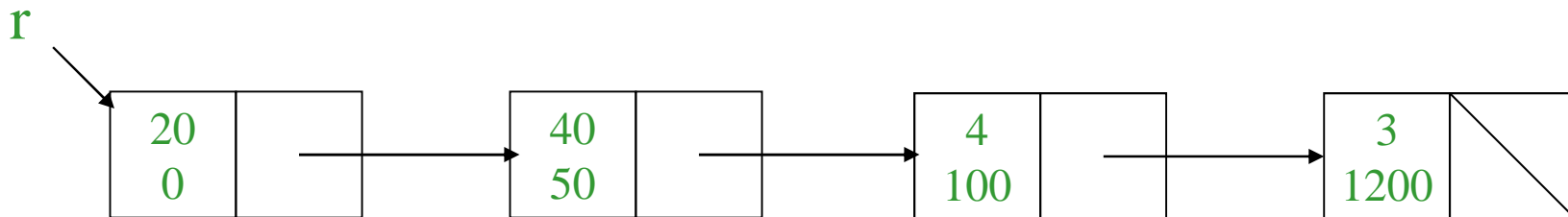
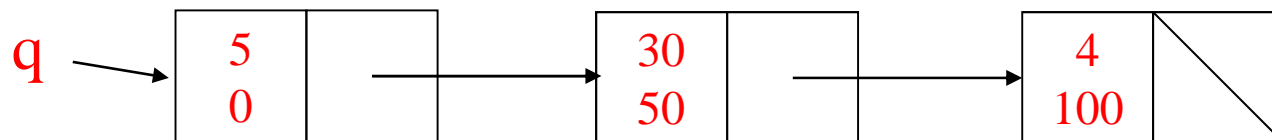
# Addition of Two Polynomials

One pass down each list:  $O(m+n)$

$$15 + 10x^{50} + 3x^{1200}$$



$$5 + 30x^{50} + 4x^{100}$$



# Sparse Matrices

18	0	33	0	0	0
0	0	0	0	0	0
0	0	0	0	0	0
0	0	0	99	0	0
0	0	0	0	0	0
0	0	0	0	0	27

```
( <row  (<column data> <column data> ...) >
  <row  (<column data> <column data> ...) >
  ... )
```

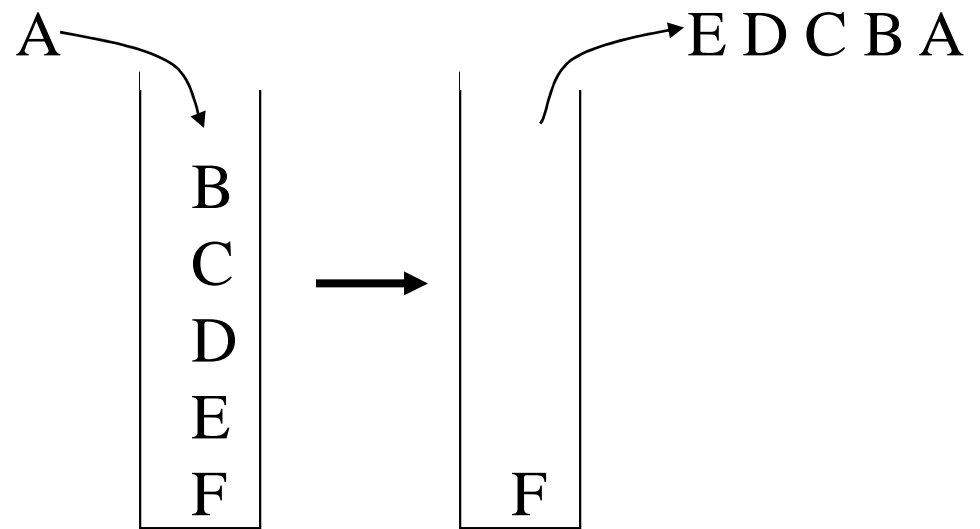
```
( <1 (<1 18> <3 33>)>, <4 (<4 99>)> <6 (<6 27>)> )
```

```
( <row column data> <row column data> ... )
```

```
( <1, 1, 18> <1, 3, 33> <4, 4, 99> <6, 6, 27> )
```

# First Example: Stack ADT

- LIFO: Last In First Out
- Stack operations
  - create
  - destroy
  - push
  - pop
  - top
  - is\_empty



# Stacks in Practice

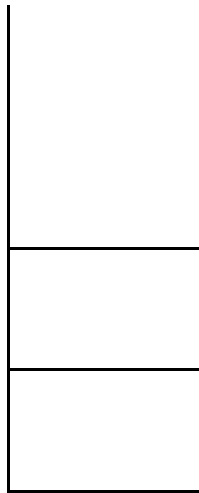
- Function call stack
- Removing recursion
- Balancing symbols (parentheses)
- Evaluating Postfix Notation

Infix:  $7 + 8 * 4 - 30$

Postfix:  $7 8 4 * + 30 -$

# Infix to Postfix with a Stack

7 + 8 \* 4 - 3 0 #



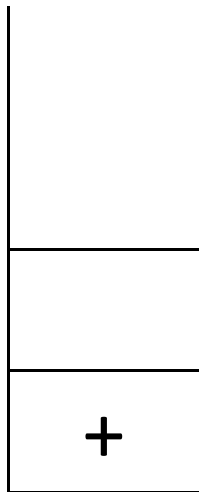
7



# Infix to Postfix with a Stack

7 + 8 \* 4 - 30 #

Stack:



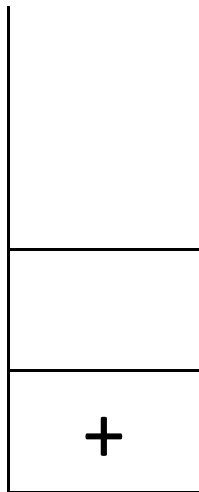
Output:

7

# Infix to Postfix with a Stack

7 + 8 \* 4 - 3 0 #

Stack:



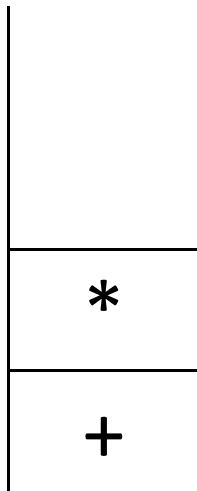
Output:

7 8

# Infix to Postfix with a Stack

7 + 8 \* 4 - 3 0 #

Stack:



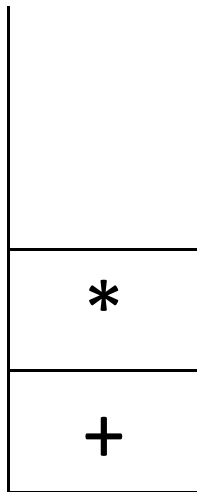
Output:

7 8

# Infix to Postfix with a Stack

7 + 8 \* 4 - 30 #

Stack:



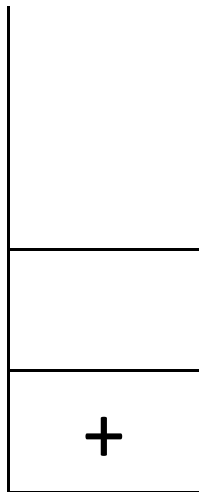
Output:

7 8 4

# Infix to Postfix with a Stack

7 + 8 \* 4 - 30

Stack:



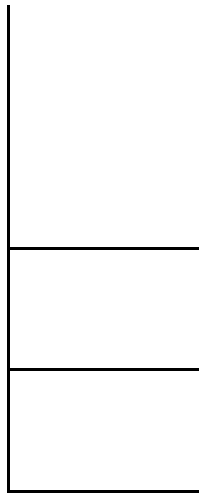
Output:

7 8 4 \*

# Infix to Postfix with a Stack

7 + 8 \* 4 - 30 #

Stack:



(empty stack)

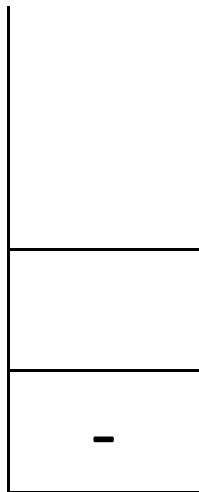
Output:

7 8 4 \* +

# Infix to Postfix with a Stack

7 + 8 \* 4 - 30 #

Stack:



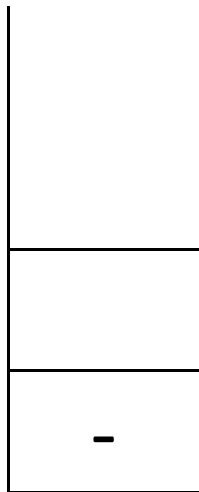
Output:

7 8 4 \* +

# Infix to Postfix with a Stack

7 + 8 \* 4 - 30 #

Stack:



Output:

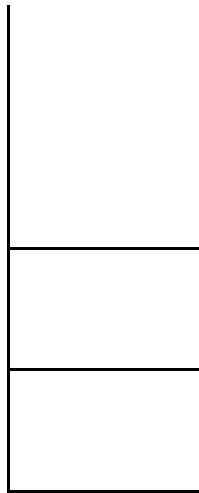
7 8 4 \* + 30



# Infix to Postfix with a Stack

7 + 8 \* 4 - 30 #

Stack:



Output:

7 8 4 \* + 30 -

# Array vs. Linked List

- Too much space, or not enough space
- Not as complex
- Could make array more robust
- Kth element accessed “easily”
- Insert requires shift
- Can grow as needed
- More memory per item in the queue
- Linked list code more complex
- Kth element access requires linear scan??

# Second Example: Queue ADT

We briefly mentioned  
it in Lecture 01...

- FIFO: First In First Out
- Queue operations

create

destroy

enqueue

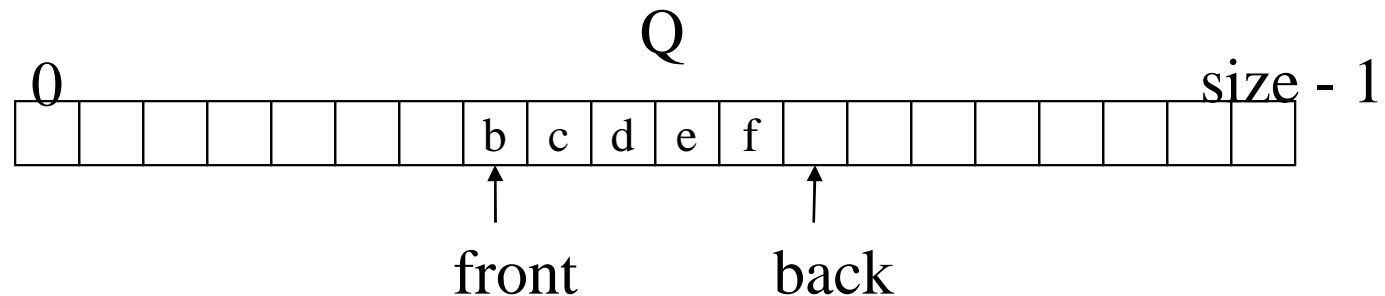
dequeue

is\_empty



<http://courses.cs.vt.edu/csonline/DataStructures/Lessons/QueuesImplementationView/applet.html>

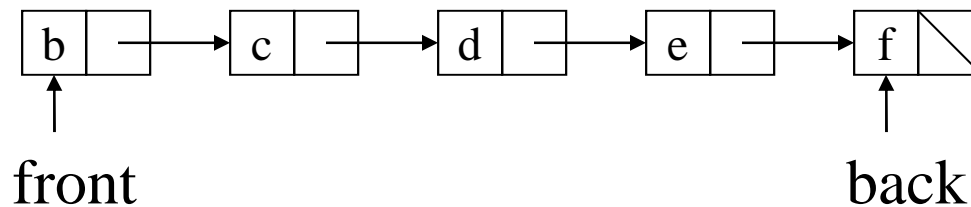
# Circular Array Queue Data Structure



```
enqueue(Object x) {  
    Q[back] = x ;  
    back = (back + 1) % size  
}  
  
dequeue() {  
    x = Q[front] ;  
    front = (front + 1) % size;  
    return x ;  
}
```

How test for empty list?  
How to find K-th  
element in the queue?  
What is complexity of  
these operations?  
Limitations of this  
structure?

# Linked List Queue Data Structure



```
void enqueue(Object x) {
    if (is_empty())
        front = back = new Node(x)
    else
        back.next = new Node(x)
        back = back.next
}
bool is_empty() {
    return front == null
}
```

```
Object dequeue() {
    assert(!is_empty)
    return_data = front.data
    temp = front
    front = front.next
    delete temp
    return return_data
}
```

# Circular Array vs. Linked List

- Too much space, or not enough space
- Not as complex
- Could make array more robust
- Can grow as needed
- More memory per item in the queue
- Linked list code more complex

# Applications of Queues

- Tree traversals, graph traversals
  - Will see in coming lectures
- Wherever fairness is needed:
  - Printer queues
  - Packets in a network
  - http requests at a web server
  - Etc.