

# CSE 373: Data Structures and Algorithms

## Lecture 3: Math Review/Asymptotic Analysis

# Motivation

- So much data!!
  - Human genome:  $3.2 * 10^9$  base pairs
    - If there are  $6.8 * 10^9$  on the planet, how many base pairs of human DNA?
  - Earth surface area:  $1.49 * 10^8$  km<sup>2</sup>
    - How many photos if taking a photo of each m<sup>2</sup>?
    - For every day of the year ( $3.65 * 10^2$ )?
- But aren't computers getting faster and faster?

# Why algorithm analysis?

- As problem sizes get bigger, analysis is becoming *more* important.
- The difference between good and bad algorithms is getting bigger.
- Being able to analyze algorithms will help us identify good ones without having to program them and test them first.

# Measuring Performance: Empirical Approach

- Implement it, run it, time it (averaging trials)
  - Pros?
  - Cons?

# Measuring Performance: Empirical Approach

- Implement it, run it, time it (averaging trials)
  - Pros?
    - Find out how the system effects performance
    - Stress testing – how does it perform in dynamic environment
    - No math!
  - Cons?
    - Need to implement code
    - Can be hard to estimate performance
    - When comparing two algorithms, all other factors need to be held constant (e.g., same computer, OS, processor, load)

# Measuring Performance: Analytical Approach

- Use a simple model for basic operation costs
- Computational Model
  - has all the basic operations:  
+, -, \*, /, =, comparisons
  - fixed sized integers (e.g., 32-bit)
  - infinite memory
  - all basic operations take exactly one time unit (one CPU instruction) to execute

# Measuring Performance: Analytical Approach

- Analyze steps of algorithm, estimating amount of work each step takes
  - Pros?
    - Independent of system-specific configuration
    - Good for estimating
    - Don't need to implement code
  - Cons?
    - Won't give you info exact runtimes optimizations made by the architecture (i.e. cache)
    - Only gives useful information for large problem sizes
    - In real life, not all operations take exactly the same time and have memory limitations

# Analyzing Performance

- General “rules” to help measure how long it takes to do things:

**Basic operations** Constant time

**Consecutive statements** Sum of number of statements

**Conditionals** Test, plus larger branch cost

**Loops** Sum of iterations

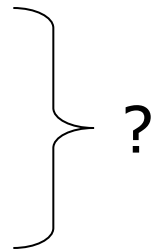
**Function calls** Cost of function body

**Recursive functions** Solve recurrence relation...

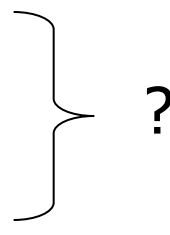


# Efficiency examples

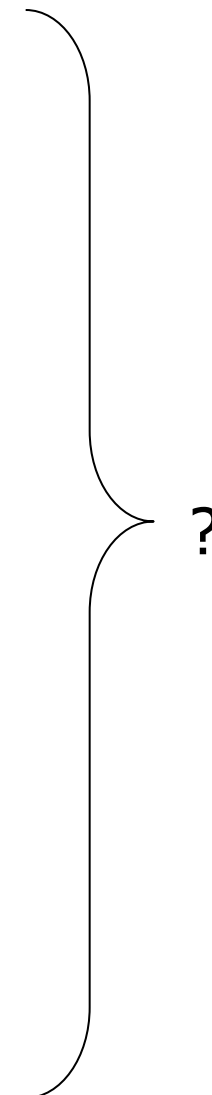
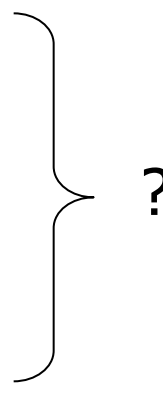
```
statement1;  
statement2;  
statement3;
```



```
for (int i = 1; i <= N; i++) {  
    statement4;  
}
```

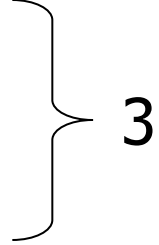


```
for (int i = 1; i <= N; i++) {  
    statement5;  
    statement6;  
    statement7;  
}
```



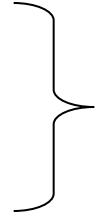
# Efficiency examples

```
statement1;  
statement2;  
statement3;
```



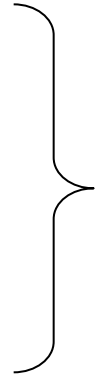
3

```
for (int i = 1; i <= N; i++) {  
    statement4;  
}
```

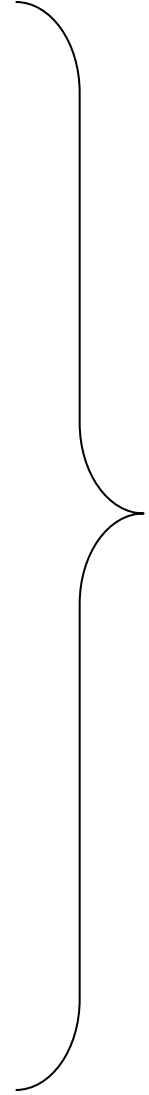


N

```
for (int i = 1; i <= N; i++) {  
    statement5;  
    statement6;  
    statement7;  
}
```



3N



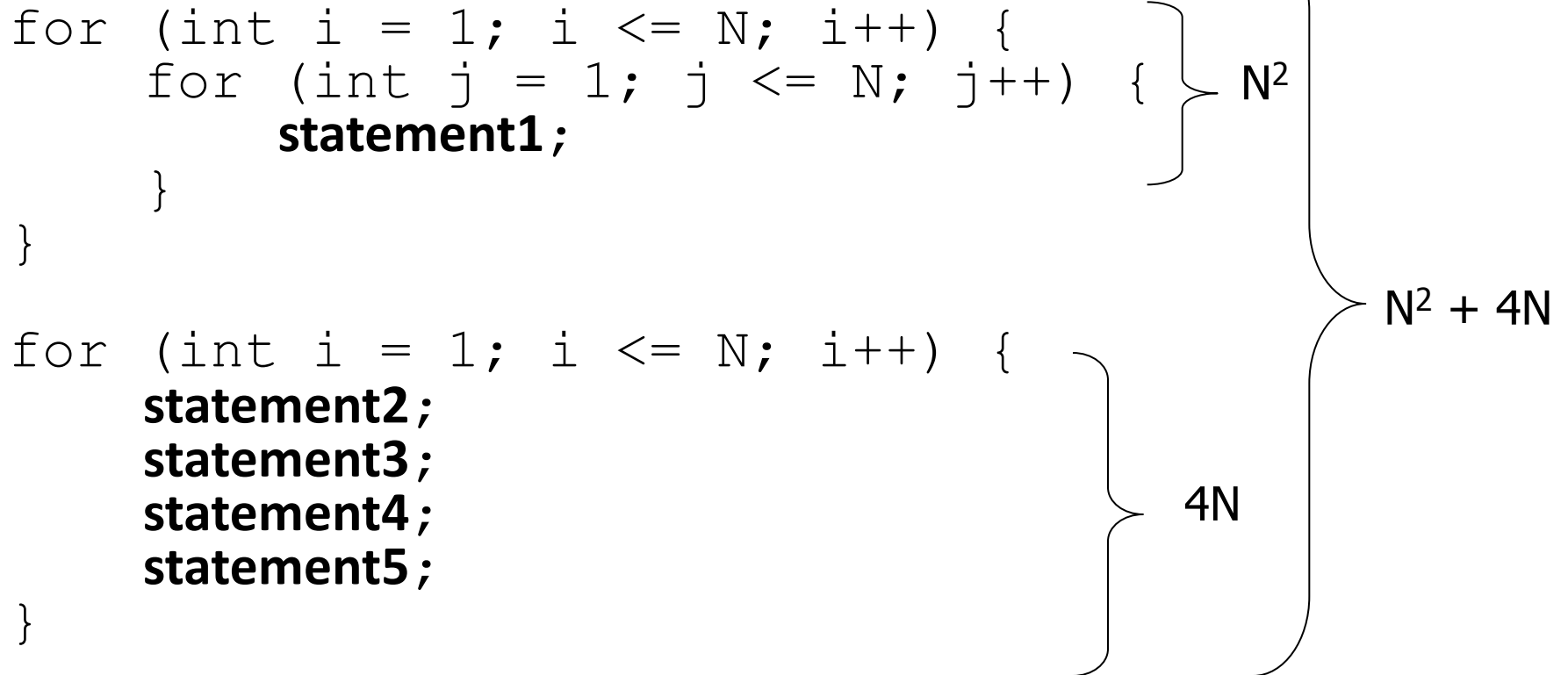
4N + 3

# Efficiency examples 2

```
for (int i = 1; i <= N; i++) {  
    for (int j = 1; j <= N; j++) {  
        statement1;  
    }  
}  
  
for (int i = 1; i <= N; i++) {  
    statement2;  
    statement3;  
    statement4;  
    statement5;  
}
```

The diagram illustrates the complexity analysis of two code blocks. The first block is a nested loop structure where the inner loop iterates over  $j$  from 1 to  $N$  for each iteration of the outer loop over  $i$  from 1 to  $N$ . A brace on the right of the inner loop is labeled with a question mark, and a larger brace on the right of the entire first block is also labeled with a question mark. The second block is a single loop structure where the loop iterates over  $i$  from 1 to  $N$ , and each iteration contains four statements: **statement2**, **statement3**, **statement4**, and **statement5**. A brace on the right of these four statements is labeled with a question mark, and a larger brace on the right of the entire second block is also labeled with a question mark.

# Efficiency examples 2



- How many statements will execute if  $N = 10$ ? If  $N = 1000$ ?

# Relative rates of growth

- most algorithms' runtime can be expressed as a *function* of the input size  $N$
- **rate of growth**: measure of how quickly the graph of a function rises
- goal: distinguish between fast- and slow-growing functions
  - we only care about very large input sizes  
(for small sizes, most any algorithm is fast enough)
  - this helps us discover which algorithms will run more quickly or slowly, for large input sizes
- most of the time interested in worst case performance; sometimes look at best or average performance

# Growth rate example

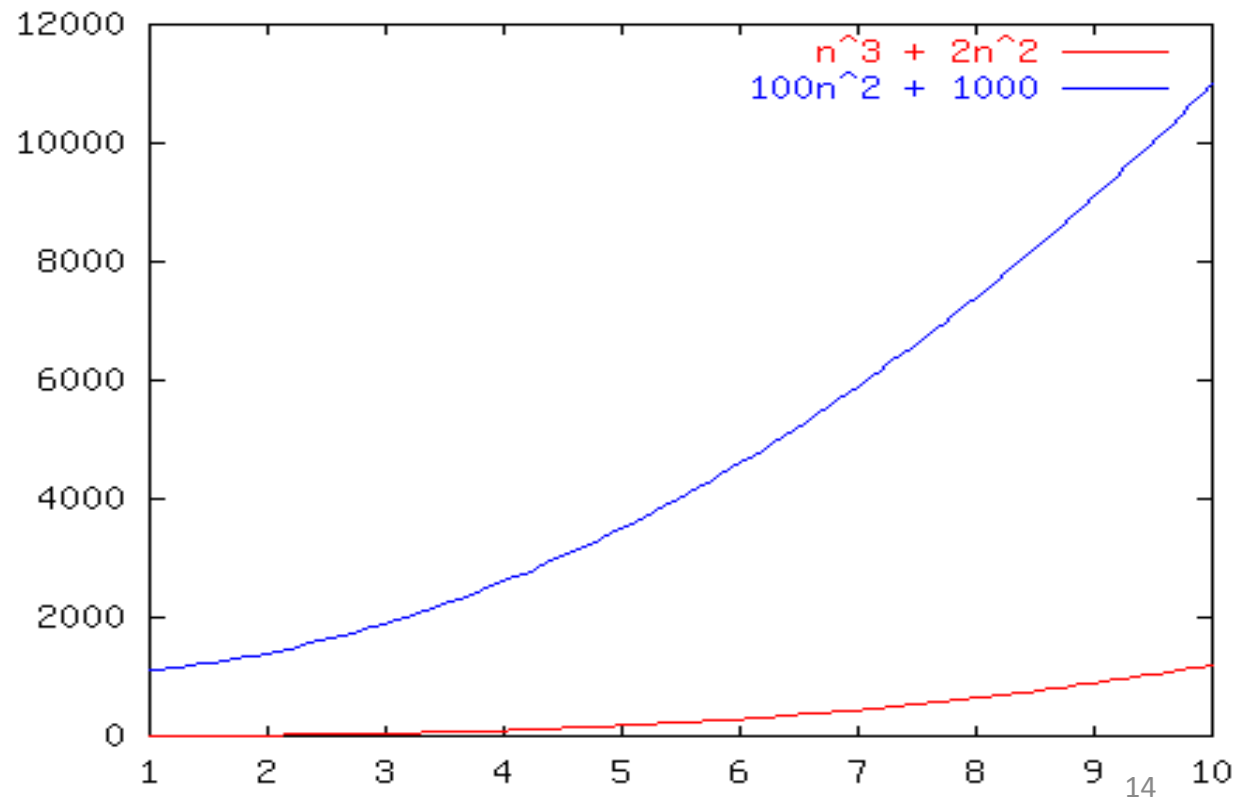
Consider these graphs of functions.

Perhaps each one represents an algorithm:

$$n^3 + 2n^2$$

$$100n^2 + 1000$$

- Which grows faster?



# Growth rate example

- How about now?

