

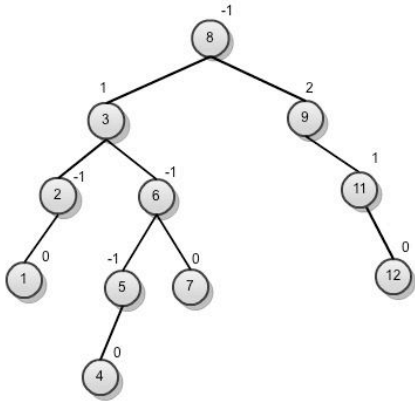
**CSE 373, Spring 2011
Final Key**

1. Sorting (12 Points)

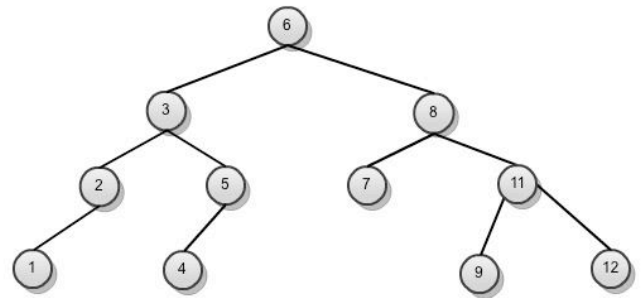
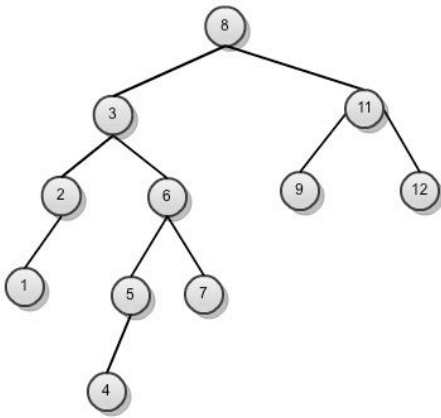
Part	Conditions	Answer	Expected Runtime
a	array size 700000, ascending order	insertion sort	$O(n)$
b	array size 350000, random order, no extra memory may be allocated	quick sort	$O(n \log n)$
c	array size 1000000, descending order	merge sort	$O(n \log n)$
d	array size 2500000 containing zip codes (i.e. values between 0 - 99999), random order	bucket sort	$O(n)$

2. AVL Trees (10 Points)

a.



b.



3. Heap Implementation (12 Points)

Part	Answer
a	<pre>public void delete(int p) { if (p <= 0 p > size) { throw new NoSuchElementException(); } array[p] = Integer.MIN_VALUE; bubbleUp(p); this.remove(); }</pre>
b	<p>Checking that the position is valid and throwing the exception is $O(1)$. Setting the element to be deleted to <code>Integer.MIN_VALUE</code> is $O(1)$. Bubbling up the value to be deleted is worst case $O(\log n)$. Performing a regular remove from the top of the heap is $O(\log n)$. Therefore, we have $O(1 + 1 + \log n + \log n) = O(\log n)$.</p>

4. Hashing (12 Points)

Part	Answer
a	<pre> value +-----+ 0 10 +-----+ 1 22 +-----+ 2 82 +-----+ 3 53 +-----+ 4 +-----+ 5 55 +-----+ 6 92 +-----+ 7 R +-----+ 8 +-----+ 9 75 +-----+ </pre>
b	<p>Yes, 86 failed to be inserted because a bucket couldn't be found after trying half of the entries. The second 55 failed to be inserted because it was already in the set.</p>
c	7
d	10
e	.7

5. Topological Sort (10 points)

A B D E C F H G I J

6. Minimum Spanning Trees (12 points)

Part	Answer
a	A, B, C, D, E, G, H, I, J, F edges: AB = 2 BC = 3 BD = 5 DE = 4 DG = 6 EH = 7 GI = 8 IJ = 1 EF = 13 Total: 49
b	IJ = 1 AB = 2 BC = 3 DE = 4 BD = 5 DG = 6 EH = 7 GI = 8 EF = 13 Total: 49

7. Graph Implementation (12 points)

```
// BFS-based Solution 1: looking at previous/source node to determine node's set
public void friendsAndEnemies(V v1, Set<V> friends, Set<V> enemies) {
    this.clearVertexInfo();
    Queue<V> queue = new LinkedList<V>();
    queue.add(v1);
    friends.add(v1);
    this.vertexInfo.get(v1).visited = true;

    while(!queue.isEmpty()) {
        V v = queue.remove();

        for (V n : this.neighbors(v)) {
            VertexInfo<V> vi = this.vertexInfo.get(n);
            if (!vi.visited) {
                vi.visited = true;
                queue.add(n);

                if (friends.contains(v)) {
                    enemies.add(n);
                } else {
                    friends.add(n);
                }
            }
        }
    }
}
```

```
// BFS-based Solution 2: looking at distance to determine node's set
public void friendsAndEnemiesDistance(V v1, Set<V> friends, Set<V> enemies) {
    this.clearVertexInfo();
    vertexInfo.get(v1).distance = 0;

    Queue<V> queue = new LinkedList<V>();
    queue.offer(v1);

    friends.add(v1);

    while(!queue.isEmpty()) {
        V v = queue.poll();

        for (V n : this.neighbors(v)) {
            VertexInfo<V> vi = this.vertexInfo.get(n);
            if (vi.distance == Integer.MAX_VALUE) {
                vi.distance = vertexInfo.get(v).distance + 1;

                queue.offer(n);

                if (vi.distance % 2 == 0) {
                    friends.add(n);
                } else {
                    enemies.add(n);
                }
            }
        }
    }
}
```

```

// DFS-based Solution 1: looking at distance to determine node's set
public void friendsAndEnemies(V v1, Set<V> friends, Set<V> enemies) {
    this.clearVertexInfo();
    friends.add(v1);
    vertexInfo.get(v1).distance = 0;

    for (V neighbor : neighbors(v1)) {
        friendsAndEnemies(neighbor, friends, enemies, 0, false);
    }
}

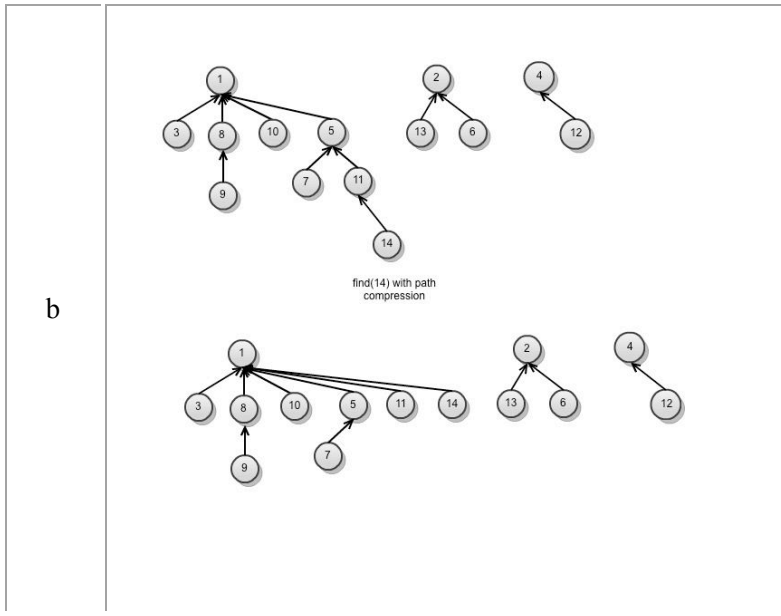
public void friendsAndEnemies(V v1, Set<V> friends, Set<V> enemies, int distance,
boolean isFriend) {
    if (distance < vertexInfo.get(v1).distance) {
        vertexInfo.get(v1).distance = distance;
        if (isFriend) {
            friends.add(v1);
            enemies.remove(v1);
        } else {
            enemies.add(v1);
            friends.remove(v1);
        }

        for (V neighbor : neighbors(v1)) {
            friendsAndEnemies(neighbor, friends, enemies, distance + 1, !isFriend);
        }
    }
}
}

```

8. Disjoint Sets (10 points)

Part	Answer
a	<p>The diagrams illustrate the following sequence of operations:</p> <ul style="list-style-type: none"> union(1, 3): Node 3 points to node 1. union(5, 7): Node 7 points to node 5. union(8, 9): Node 9 points to node 8. union(1, 8): Node 8 points to node 1. union(2, 13): Node 13 points to node 2. union(1, 10): Node 10 points to node 1. union(11, 14): Node 14 points to node 11. union(5, 11): Node 11 points to node 5. union(2, 6): Node 6 points to node 2. union(4, 12): Node 12 points to node 4. union(x, y) = union(1, 5): Node 5 points to node 1.



9. BTrees (10 points)

Part	Answer
a	
b	<p>The point of B-trees is to store data on disk. For them to work efficiently, nodes should all fit within one disk block, which is a fixed size. Internal nodes hold keys and links. External nodes hold keys and data. Therefore, a different number of keys/links than keys/data may fit into a single disk block. This means that different values of M and L may be required for internal and external nodes.</p>