

1. Big-Oh

Calculate the exact value of the variable `sum` after the following code fragment, in terms of variable `n`. Use summation notation to compute a closed-form solution. Then use this value to give a tightly bounded Big-Oh analysis of the runtime of the code fragment.

```
int sum = 0;
for (int i = 1; i <= n - 3; i++) {
    for (int j = 1; j <= i + 4; j += 2) {
        sum += 3;
    }

    sum++;
}

for (int i = 1; i <= 100; i++) {
    sum++;
}
```

2. Sorting

Consider the following array of `int` values.

```
[7, 1, 6, 12, -3, 8, 4, 21, 2, 30, -1, 9]
```

- Write the contents of the array after 3 passes of the outermost loop of **bubble** sort.
- Write the contents of the array after 4 passes of the outermost loop of **selection** sort.
- Write the contents of the array after 5 passes of the outermost loop of **insertion** sort.
- Write the contents of the array after all the recursive calls of **merge** sort have finished (before the very last merge begins).
- Write the contents of the array after the first partitioning of **quick** sort has finished (before recursive calls). Assume that the first element is chosen as the pivot.

Please show your work. You do not have to write an entirely new array after each pass of the algorithm, but since the final answer depends on every add/remove being done correctly, you may wish to show the array at various important stages (i.e. after each pass or recursive call) to help earn partial credit in case of an error.

3. Trees and Heaps

Given the following String elements:

"m", "x", "z", "s", "d", "b", "i", "t", "r", "g", "w", "k", "h"

Draw the tree that results when all of the above elements are added (in the given order) to each of the following initially empty data structures:

- a. A binary search tree (BST).
 - i. Draw the tree that results after adding all the elements.
 - ii. Draw the tree that results after removing "k".
 - iii. Draw the tree that results after removing "t".
 - iv. Draw the tree that results after removing "x".
 - v. Draw the tree that results after removing "m".
- b. An AVL tree. Draw the tree that results after inserting all elements. Draw a new tree each time a rotation is necessary and say which kind of rotation was needed.
- c. A minimum binary heap.
 - i. Draw the heap after adding all the elements.
 - ii. Perform three removes on the heap. Draw a new heap after each remove.

Please show your work.

4. Set Programming

Part A: Implementation

One of the operations commonly performed on sets is intersection. The intersection of two sets contains all the items that the two sets have in common. The `StringSet` interface has been altered to have an `intersect` method. This method takes a `StringSet` as a parameter and returns a new `StringSet` that contains all the `Strings` that are in both the `StringSet` on which the method is called and the `StringSet` parameter. This method should not alter in any way the `StringSet` on which the method is called or the `StringSet` parameter.

The `StringTreeSet` class must now implement the `intersect` method. The public method has been written below. Write the recursive helper method `intersect` that this public method calls to populate the intersection `StringSet` to contain all the `Strings` that both this `StringTreeSet` instance and the other `StringSet` have. You may assume all of the other methods are implemented as discussed in lecture.

```
public interface StringSet {
    public boolean add(String value);
    public boolean contains(String value);
    public StringSet intersect(StringSet other);
    public void print();
    public boolean remove(String value);
    public int size();
}

// A binary search tree implementation of a Set for Strings.
public class StringTreeSet implements StringSet {
    protected StringTreeNode root;

    ...

    public StringSet intersect(StringSet other) {
        StringSet intersection = new StringTreeSet();
        intersect(root, other, intersection);
        return intersection;
    }

    // YOUR RECURSIVE INTERSECT HELPER SHOULD GO HERE

}
```

Part B: Analysis

Consider the following code fragment:

```
StringSet s3 = s1.intersect(s2);
```

Assuming that `s2` has about half of the items that `s1` has, give the worst case running time of this method call if both `s1` and `s2` are of type `StringTreeSet` (i.e. a binary search tree implementation). Explain how you arrived at this running time.