

# CSE 373 Optional Section

Led by Yuanwei, Luyi

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# Today

- Proof by Induction
- Big-Oh
- Algorithm Analysis

# Proof by Induction

## Base Case:

1. Prove  $P(0)$  (sometimes  $P(1)$ )

## Inductive Hypothesis:

2. Let  $k$  be an arbitrary integer  $\geq 0$

3. Assume that  $P(k)$  is true

## Inductive Step

4. have  $P(k)$  is true, Prove  $P(k+1)$  is true

## Conclusion:

5.  $P(n)$  is true for  $n \geq 0$  (or  $1\dots$ )

# Examples

$$\sum_{i=1}^N i^2 = 1 + 2^2 + 3^2 + 4^2 + \dots + n^2 = \frac{N(N+1)(2N+1)}{6} \quad \text{for all } n \geq 1$$

$$\sum_{i=0}^N 2^i = 2^0 + 2^1 + 2^2 + \dots + 2^n = 2^{n+1} - 1$$

Extra

$$\sum_{i=1}^n \frac{1}{i(i+1)} = \frac{n}{n+1} \quad \text{where } n \in \mathbb{Z}^+$$

# Logarithms

- log in CS means log base of 2
- log grows very slowly
- $\log AB = \log A + \log B$ ;  $\log(A/B) = \log A - \log B$
- $\log(N^k) = k \log N$ 
  - Eg.  $\log(A^2) = \log(A * A) = \log A + \log A = 2 \log A$
- distinguish  $\log(\log x)$  and  $\log^2 x$  --  $(\log x)(\log x)$

# Big-Oh

- We only look at worst case
- Big input
- Ignore constant factor and lower order terms
  - Why?
- Definition:

*$g(n)$  is in  $O(f(n))$  if there exist constants  $c$  and  $n_0$  such that  $g(n) \leq c f(n)$  for all  $n \geq n_0$*

- Also lower bound and tight bound

We use  $O$  on a function  $f(n)$  (for example  $n^2$ ) to mean the **set of functions** with asymptotic behavior **less than or equal to**  $f(n)$

# Big-Oh Practice

- Prove that  $5n^2+3n$  is  $O(n^2)$ 
  - Key point
    - Find constant  $c$  and  $n_0$

# Big-Oh Practice

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Possible  $c$  and  $n_0$ :

$$c = 8 \text{ and } n_0 = 1$$

$$c = 6 \text{ and } n_0 = 3$$

...



# Math Related

- Series

$$\sum_{i=1}^N A^i = A + A^2 + A^3 + A^4 + \dots = \frac{A^{N+1} - 1}{A - 1}$$

$$\sum_{i=1}^N i = 1 + 2 + 3 + 4 + \dots = \frac{N(N+1)}{2} \approx \frac{N^2}{2}$$

$$\sum_{i=1}^N i^2 = 1 + 2^2 + 3^2 + 4^2 + \dots = \frac{N(N+1)(2N+1)}{6} \approx \frac{N^3}{3}$$

- Very useful for runtime analysis
- On your textbook, p4

# How to analyze the code?

Consecutive statements	Sum of times
Conditionals	Time of test plus slower branch
Loops	Sum of iterations
Calls	Time of call's body
Recursion	Solve recurrence equation

# Examples

```
1.int sunny (int n) {  
    if (n < 10)  
        return n - 1;  
    else {  
        return sunny (n / 2);  
    }  
}
```

```
2.int funny (int n, int sum) {  
    for (int k = 0; k < n * n; +  
+k)  
        for (int j = 0; j < k; j+  
+)  
            sum++;  
    return sum;  
}
```

```
3.int happy (int n, int sum) {  
    for (int k = n; k > 0; k = k - 1) {  
        for (int i = 0; i < k; i++)  
            sum++;  
        for (int j = n; j > 0; j--)  
            sum++;  
    }  
    return sum;  
}
```

# Examples

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3.int happy (int n, int sum) {
    for (int k = n; k > 0; k = k - 1) {
        for (int i = 0; i < k; i++)
            sum++;
        for (int j = n; j > 0; j--)
            sum++;
    }
    return sum;
}
```

Answer:

1.  $O(\log n)$

2.  $O(n^4)$

3.  $O(n^2)$