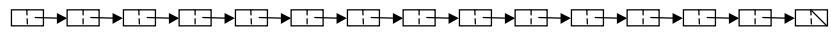




CSE 373: Hash Tables

Chapter 5



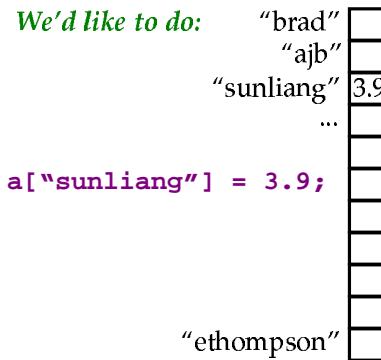
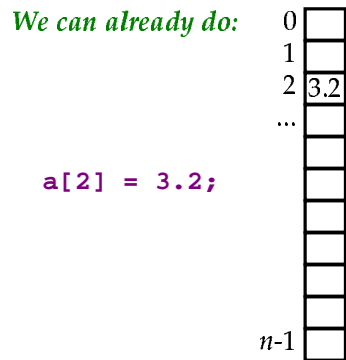
Motivation



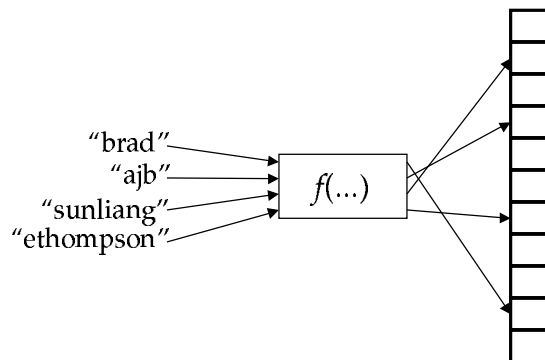
Goal: The ability to store and retrieve information in $O(1)$ time

Current Solutions:

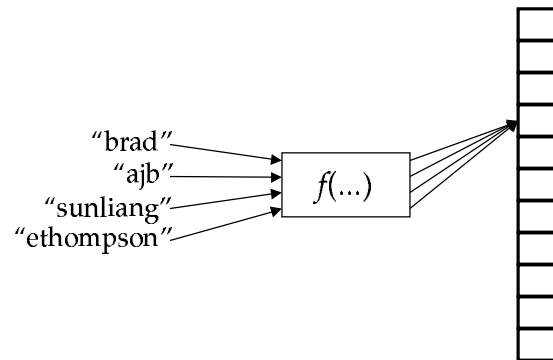
Hash Table Goal



Hash Table Approach



Hashing Conflicts



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Hashing Integers



What's a simple scheme for hashing integers?

```
int HashInteger(int key,int tablesize) {  
  
}
```

- advantages?
- disadvantages?

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General Strategies



- Selecting a good hash function often depends on the set of possible keys
- Using a hash table whose size is a prime number tends to help reduce conflicts

Hashing Strings



General approach:

- convert string to integer
- “mod” integer by table size

Naive approach:

```
int HashString(char* key,int tablesize) {
    int total=0;
    while (*key) {
        total += *key;
    }
    return total%tablesize;
}
```

Problem with Naive Approach



```
HashString("bat",n)
= HashString("tab",n);
= HashString("rad",n);
```

Total Number of Possibilities $\approx 127 \times 12 = 1524$

Useful Number of Possibilities $\approx 26 \times 12 = 312$

Probably not good if hash table size or number of keys is greater than this...

Improved Approach



Read string as base 27 number:

1 27 729

$$\text{b a t} = 2 \times 1 + 1 \times 27 + 20 \times 729 = 14,609$$

$$\text{t a b} = 20 \times 1 + 1 \times 27 + 2 \times 729 = 1,505$$

$$\text{r a d} = 18 \times 1 + 1 \times 27 + 4 \times 729 = 2,961$$

Advantages?

Disadvantages?

Other Ideas



Hash using only a subset of the characters...

- first three?
- last three?
- middle three?
- first, middle, last?
- etc.

Hash Function Design Goals



- Hash to all slots in your table
- Avoid collisions
- Hash as evenly as possible
- Hash as quickly as possible

(Again, note that much of this may depend on the set of possible keys...)

Harsh Hash Reality



- No matter how good your hash function is, collisions will probably occur
- Thus, we also need a collision resolution strategy...
 - separate chaining
 - resizable hash table
 - open addressing

Hash Table Operations



- Main Operations:
 - `void insert(HashedObj& key);`
 - `HashedObj& find(HashedObj& key);`
 - `void remove(HashedObj& key);`
- Normal Creation/Deletion Operations
- No iteration, `FindMin()`/`Max()`, etc.
(why?)

Alternate Hash Table Operations



- Main Operations:

```
void insert(HashedObj& key, Object& data);  
Object& find(HashedObj& key);  
void remove(HashedObj& key);
```

- Similar to the normal interface, but associates some data with each key

Operator Analysis



Hash Table *List* *BST*

problem size

space

`insert()`

`find()`

`remove()`

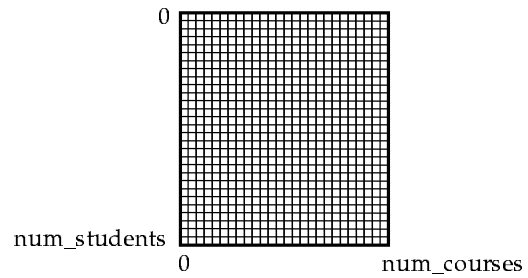
other

Flashback to Day 1



```
const int num_courses = 7000;  
const int num_students = 33000;
```

```
typedef int registry[num_students][num_courses];
```



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