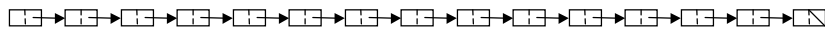


# CSE 373: Graphs

## Chapter 9

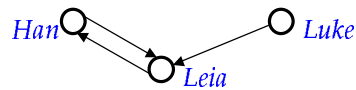


# Graphs



### Graphs:

- A data structure useful for representing relationships between things
- A graph  $G$  is represented as  $G = (V, E)$ 
  - $V$  = a set of *vertices (nodes)*
  - $E$  = a set of *edges* connecting vertices from  $V$



- More general and arbitrary than trees (trees are a restricted type of graph)

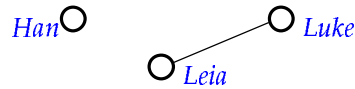
## Directed/Undirected Graphs



- In *directed* graphs, edges have a specific direction:



- In *undirected* graphs, they don't:



- Vertices  $u$  and  $v$  are *adjacent* if  $(u,v) \in E$

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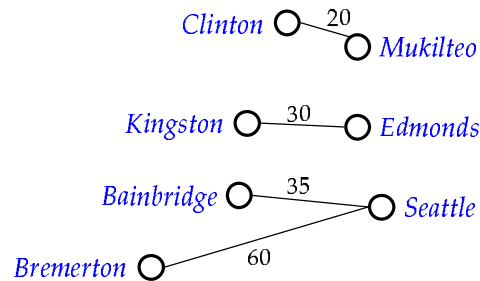
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## Weighted Graphs



*weighted graphs* store a weight with each edge:



(nodes sometimes have weights too...)

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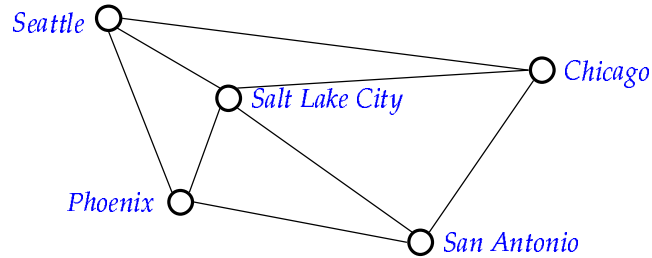
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## Paths



A *path* is a list of vertices  $v_1, v_2, \dots, v_n$  such that  $(v_i, v_{i+1}) \in E$ .



$p = \text{Seattle, Salt Lake City, Phoenix, San Antonio, Phoenix, Seattle}$

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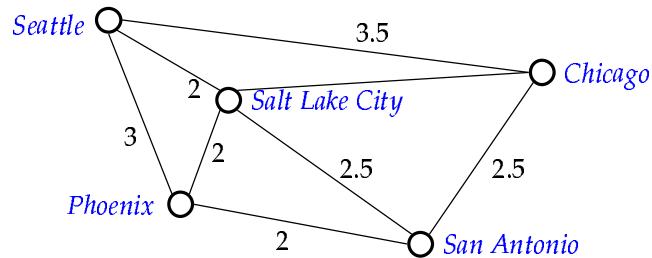
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## Path Length and Cost



- The *length* of a path is the number of edges
- Its *cost* is the sum of the edges' weights



$\text{length}(p) = 5$

$\text{cost}(p) = 11$

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## Simple Paths and Cycles



A *simple path* repeats no vertices (except the first can also be the last):

- Seattle, Salt Lake City, Phoenix, San Antonio
- Seattle, Salt Lake City, San Antonio, Phoenix, Seattle

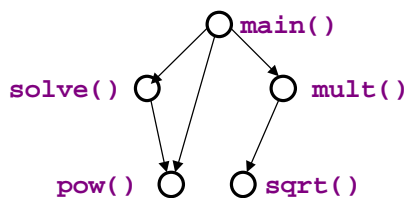
A *cycle* is a path that starts and ends at the same node

- Seattle, Salt Lake City, San Antonio, Phoenix, Seattle  
(For undirected graphs edges cannot appear twice)

## Directed Acyclic Graphs



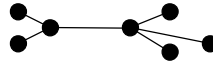
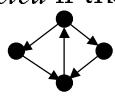

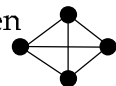
*Directed Acyclic Graphs (DAGs)* are directed graphs with no cycles



$trees \subset DAGs \subset directed\ graphs$

## Connectivity



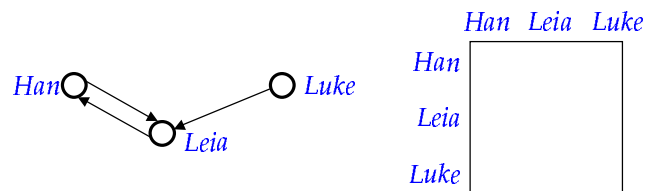
- Undirected graphs are *connected* if there is a path between any two vertices 
- Directed graphs are *strongly connected* if there is a path between any two vertices 
- It is *weakly connected* if it's connected when direction is ignored 
- A *complete* graph is one that has an edge between every pair of vertices 

## Graph Implementations



**Adjacency Matrix:** A  $|V| \times |V|$  array in which:

- element  $(u,v)$  is 1 if there is an edge  $(u,v)$
- it is 0 otherwise
- for weighted graphs, store weights rather than 1/0

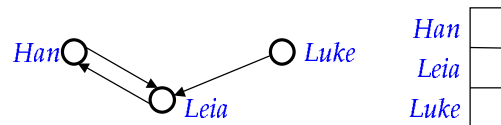


- space requirements?

## Graph Implementations



**Adjacency Lists:** A  $|V|$ -ary array in which each entry stores a list of all adjacent vertices



Space requirements?

How could we index into an adjacency list or matrix when nodes are named?

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## Graph Applications



- Storing things that are graphs by nature (AI)
  - distances between cities
  - airline flights, travel options
  - relationships between people, things
  - distances between rooms in the game Clue
- Compilers
  - *callgraph* – which functions call which other ones
  - *dependence graphs* – which variables are defined and used at which statements

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