

Virtual Memory

Virtual Memory is the process of mapping a logical address space numbered from 0 to the physical address space of the computer so that the RAM serves as a cache for the program's memory stored on the disk. The mapping is performed on a page basis.

© Larry Snyder, 2000. All rights reserved.

Virtual Memory

Virtual memory solves several critical problems:

- Simplifying RAM memory management by separating the address of memory from its physical location
- Providing protection for users by giving them their own address spaces
- Simplifying compilation and usage of libraries by allowing all programs to begin addressing memory from zero

© Larry Snyder, 2000. All rights reserved.

Virtual to Physical Translation

Physical memory (RAM) is divided into pages -- contiguous sequences of memory typically in the 4KB -- 16KB range

The pages are blocks of a fully associative cache for the memory of the program which is stored in the next lower level of the memory hierarchy (disk)

Virtual Address: 31 30 29 | 12 11 | 2 1 0
Virtual Page Number | Page Offset

Translation

Physical Address: 29 28 27 | 12 11 | 2 1 0
Physical Page Number | Page Offset

© Larry Snyder, 2000. All rights reserved.

Page Table

V Page Table

Layout in Physical Memory

Page table and disk addresses may be kept in separate tables

© Larry Snyder, 2000. All rights reserved.

VM Considerations

- Large page sizes allow the huge miss penalties to be amortized over many references
- Avoiding premature page replacement, e.g. by address collisions, extends page life and reduces fault rate ... thus full associativity
- Page faults can be handled in software, which can use cleverness to reduce fault rate
- Write back is needed because of the high cost of writes

© Larry Snyder, 2000. All rights reserved.

Implementing Address Translation

Virtual Address: 31 30 | 12 11 | 1 0
Virtual Page Number | Offset

Pg Table Reg → V → Page Table

Physical page number

0 = fault

Physical Address: 29 28 | 12 11 | 1 0
Physical Page Number | Offset

Write Back requires a "dirty" bit

© Larry Snyder, 2000. All rights reserved.

Page Replacement Strategy

When the Valid = 0, a page fault is signaled.

Some page must be replaced -- pick the page that will be used furthest into the future: Opt

Replace the least recently used (LRU) page.

LRU strategies are effective, but expensive.

"Use" bits can be a decent approximation.

Space required for page tables can be substantial:

4K pages imply 20 bits of virtual address.

At x bytes per entry implies xMB for full page table.

Represent only the prefix of the table using base + extent.

Grow low addresses and high addresses separately by using the msb to indicate which part of the VM space is being used.