

# MIPS

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MIPS is a “computer family”

- R2000/R3000 (32-bit)
- R4000/4400 (64-bit)
- R8000 (for scientific & graphics applications)
- R10000 & R12000 (64-bit & out-of-order execution)

MIPS originated as a Stanford research project

**M**icroprocessor without **I**nterlocked **P**ipe **S**tages

MIPS was bought by Silicon Graphics (SGI) & is now independent

- focuses on embedded processors for game machines (e.g., Nintendo)

MIPS is a RISC

## MIPS Registers

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Part of the state of a process

Thirty-two 32-bit **general purpose registers (GPRs)**: \$0, \$1, ..., \$31

- integer arithmetic
- address calculations
- temporary values

By convention software uses different registers for different purposes  
(*next slide*)

A 32-bit **program counter (PC)**

Two 32-bit registers **HI** and **LO** used specifically for multiply and divide

- HI & LO concatenated for the product
- LO for the quotient; HI for the remainder

Thirty-two 32-bit registers used for floating-point arithmetic: **\$f0**, **\$f1**, ..., **\$f31**

- often used as 16 64-bit registers for double precision FP

Other special-purpose registers (*later*)

## MIPS Register Names and gcc Conventions

Register	Name	Use	Comment
\$0	zero	always 0	cannot be written
\$1	\$at	reserved for assembler	don't use it!
\$2, \$3	\$v0, \$v1	function return	
\$4 - \$7	\$a0 - \$a3	pass first 4 procedure/function arguments	
\$8 - \$15	\$t0 - \$t7	temporaries	caller saved (callee uses them without saving them)
\$16 - \$23	\$s0 - \$s7	temporaries	callee saved (caller assumes they will be available on function return)
\$24, \$25	\$t8, \$t9	temporaries	caller saved
\$26, \$27	\$k0, \$k1	reserved for the OS	don't use them!
\$28	\$gp	pointer to global static memory	points to the middle of a 64KB block in the static data segment ( <i>next slide</i> )
\$29	\$sp	stack pointer	points to the last allocated stack location ( <i>next slide</i> )
\$30	\$fp	frame pointer	points to the stack frame ( <i>later</i> )
\$31	\$ra	procedure/function return address	

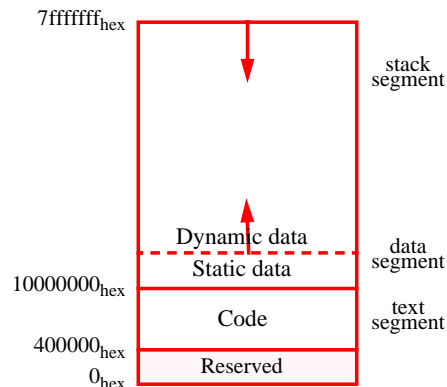
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## Memory Usage

A software convention



**text segment:** the code

**data segment**

- **static data:** objects whose size is known to the compiler & whose lifetime is the whole program execution
- **dynamic data:** objects allocated as the program executes (`malloc`, `new`)

**stack segment:** LIFO process-local storage

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## MIPS Load-Store Architecture

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Most instructions compute on operands stored in registers

- load data into a register from memory
- compute in registers
- the result is stored into memory

For example:

```
a = b + c
d = a + b
```

is “compiled” into:

```
load b into register $x
load c into register $y
$z ← $x + $y
store $z into a
$z ← $z + $x
store $z into d
```

## MIPS Information Units

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Data types and sizes

- byte
- half-word (2 bytes)
- word (4 bytes)
- float (4 bytes using single-precision floating-point format)
- double (8 bytes using double-precision floating-point format)

Memory is byte-addressable

A data type must start on an address evenly divisible by its size in bytes

## Big & Little Endian Byte Order

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Every word starts at an address that is divisible by 4.  
 Which byte in the word is byte 0?  
 How is the data in `.byte a,b,c,d` stored?

**Big Endian**

31	24	23	16	15	8	7	0	byte address
a	b	c	d					0
e	f	g	h					4
i	j	k	l					8
0's	0's	0's	110 <sub>2</sub>					12

**Most** significant byte is the lowest byte address.  
 Word is addressed by the byte address of the **most** significant byte.

**Little Endian**

31	24	23	16	15	8	7	0	byte address
d	c	b	a					0
h	g	f	e					4
l	k	j	i					8
0's	0's	0's	110 <sub>2</sub>					12

**Least** significant byte is the lowest byte address.  
 Word is addressed by the byte address of the **least** significant byte.

## Big & Little Endian Byte Order

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Problems when transferring data structures that contain a mixture of integers & characters between big & little endian computers

	byte address		byte address						
<b>Big Endian</b>	0	J	I	M			0	<b>Little Endian</b>	
	4	S	M	I	T		4		
	8	H	0	0	0		8		
	12	0's	0's	0's	21		12		

Transfer from big endian computer to little endian computer

	byte address							
	0	M	I	J			0	<b>Little Endian</b>
	4	T	I	M	S		4	
	8	0	0	0	H		8	
	12	21	0's	0's	0's		12	

## MIPS Information Units

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MIPS supports both **big-** and **little-endian** byte orders

SPIM uses the byte order of the machine its running on

- Intel: little-endian
- Alpha, SPARC, Mac: big-endian

Words in SPIM are listed from left to right, but byte addresses are little-endian within a word

