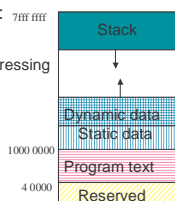


Program and memory layout

- By convention the layout is:
 - Note that only half of the addressing space is taken by user
Other half is O.S.



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1

Procedures

- Procedures/functions are the major program structuring mechanism
- Calling and returning from a procedure requires a *protocol* between *caller* and *callee*
- Protocol is based on conventions

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2

Procedures/Functions -- Protocol

- Each machine (compiler?) has its own set of protocol(s)
- Protocol: combination of hardware/software
 - e.g., "jal" is hardware
 - use of register \$29 as \$sp is software
- Protocol: sequence of steps to be followed at each call and each return
 - controlled by hardware and/or software
- In RISC machines
 - hardware performs simple instructions
 - software (compiler/assembler) controls sequence of instructions

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3

Program stack

- Each executing program (process) has a *stack*
- Stack = dynamic data structure accessed in a LIFO manner
- Program stack automatically allocated by O.S.
- At the start of the program, register \$sp (\$29 in MIPS) is automatically loaded to point to the *first empty slot* on top of stack
 - After that it will be your responsibility to manage \$sp
- By convention, stack grows towards lower addresses
 - to allocate new space (i.e., when you *push*), decrement \$sp
 - to free space on top of stack (*pop*), increment \$sp

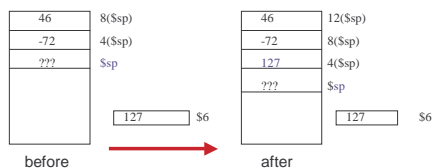
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4

Push operation

- push* adds an item on top of stack
 - one instruction to manipulate the data, e.g. "sw \$6,0(\$sp)"
 - one instruction to adjust the stack pointer e.g., "subu \$sp,\$sp,4"



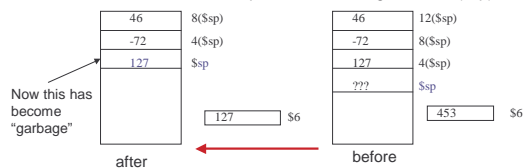
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5

Pop operation

- pop* removes the item on top of stack and stores it in a register
 - one instruction to adjust the stack pointer e.g., "addu \$sp,\$sp,4"
 - one instruction to manipulate the data, e.g. "lw \$6,0(\$sp)"



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6

Procedure call requirements (caller/callee)

- Caller must pass the return address to the callee
- Caller must pass the parameters to the callee
- Caller must save what is in *volatile* (registers) that could be used by the callee
- Callee must save the return address (in case it becomes a caller)
- Callee must provide (stack) storage for its own use
- Caller/callee should support recursive calls

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7

Mechanism

- Registers are used for
 - passing return address in \$ra
 - jal target
 - passing a small number of parameters (up to 4 in \$a0 to \$a3)
 - keeping track of the stack (\$sp)
 - returning function values (in \$v0 and \$v1)
- Stack is used for
 - saving registers to be used by callee
 - saving info about the caller (return address)
 - passing parameters if needed
 - allocating local data for the called procedure

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8

Procedure calls and register conventions

Register	Name	Function	Comment
\$0	Zero	Always 0	No-op on write
\$1	\$at	Reserved for assembler	Don't use it
\$23	\$v0-v1	Expr. Eval/funct. Return	
\$47	\$a0-a3	Proc./func. Call parameters	
\$8-15	\$t0-t7	Temporaries; volatile	Not saved on proc. Calls
\$16-23	\$s0-s7	Temporaries	Should be saved on calls
\$24-25	\$t8-t9	Temporaries; volatile	Not saved on proc. Calls
\$26-27	\$k0-k1	Reserved for O.S.	Don't use them
\$28	\$gp	Pointer to global static memory	
\$29	\$sp	Stack pointer	
\$30	\$fp	Frame pointer	
\$31	\$ra	Proc./func return address	

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9

Who does what on a call (one sample protocol)

- Caller
 - Saves any volatile register (\$t0-\$t9) having contents that need to be kept
 - Puts up to 4 arguments in \$a0-\$a3
 - If more than 4 arguments, pushes the rest on the stack
 - calls with jal instruction
- Callee
 - saves \$ra on stack
 - saves any non-volatile register (\$s0-s7) that it will use

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10

Who does what on return

- Callee
 - restores any non-volatile register (\$s0-\$s7) it has used
 - restores \$ra
 - puts function results in \$v0-\$v1
 - adjusts \$sp
 - returns to caller with "jr \$ra"
- Caller
 - restores any volatile register it had saved
 - examines \$v0-\$v1 if needed

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11

Example of a call sequence

- Assume 2 arguments in \$t0 and \$t3 and we want to save the contents of \$t6 and \$t7


```

move    $a0,$t0    #1st argument in $a0
move    $a1,$t3    #2nd argument in $a1
subu    $sp,$sp,8  #room for 2 temps on stack
sw      $t6,8($sp) #save $t6 on stack
sw      $t7,4($sp) #save $t7 on stack
jal     target
            
```
- Assume the callee does not need to save registers


```

target: sw  $ra,0($sp) #save return address
subu    $sp,$sp,4  # on stack
            
```

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12

Return from the previous sequence

- The callee will have put the function results in \$v0-\$v1

```
addu    $sp,$sp,4      #pop
lw      $ra,0($sp)     #return address in $ra
jr      $ra            #to caller
```

- The caller will restore \$t6 and \$t7 and adjust stack

```
lw      $t6,8($sp)
lw      $t7,4($sp)
addu    $sp,$sp,8
```