

MIPS Procedure Calls JVM and Assignment 3

CSE 378 – Section 3

Procedure Call Basics

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|--|--|
| <ol style="list-style-type: none"> 1 Jump to procedure: <ul style="list-style-type: none"> jal <label> - Saves return address to \$ra 1 Return from a procedure: <ul style="list-style-type: none"> jr \$ra 1 \$a0 - \$a3 to pass arguments 1 \$v0 and \$v1 to return values 1 Save certain registers to preserve across procedure calls. <ul style="list-style-type: none"> - Use the stack | <ol style="list-style-type: none"> 1 \$t0-\$t9, \$a0-a3, \$v0-v1 – <i>caller-saved</i>. <ul style="list-style-type: none"> - Caller's responsibility to save if expects to use these after a call. 1 \$s0-\$s7, \$ra, \$fp – <i>callee-saved</i>. <ul style="list-style-type: none"> - Callee's responsibility to save if callee uses them. - Save at beginning of function, restore at end |
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Calling procedures

To call a procedure:

1. Put arguments into \$a0-a3
2. Save caller-saved registers
3. jal <proc>
4. Restore caller-saved registers

Example:

```
<some stuff here, uses $t2>
# set up a call to myproc(4)
addi $a0, $0, 4
subu $sp, $sp, 4
sw $t2, 0($sp)
jal myproc
lw $t2, 0($sp)
addiu $sp, $sp, 4
<use $t2 again>
```

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Setup at the start/end of procedure

Before any procedure starts running, it must:

1. Allocate memory for callee-saved registers
2. Save callee-saved registers
 - 1 If calling another procedure inside, must save \$ra! (why?)

At the end of procedure:

1. Place return value into \$v0
2. Restore callee-saved regs
3. jr \$ra

```
myproc: # wants to use $s0 inside
subu $sp, $sp, 8
sw $ra, 4($sp)
sw $s0, 0($sp)
...
<do some computation in $s0>
...
addi $v0, $s0, 42
lw $s0, 0($sp)
lw $ra, 4($sp)
addiu $sp, $sp, 8
jr $ra
```

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Miscellaneous

- 1 MIPS stack conventions:
 - \$sp double-word aligned
 - Minimum frame size is 24 bytes (fits four arguments and return address)
 - Don't use it for projects
 - Other rules flexible too: have to use common sense for what you need to save
- 1 If >4 arguments, use the stack to pass them
 - Caller, callee must agree on where they go in the stack and who pops them off.

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JVM and Assignment 3

- 1 JVM is a stack machine
 - Portability
 - Compactness
- 1 Our simplified JVM consists of:
 - execution stack
 - 1 Instructions take parameters from the stack
 - 1 Instructions place results onto the stack
 - Pointer to top of the stack
 - local storage
 - 1 Just a big array for storing data
 - Java bytecode program
 - Program counter

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Emulating JVM

- 1 Interpreter:
 - Get next instruction
 - Decode it
 - Execute
 - Store results
 - Repeat

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Emulating JVM

- 1 Probably need SPIM registers for:
 - Pointer to top of JVM stack
 - Pointer to current JVM instruction (PC)
 - Holding a couple of values from the stack (when pushing/popping) - v1, v2
- 1 Use SPIM static data section for:
 - The entire execution stack (1024 bytes maximum)
 - the local storage area
 - The program itself
 - 1 sequence of instruction opcodes and parameters

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JVM Instructions

- 1 Load constant (BIPUSH for 8-bit, SIPUSH for 16)



- 1 Arithmetic (IADD, ISUB, IMUL, IDIV)



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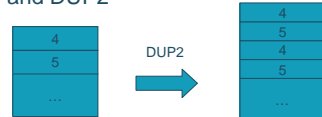
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JVM Instructions 2

- 1 POP



- 1 DUP and DUP2



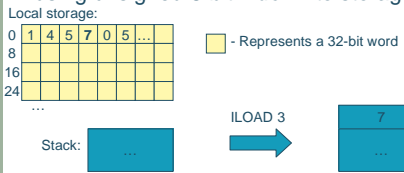
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Loading from local storage

- 1 ILOAD, ISTORE - load/store 32-bit word using unsigned 8-bit index into storage



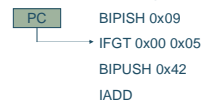
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Branches

- 1 Pop one thing off stack, compare with zero using specified condition, update PC if true
- 1 Take a signed 2-byte offset from current PC
 - No "labels" in bytecodes, just offsets



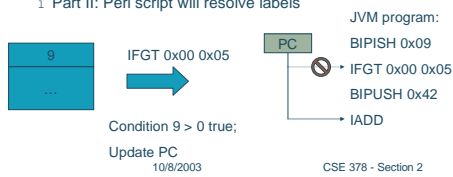
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Branches

- To understand offset destinations, add up opcodes (1 byte), along with any arguments
 - E.g. IFEQ 0x00 0x05 takes 3 bytes, IADD takes 1.



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Example: a=a+b+c

Add first 3 words in local storage
Store the result into the first local storage word

```

ILOAD 0
ILOAD 1
ILOAD 2
IADD
IADD
ISTORE 0
    
```

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Example 2: if (b==0) a=3; else a=5;

- Assume a is local word 0,
b is local word 1:

<pre> ILOAD 1 IFEQ skip BIPUSH 3 ISTORE 0 GOTO endif skip: BIPUSH 5 ISTORE 0 endif: ... </pre>	<p>To bytecodes (use perl script)</p> <p>→</p>	<pre> -align 2 test2: .byte 0x15, 0x01 # iload .byte 0x39, 0x00, 0x0a # ifeq .byte 0x10, 0x03 # bipush .byte 0x36, 0x00 # istore .byte 0xa7, 0x00, 0x07 # goto .byte 0x10, 0x05 # bipush .byte 0x36, 0x00 # istore .byte 0x00 # -align 2 end_test2: </pre>
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