CSE 401 – Compilers

LR Parsing

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Agenda

- LR Parsing
- Table-driven Parsers
- Parser States
- Shift-Reduce and Reduce-Reduce conflicts

LR(1) Parsing

- We'll look at LR(1) parsers
  - Left to right scan, Rightmost derivation, 1 symbol lookahead
  - Almost all practical programming languages have an LR(1) grammar
  - LALR(1), SLR(1), etc. – subsets of LR(1)
    - LALR(1) can parse most real languages, is more compact, and is used by YACC/Bison/CUP/etc.

Bottom-Up Parsing

- Idea: Read the input left to right
- Whenever we've matched the right hand side of a production, reduce it to the appropriate non-terminal and add that non-terminal to the parse tree
- The upper edge of this partial parse tree is known as the *frontier*

Example

- Grammar

  S ::= a*AB* e

  A ::= *A*bc | b

  B ::= d

- Bottom-up Parse

  a b b c d e

Details

- The bottom-up parser reconstructs a reverse rightmost derivation
- Given the rightmost derivation

  $S => \beta_1 => \beta_2 => \ldots => \beta_{n-2} => \beta_{n-1} => \beta_n = w$

  the parser will first discover $\beta_{n-1} => \beta_n$, then $\beta_{n-2} => \beta_{n-1}$, etc.
- Parsing terminates when
  - β_1 reduced to S (start symbol, success), or
  - No match can be found (syntax error)

How Do We Parse with This?

- Key: given what we've already seen and the next input symbol, decide what to do.
- Choices:
  - Perform a reduction
  - Look ahead further
- Can reduce $A => \beta$ if both of these hold:
  - $A => \beta$ is a valid production
  - $A => \beta$ is a step in *this* rightmost derivation
- This is known as a *shift-reduce* parser

Sentential Forms

- If $S =>^* \alpha$, the string α is called a *sentential form* of the of the grammar
- In the derivation
  $S => \beta_1 => \beta_2 => ... => \beta_{n-2} => \beta_{n-1} => \beta_n = w$
  each of the β_i are sentential forms
- A sentential form in a rightmost derivation is called a right-sentential form (similarly for leftmost and left-sentential)

Handles

- Informally, a substring of the tree frontier that matches the right side of a production
  - Even if $A ::= \beta$ is a production, β is a handle only if it matches the frontier at a point where $A ::= \beta$ was used in that derivation
  - β may appear in many other places in the frontier without being a handle for that particular production

Handles (cont.)

- Formally, a *handle* of a right-sentential form γ is a production $A ::= \beta$ and a position in γ where β may be replaced by A to produce the previous right-sentential form in the rightmost derivation of γ

Handle Examples

- In the derivation

  S => a*AB*e => a*A*de => a*A*bcde => abbcde

  - abbcde is a right sentential form whose handle is *A*::=b at position 2
  - a*A*bcde is a right sentential form whose handle is *A*::=*A*bc at position 4
    - Note: some books take the left of the match as the position

Implementing Shift-Reduce Parsers

- Key Data structures
  - A stack holding the frontier of the tree
  - A string with the remaining input

Shift-Reduce Parser Operations

- *Reduce* – if the top of the stack is the right side of a handle $A ::= \beta$, pop the right side β and push the left side A.
- *Shift* – push the next input symbol onto the stack
- *Accept* – announce success
- *Error* – syntax error discovered

S ::= a*AB*e
A ::= *A*bc | b
B ::= d

Shift-Reduce Example

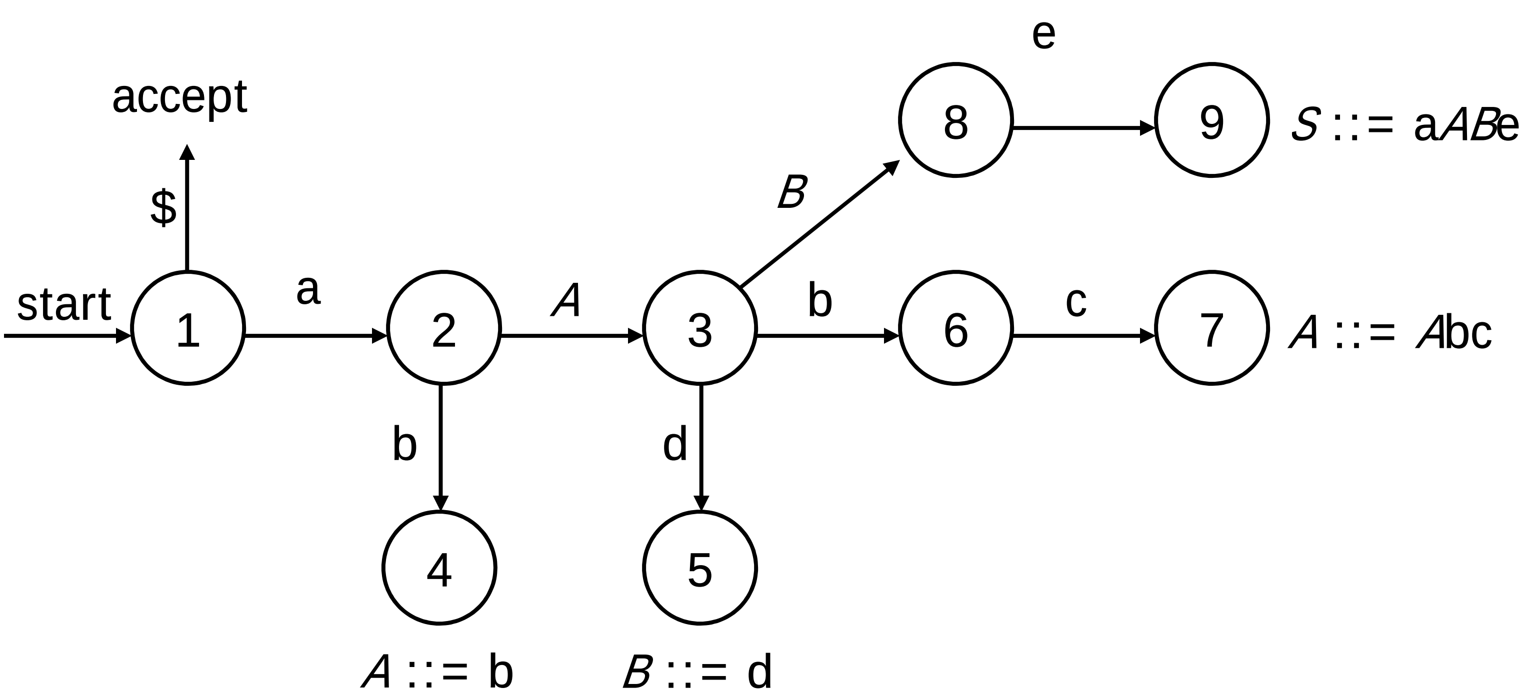
Stack	Input	Action
$	abbcde$	*shift*

How Do We Automate This?

- Def. *Viable prefix* – a prefix of a right-sentential form that can appear on the stack of the shift-reduce parser
  - Equivalent: a prefix of a right-sentential form that does not continue past the rightmost handle of that sentential form
- Idea: Construct a DFA to recognize viable prefixes given the stack and remaining input
  - Perform reductions when we recognize them

DFA for prefixes of

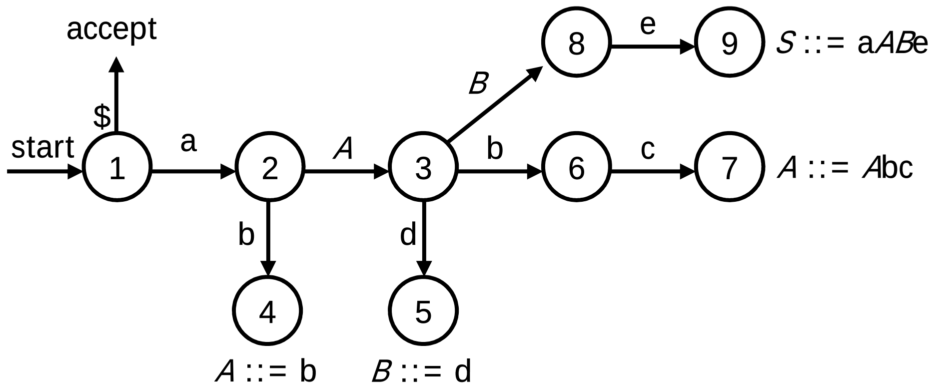
S ::= a*AB*e
A ::= *A*bc | b
B ::= d



Trace

$S ::= aABe$
$A ::= Abc \mid b$
$B ::= d$

Stack	Input
$	abbcde$



Observations

- Way too much backtracking
  - We want the parser to run in time proportional to the length of the input
- Where the heck did this DFA come from anyway?
  - From the underlying grammar
  - We'll defer construction details for now

Avoiding DFA Rescanning

- Observation: after a reduction, the contents of the stack are the same as before except for the new non-terminal on top
  - ∴ Scanning the stack will take us through the same transitions as before until the last one
  - ∴ If we record state numbers on the stack, we can go directly to the appropriate state when we pop the right hand side of a production from the stack

Stack

- Change the stack to contain pairs of states and symbols from the grammar

  $\$s_0\ X_1\ s_1\ X_2\ s_2\ \ldots\ X_n\ s_n$
  - State s_0 represents the accept state
    - (Not always added – depends on particular presentation)

- Observation: in an actual parser, only the state numbers need to be pushed, since they implicitly contain the symbol information, but for explanations, it's clearer to use both.

Encoding the DFA in a Table

- A shift-reduce parser's DFA can be encoded in two tables
  - One row for each state
  - *action* table encodes what to do given the current state and the next input symbol
  - *goto* table encodes the transitions to take after a reduction

Actions (1)

- Given the current state and input symbol, the main possible actions are
  - s*i* – shift the input symbol and state *i* onto the stack (i.e., shift and move to state *i*)
  - r*j* – reduce using grammar production *j*
    - The production number tells us how many <symbol, state> pairs to pop off the stack

Actions (2)

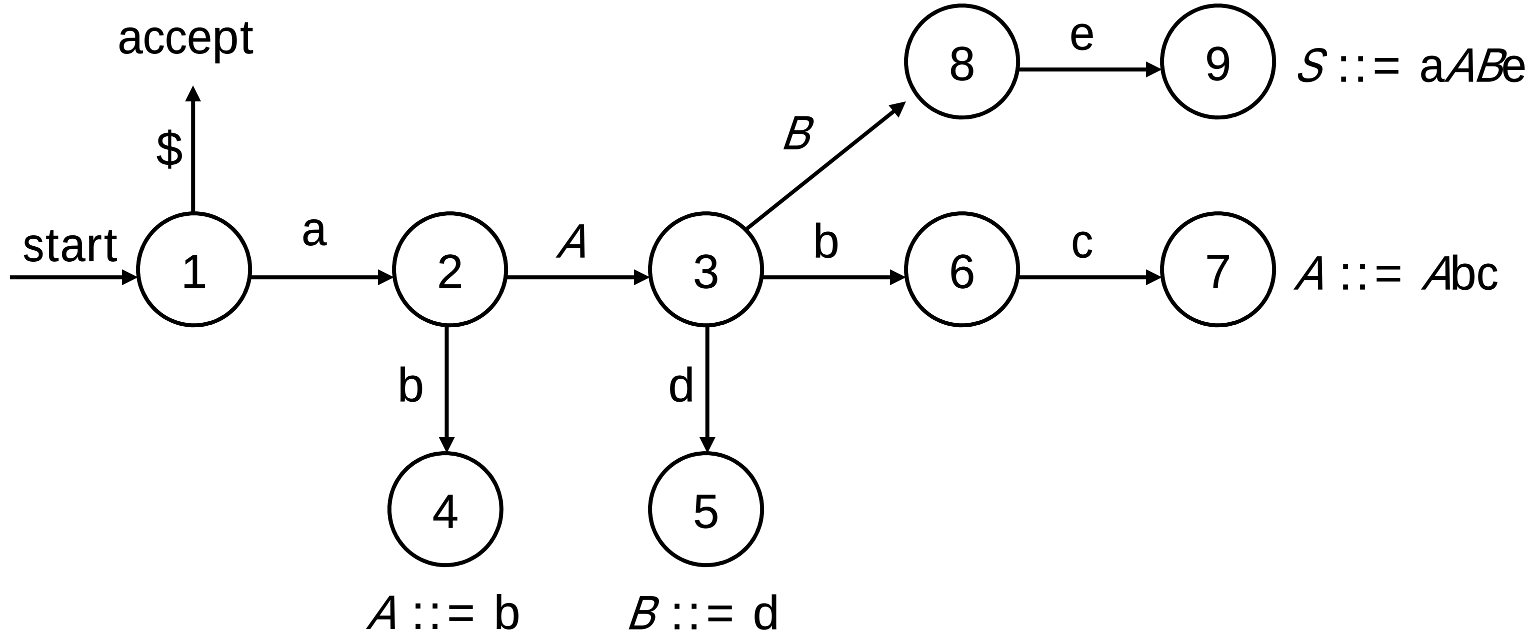
- Other possible *action* table entries
  - *accept*
  - blank – no transition – syntax error
    - A LR parser will detect an error as soon as possible on a left-to-right scan
    - A real compiler needs to produce an error message, recover, and continue parsing when this happens

Goto

- When a reduction is performed, <symbol, state> pairs are popped from the stack revealing a state *uncovered_s* on the top of the stack
- goto[*uncovered_s* , *A*] is the new state to push on the stack when reducing production *A* ::= β (after popping β and finding state *uncovered_s* on top)

Reminder: DFA for

S ::= a*AB*e
A ::= *A*bc | b
B ::= d



LR Parse Table for

1. *S* ::= a*ABe*
2. *A* ::= *A*bc
3. *A* ::= b
4. *B* ::= d

State	*action*						*goto*		
	a	b	c	d	e	$	A	B	S
1	s2					acc			g1
2		s4					g3		
3		s6		s5				g8	
4	r3	r3	r3	r3	r3	r3			
5	r4	r4	r4	r4	r4	r4			
6			s7						
7	r2	r2	r2	r2	r2	r2			
8					s9				
9	r1	r1	r1	r1	r1	r1			

LR Parsing Algorithm (1)

```
word = scanner.getToken();
while (true) {
    s = top of stack;
    if (action[s, word] = si ) {
      push word; push i (state);
      word = scanner.getToken();
    } else if (action[s, word] = rj ) {
      pop 2 * length of right side of
          production j (2*|β|);
      uncovered_s = top of stack;
      push left side A of production j ;
      push state goto[uncovered_s, A];
    }
} else if (action[s, word] = accept ) {
      return;
} else {
      // no entry in action table
      report syntax error;
      halt or attempt recovery;
}
```

1. *S* ::= a*AB*e
2. *A* ::= *A*bc
3. *A* ::= b
4. *B* ::= d

Example

Stack

$

Input

abbcde$

S	action						goto		
	a	b	c	d	e	$	A	B	S
1	s2					ac			g1
2		s4					g3		
3		s6		s5				g8	
4	r3	r3	r3	r3	r3	r3			
5	r4	r4	r4	r4	r4	r4			
6			s7						
7	r2	r2	r2	r2	r2	r2			
8					s9				
9	r1	r1	r1	r1	r1	r1			

LR States

- Idea is that each state encodes
  - The set of all possible productions that we could be looking at, given the current state of the parse, and
  - *Where* we are in the right hand side of each of those productions

Items

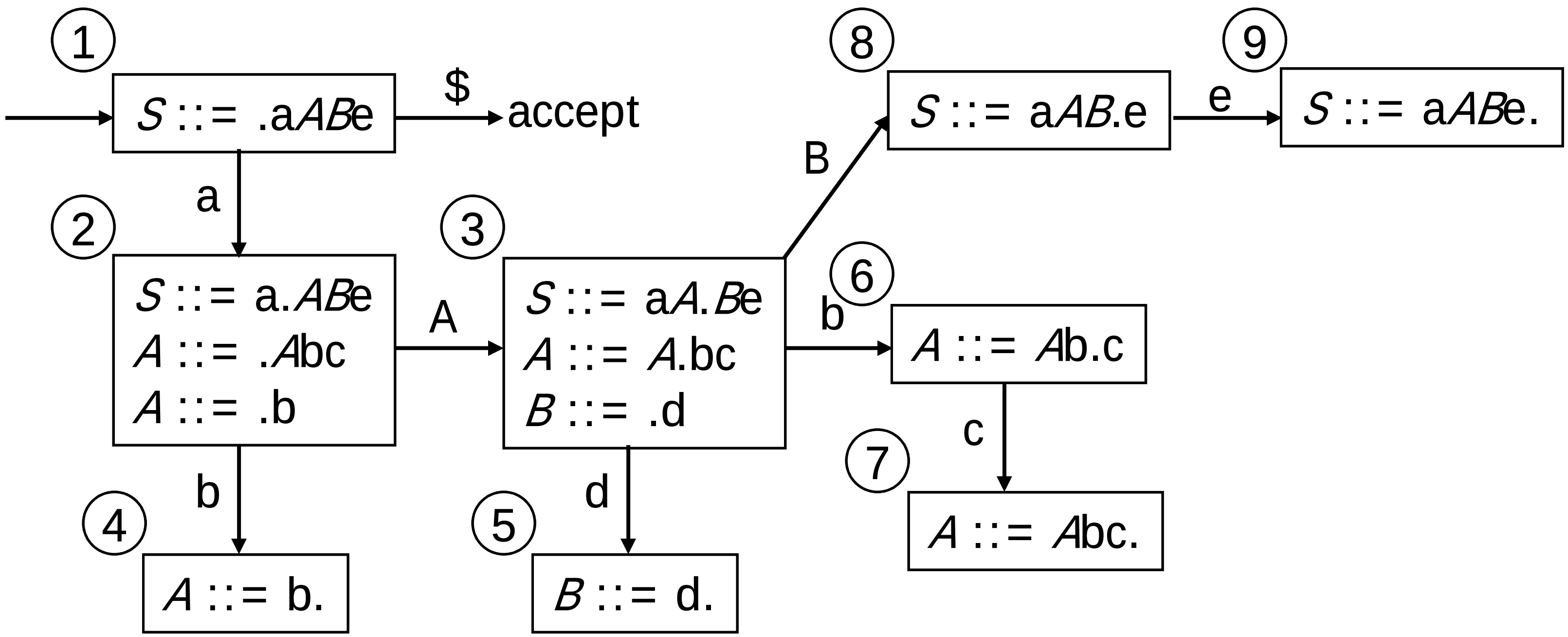
- An *item* is a production with a dot in the right hand side
- Example: Items for production $A ::= XY$

  $A ::= .XY$

  $A ::= X.Y$

  $A ::= XY.$
- Idea: The dot represents a position in the production

DFA for $S ::= aABe$ $A ::= Abc \mid b$ $B ::= d$



Problems with Grammars

- Grammars can cause problems when constructing a LR parser
  - Shift-reduce conflicts
  - Reduce-reduce conflicts

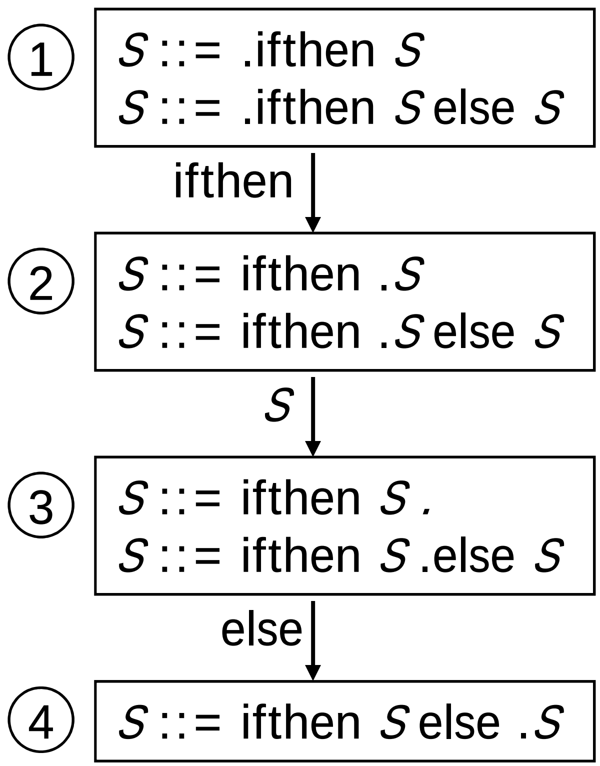
Shift-Reduce Conflicts

- Situation: both a shift and a reduce are possible at a given point in the parse (equivalently: in a particular state of the DFA)
- Classic example: if-else statement

  S ::= ifthen *S* | ifthen *S* else *S*

Parser States for

1. S ::= ifthen S
2. S ::= ifthen S else S



- State 3 has a shift-reduce conflict
  - Can shift past else into state 4 (s4)
  - Can reduce (r1)

    S ::= ifthen S

(Note: other S ::= .ifthen items not included in states 2-4 to save space)

Solving Shift-Reduce Conflicts

- Fix the grammar
  - Done in Java reference grammar, others
- Use a parse tool with a "longest match" rule – i.e., if there is a conflict, choose to shift instead of reduce
  - Does exactly what we want for if-else case
  - Guideline: a few shift-reduce conflicts are fine, but be sure they do what you want

Reduce-Reduce Conflicts

- Situation: two different reductions are possible in a given state
- Contrived example

  $S ::= A$

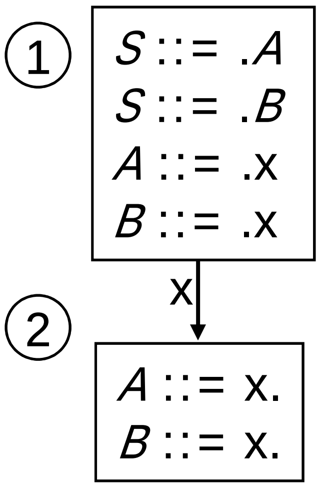
  $S ::= B$

  $A ::= \mathrm{x}$

  $B ::= \mathrm{x}$

Parser States for

1. $S ::= A$
2. $S ::= B$
3. $A ::= x$
4. $B ::= x$



- State 2 has a reduce-reduce conflict (r3, r4)

Handling Reduce-Reduce Conflicts

- These normally indicate a serious problem with the grammar.
- Fixes
  - Use a different kind of parser generator that takes lookahead information into account when constructing the states (LR(1) instead of SLR(1) for example)
    - Most practical tools use this information
  - Fix the grammar

Another Reduce-Reduce Conflict

- Suppose the grammar separates arithmetic and boolean expressions

  expr ::= *aexp* | *bexp*

  aexp ::= *aexp* * *aident* | *aident*

  bexp ::= *bexp* && *bident* | *bident*

  aident ::= *id*

  bident ::= *id*

- This will create a reduce-reduce conflict

Covering Grammars

- A solution is to merge *aident* and *bident* into a single non-terminal (or use *id* in place of *aident* and *bident* everywhere they appear)
- This is a *covering grammar*
  - Includes some programs that are not generated by the original grammar
  - Use the type checker or other static semantic analysis to weed out illegal programs later

Coming Attractions

- Constructing LR tables
  - We'll present a simple version (SLR(0)) in lecture, then talk about extending it to LR(1)
- LL parsers and recursive descent
- Continue reading ch. 3