#### **Procedures**

# CSE 413, Autumn 2002 **Programming Languages**

http://www.cs.washington.edu/education/courses/413/02au/

2-October-2002

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#### **Combinations**

- (operator operand operand)
- There are numerous pre-defined operators
- We can define our own, arbitrarily complex operators (functions, procedures) as well
- This is a key capability by which we can operate at higher levels of abstraction

# Readings and References

#### Reading

» Sections 1.1.6-1.1.8, Structure and Interpretation of Computer Programs, by Abelson, Sussman, and Sussman

#### Other References

» Section 4.1, Revised Report on the Algorithmic Language Scheme (R5RS)

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# Define and name a procedure

- (define ( $\langle name \rangle \langle formal \ params \rangle$ )  $\langle body \rangle$ )
  - » define special form
  - » name the name that the procedure is bound to
  - » formal params names used within the body of procedure
  - » body expression (or sequence of expressions) that will be evaluated when the procedure is called.
  - » The result of the last expression in the body will be returned as the result of the procedure call

#### Example definitions

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#### Defined procedures are "first class"

- Compound procedures that we define are used exactly the same way the primitive procedures provided in Scheme are used
  - » names of built-in procedures are not treated specially; they are simply names that have been pre-defined
  - » you can't tell whether a name stands for a primitive (built-in) procedure or a compound (defined) procedure by looking at the name or how it is used

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# Evaluation example

- (area-of-ring 4 1)
  - » evaluate operator area-of-ring => procedure
    definition
  - $\rightarrow$  evaluate 4 => 4
  - $\rightarrow$  evaluate 1 => 1
  - » apply the procedure to the arguments

#### **Booleans**

- Recall that one type of data object is boolean
  - » #t (true) or #f (false)
- We can use these explicitly or by calculating them in expressions that yield boolean values
- An expression that yields a true or false value is called a predicate
  - » #t => #t
  - > (< 5 5) => #f
  - » (> pi 0) => #t

# Conditional expressions

- As in all languages, we need to be able to make decisions based on inputs and do something depending on the result
- A predicate expression is evaluated » true or false
- The consequent expression is evaluated if the predicate is true

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#### Special form: cond

- (cond  $\langle clause_1 \rangle \langle clause_2 \rangle$  ...  $\langle clause_n \rangle$ )
- each clause is of the form
  - » (⟨predicate⟩ ⟨expression⟩)
  - » where  $\langle predicate \rangle$  is a boolean expression and  $\langle expression \rangle$ is the consequent expression to execute if (predicate) is true
- the last clause can be of the form
  - » (else \( \langle expression \rangle )
  - » in which case (expression) is executed if none of the preceding (predicates) were true

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# Example: sign.scm

```
; return the sign of x as -1, 0, or 1
(define (sign x)
  (cond
     ((< x 0) -1)
     ((= \mathbf{x} \ 0) \ 0)
     ((> x 0) +1))
```

# Special form: if

- (if \( \text{predicate} \) \( \text{consequent} \) \( \text{alternate} \)
- (if \( \rho predicate \) \( \consequent \) )
- *(predicate)* is a boolean expression
- *⟨consequent⟩* is the expression to execute if *⟨predicate⟩* is true
- *(alternate)* is the expression to execute if *(predicate)* is false

#### Examples: abs.scm, true-false.scm

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# Logical composition

- (and  $\langle e_1 \rangle \langle e_2 \rangle ... \langle e_n \rangle$ )
- (or  $\langle e_1 \rangle \langle e_2 \rangle ... \langle e_n \rangle$ )
- (not  $\langle e \rangle$ )
- Scheme interprets the expressions  $e_i$  one at a time in left-to-right order until it can tell the correct answer

» ie, these are short-circuit operators

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#### in-range.scm

#### Newton's method for square root

- Guess a value y for the square root of x
- Is it close enough to the desired value  $\sqrt[2]{x}$ ? » ie, is  $y^2$  close to x?
- If yes, then done. Return recent guess.
- If no, then new guess is average of current guess and  $\frac{x}{guess}$
- Repeat with new guess

#### sqrta.scm

```
; Square root using Newton's method

(define (average a b)
   (/ (+ a b) 2.0))

(define (good-enough? guess x)
   (< (abs (- (* guess guess) x)) 0.001))

(define (improve guess x)
   (average guess (/ x guess)))

(define (sqrt-iter guess x)
   (if (good-enough? guess x)
        guess
        (sqrt-iter (improve guess x) x )))

(define (sqrta x)
   (sqrt-iter 1.0 x))</pre>
```

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# auxiliary functions

```
; Square root using Newton's method

(define (average a b)
   (/ (+ a b) 2.0))

(define (good-enough? guess x)
   (< (abs (- (* guess guess) x)) 0.001))

(define (improve guess x)
   (average guess (/ x guess)))</pre>
```

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#### iterator and main functions

```
(define (sqrt-iter guess x)
  (if (good-enough? guess x)
        guess
        (sqrt-iter (improve guess x) x )))
(define (sqrta x)
  (sqrt-iter 1.0 x))
```

#### sqrt-iter

- Our first example of recursion
- Note that this recursion is used to implement a loop (an iteration)
  - » We will see this over and over in Scheme
- Iteration is calling the same block of code with a changing set of parameters
- The syntax of the procedure is recursive but the resulting process is iterative
  - » more on this next lecture