Lecture08

CSE 421
Introduction to Algorithms

Winter 2024

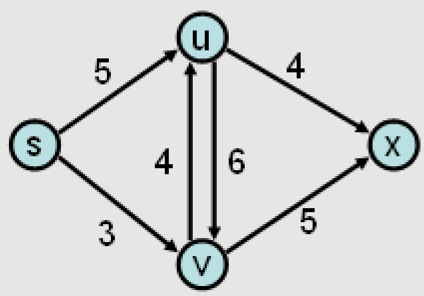
Lecture 8

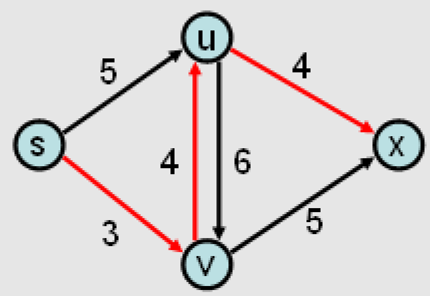
Minimum Spanning Trees

1

Bottleneck Shortest Path

- Define the bottleneck distance for a path P, $Len_B(P)$ to be the maximum cost edge along the path





$Len_B(x,y) = \text{Min} \{P \text{ from } x \text{ to } y \mid Len_B(P)\}$

Dijkstra's Algorithm for Bottleneck Shortest Paths

S = { }; d[s] = negative infinity; d[v] = infinity for v != s

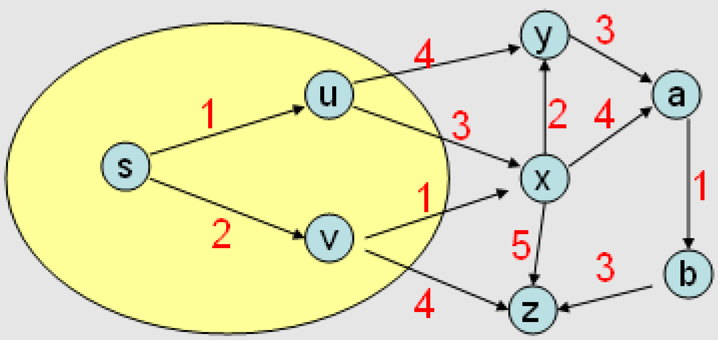
While S != V

Choose v in V-S with minimum d[v]

Add v to S

For each w in the neighborhood of v

d[w] = min(d[w], max(d[v], c(v, w)))



3

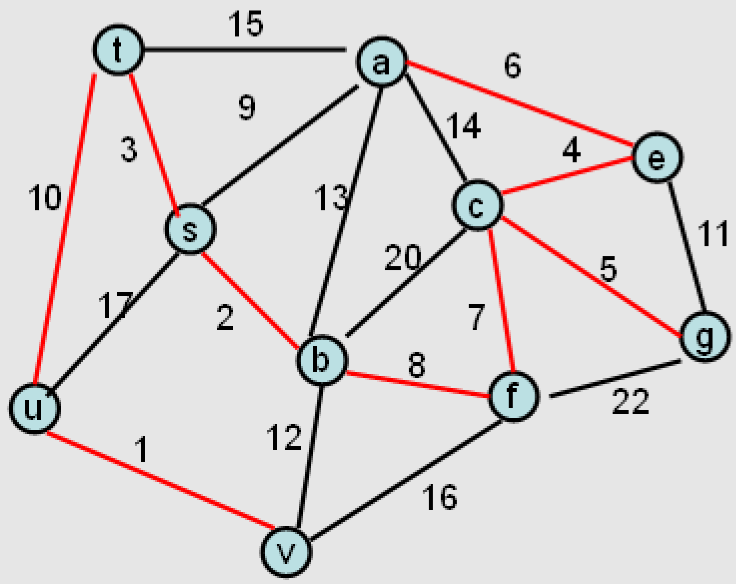
Minimum Spanning Tree

- Introduce Problem
- Demonstrate Prim's and Kruskal's algorithms
- Provide proofs that the algorithms work

4

Minimum Spanning Tree

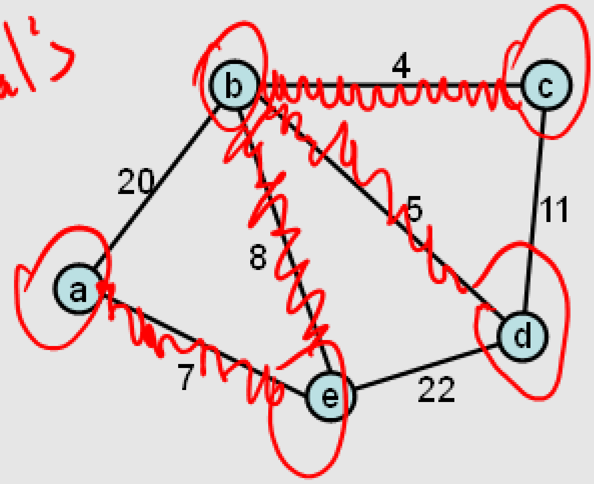
Undirected, Connected



5

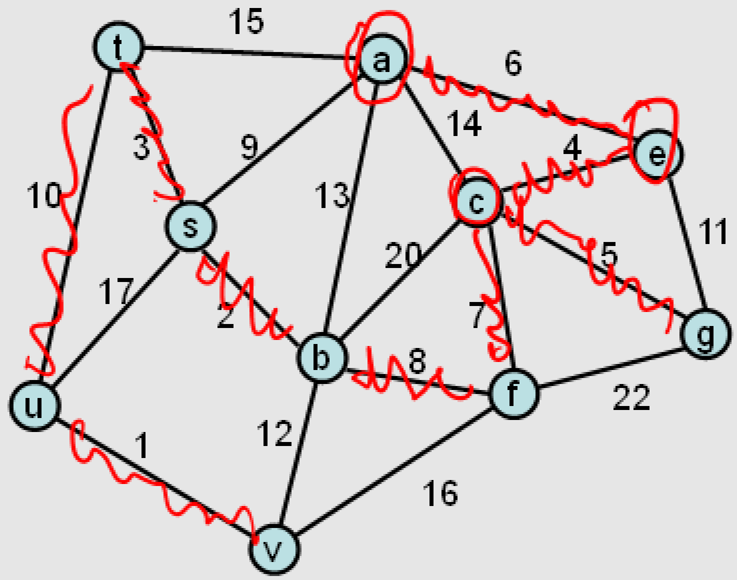
Greedy Algorithms for Minimum Spanning Tree

- Extend a tree by including the cheapest out going edge – Prim's
- Add the cheapest edge that joins disjoint components – Kruskal's
- ~~Delete the most expensive edge that does not disconnect the graph~~



Greedy Algorithm 1
Prim's Algorithm

- Extend a tree by including the cheapest out going edge



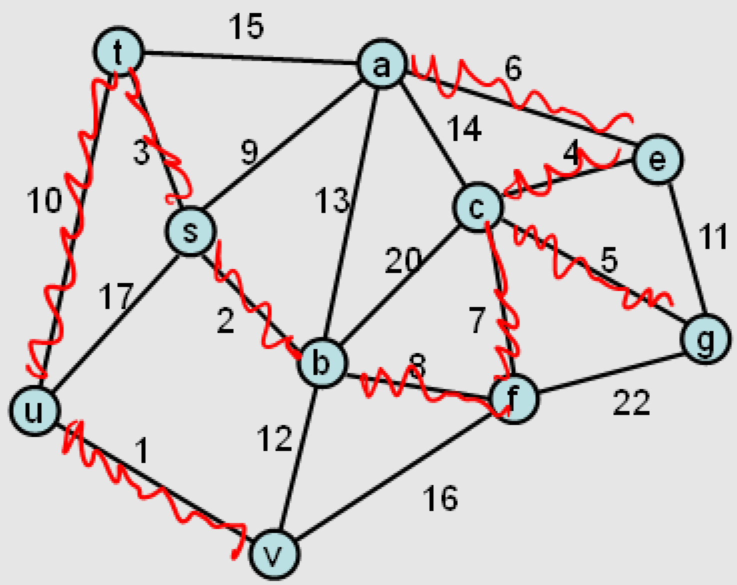
Construct the MST with Prim's algorithm starting from vertex a

Label the edges in order of insertion

7

Greedy Algorithm 2
Kruskal's Algorithm

- Add the cheapest edge that joins disjoint components



Construct the MST with Kruskal's algorithm

Label the edges in order of insertion

Dijkstra's Algorithm for Minimum Spanning Trees

S = { }; d[s] = 0; d[v] = infinity for v != s

While S != V

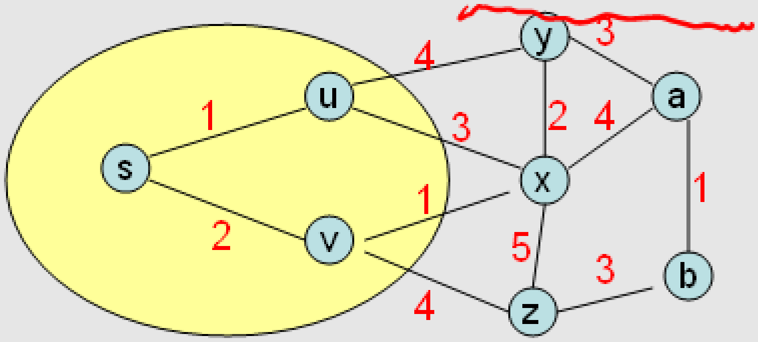
    Choose v in V-S with minimum d[v]

    Add v to S

    For each w in the neighborhood of v

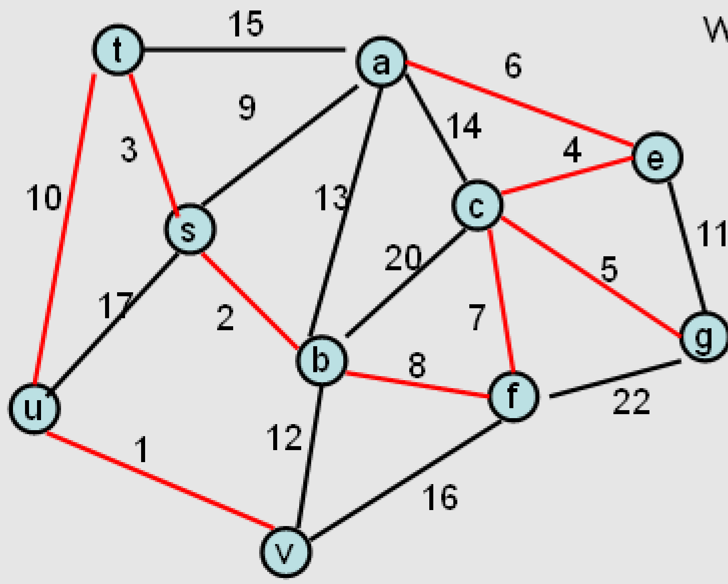
        d[w] = min(d[w], c(v, w))

m log n



9

Minimum Spanning Tree

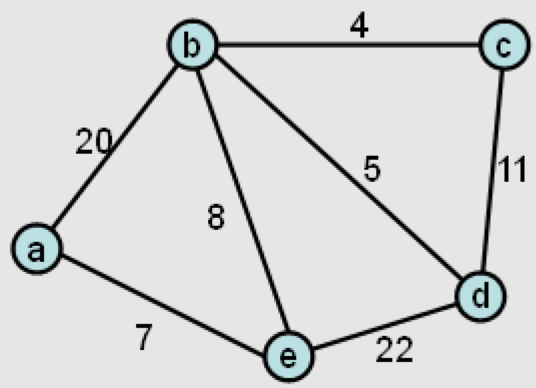


Undirected Graph G=(V,E) with edge weights

10

Greedy Algorithms for Minimum Spanning Tree

- [Prim] Extend a tree by including the cheapest out going edge
- [Kruskal] Add the cheapest edge that joins disjoint components



11

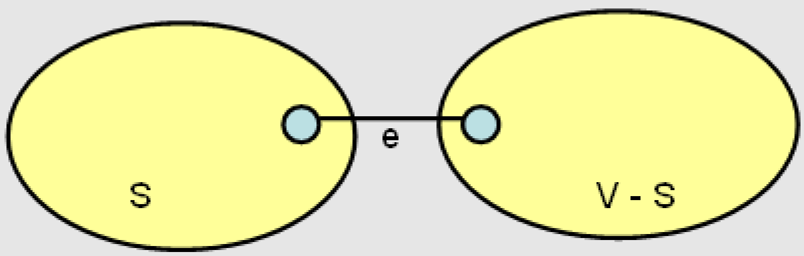
Why do the greedy algorithms work?

- For simplicity, assume all edge costs are distinct

12

Edge inclusion lemma

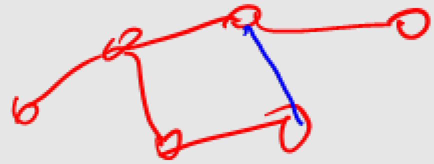
- Let S be a subset of V, and suppose e = (u, v) is the minimum cost edge of E, with u in S and v in V-S
- e is in every minimum spanning tree of G
  - Or equivalently, if e is not in T, then T is not a minimum spanning tree



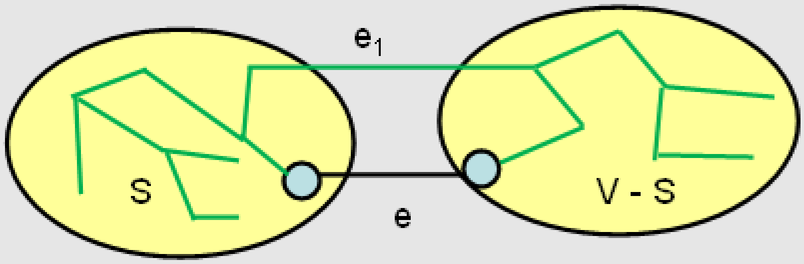
13

e is the minimum cost edge between S and V-S

Proof



- Suppose T is a spanning tree that does not contain e
- Add e to T, this creates a cycle
- The cycle must have some edge $e_1 = (u_1, v_1)$ with u_1 in S and v_1 in V-S



- $T_1 = T - \{e_1\} + \{e\}$ is a spanning tree with lower cost
- Hence, T is not a minimum spanning tree

14

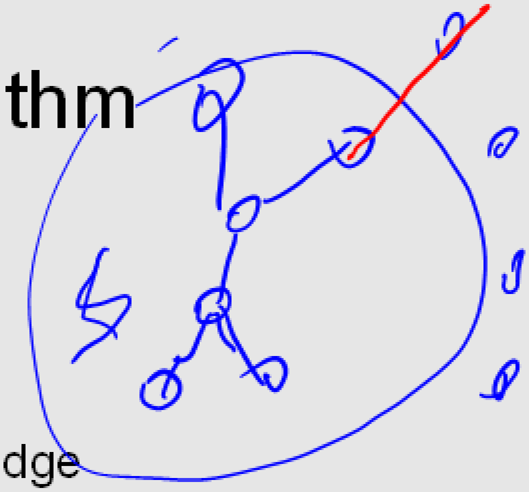
Optimality Proofs

- Prim's Algorithm computes a MST
- Kruskal's Algorithm computes a MST

- Show that when an edge is added to the MST by Prim or Kruskal, the edge is the minimum cost edge between S and V-S for some set S.

15

Prim's Algorithm



S = { }; T = { };

while S != V

choose the minimum cost edge
e = (u,v), with u in S, and v in V-S

add e to T

add v to S

e is in MST

16

Prove Prim's algorithm computes an MST

- Show an edge e is in the MST when it is added to T

17

Kruskal's Algorithm

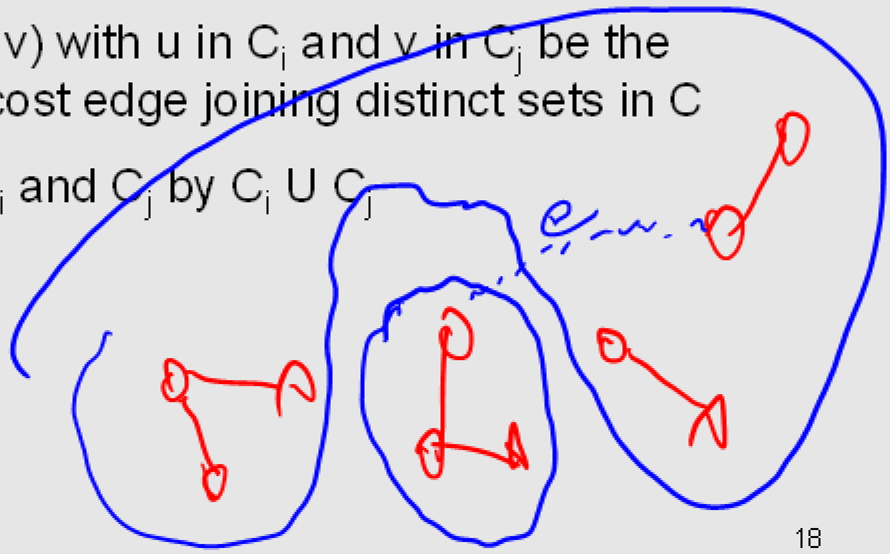
Let C = {{v_1}, {v_2}, . . ., {v_n}}; T = { }

while |C| > 1

Let e = (u, v) with u in C_i and v in C_j be the minimum cost edge joining distinct sets in C

Replace C_i and C_j by C_i U C_j

Add e to T



18

Prove Kruskal's algorithm computes an MST

- Show an edge e is in the MST when it is added to T

19

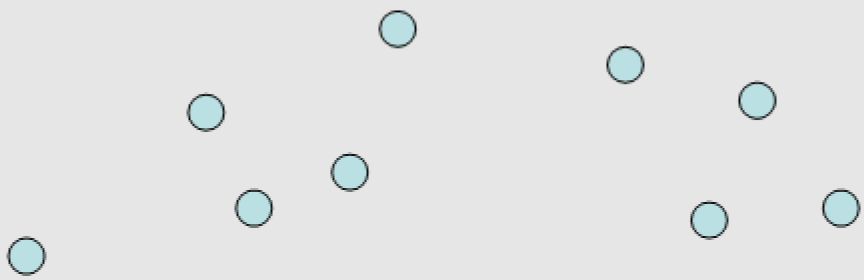
Dealing with the assumption of no equal weight edges

- Force the edge weights to be distinct
  - Give a tie breaking rule for equal weight edges

20

Application: Clustering

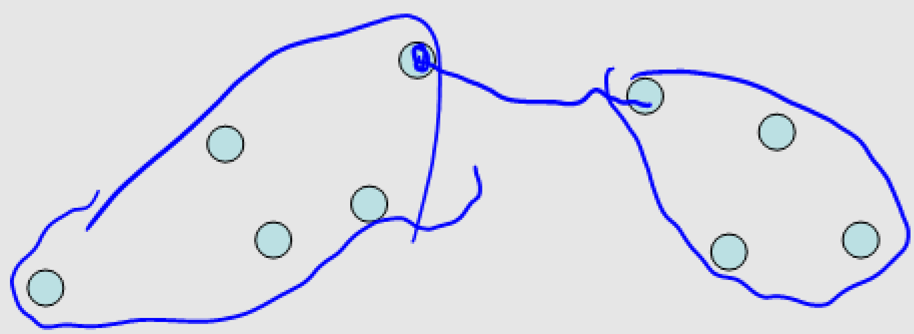
- Given a collection of points in an r-dimensional space and an integer K, divide the points into K sets that are closest together



21

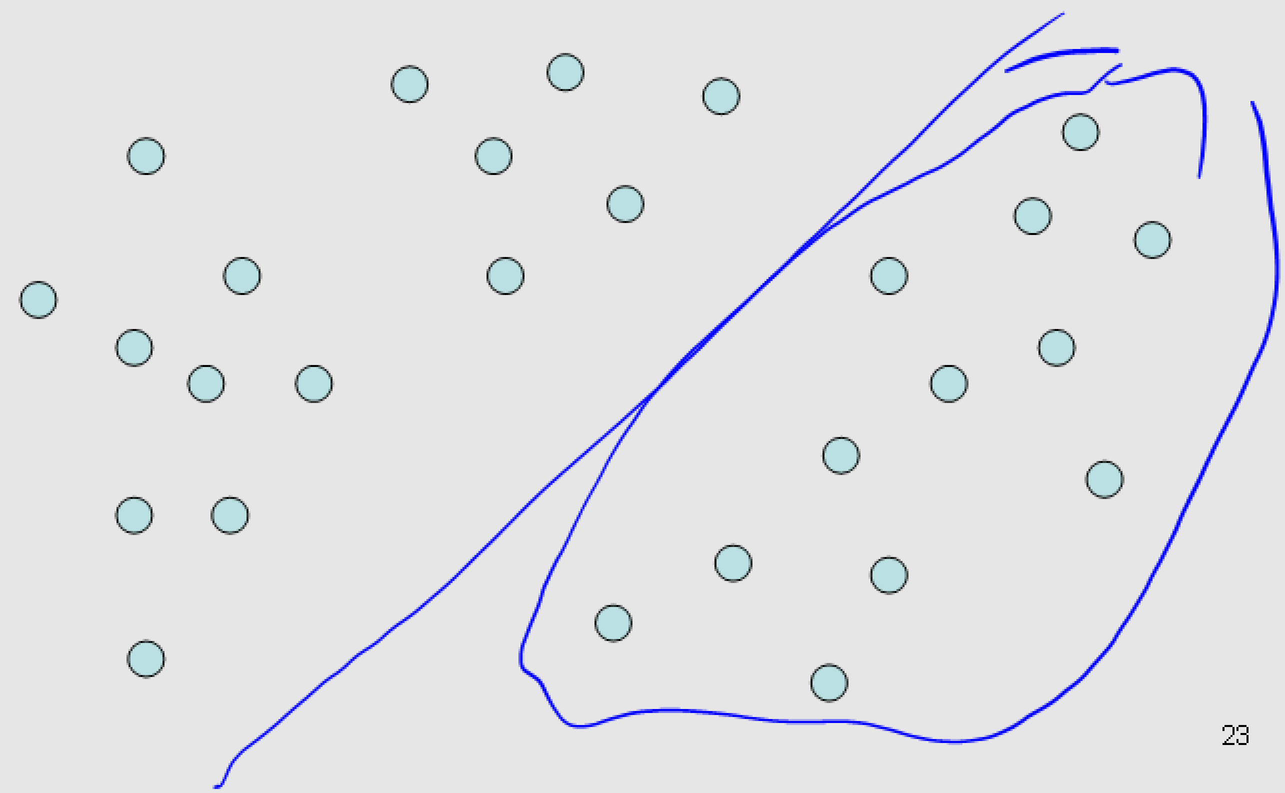
Distance clustering

- Divide the data set into K subsets to maximize the distance between any pair of sets
  - dist $(S_1, S_2) = \min \{dist(x, y) \mid x \text{ in } S_1, y \text{ in } S_2\}$



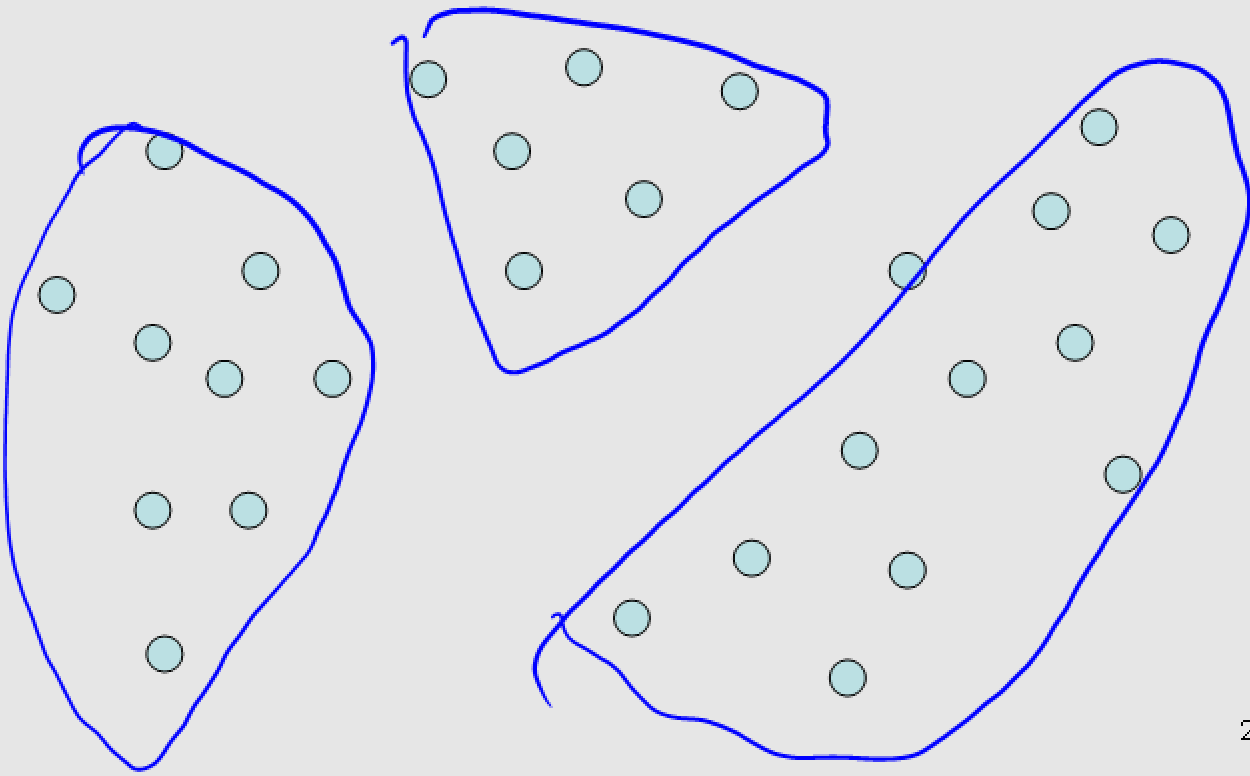
22

Divide into 2 clusters



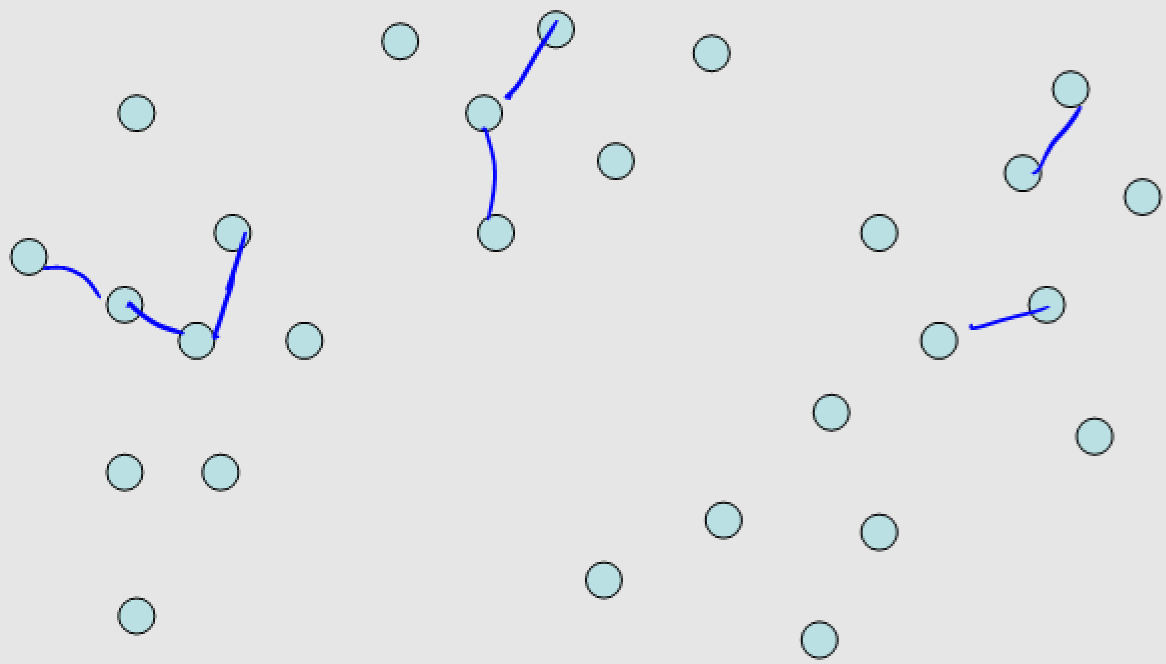
23

Divide into 3 clusters



24

Divide into 4 clusters



25

Distance Clustering Algorithm

Let C = $\{\{v_1\}, \{v_2\}, \ldots, \{v_n\}\}$; T = { }

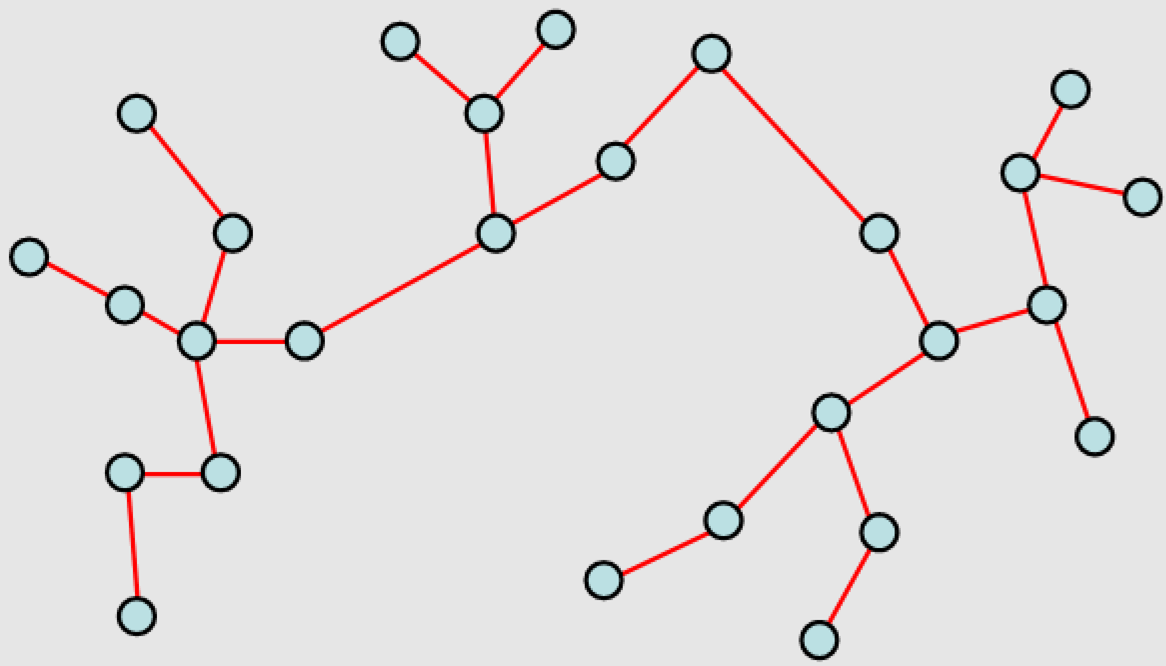
while |C| > K

Let e = (u, v) with u in C_i and v in C_j be the minimum cost edge joining distinct sets in C

Replace C_i and C_j by $C_i \cup C_j$

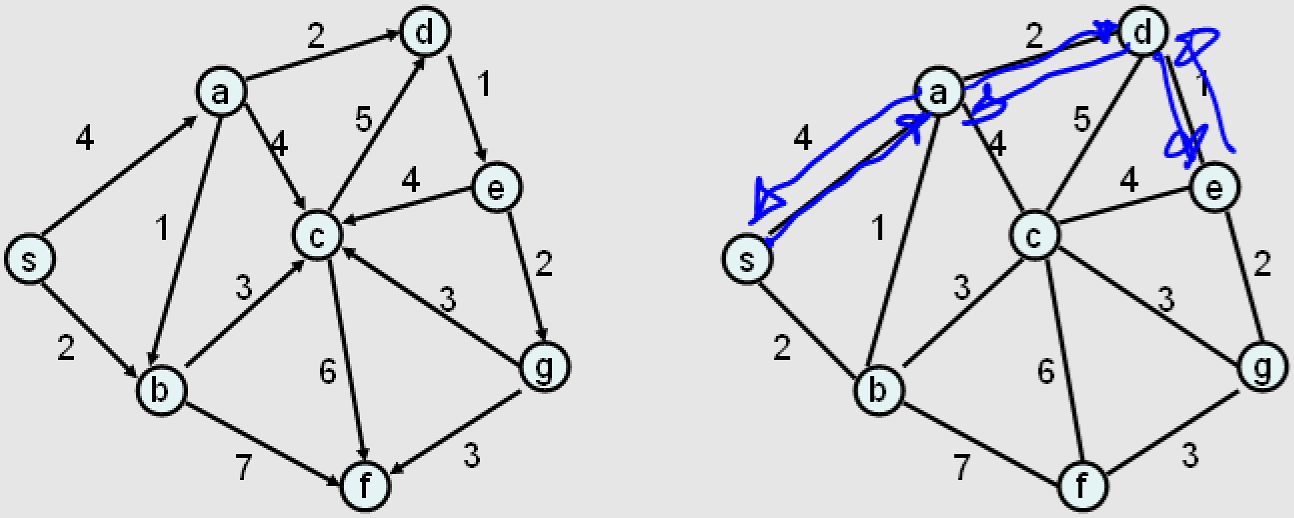
26

K-clustering



27

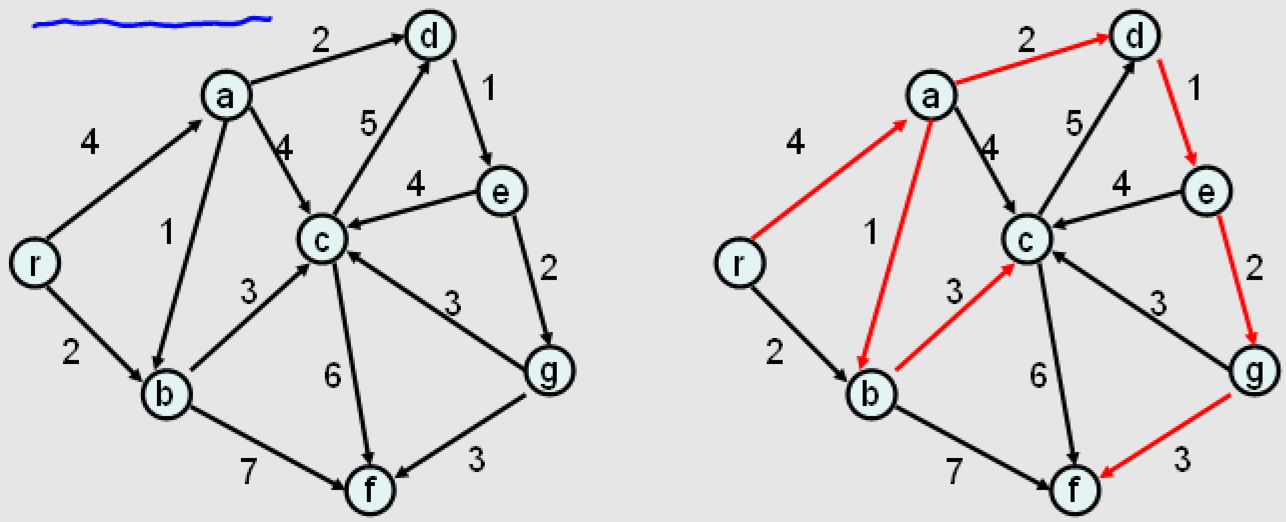
Shortest paths in directed graphs vs undirected graphs



28

What about the minimum spanning tree of a directed graph?

- Must specify the root r
- Branching: Out tree with root r

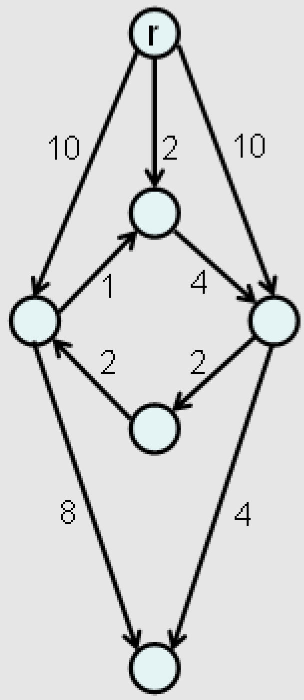


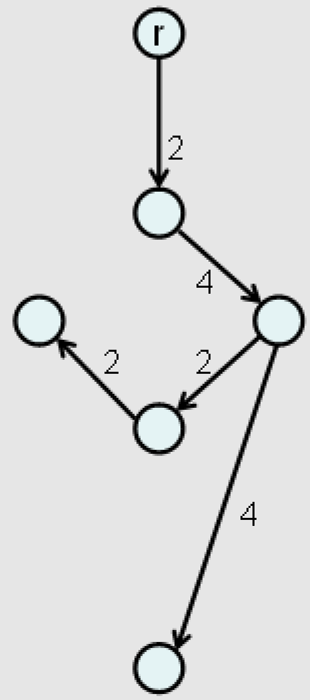
Assume all vertices reachable from r

Also called an arborescence

29

Finding a minimum branching

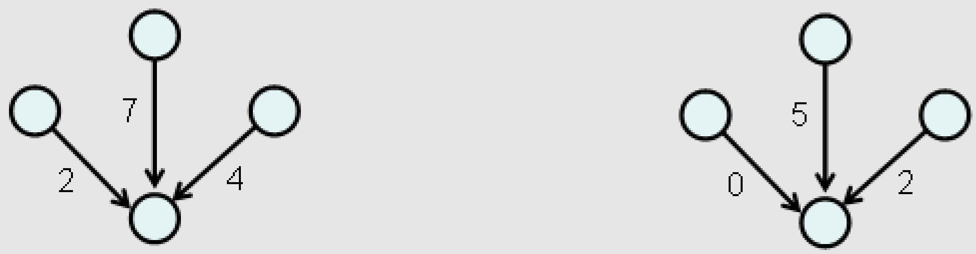




30

Finding a minimum branching

- Remove all edges going into r
- Normalize the edge weights, so the minimum weight edge coming into each vertex has weight zero



This does not change the edges of the minimum branching

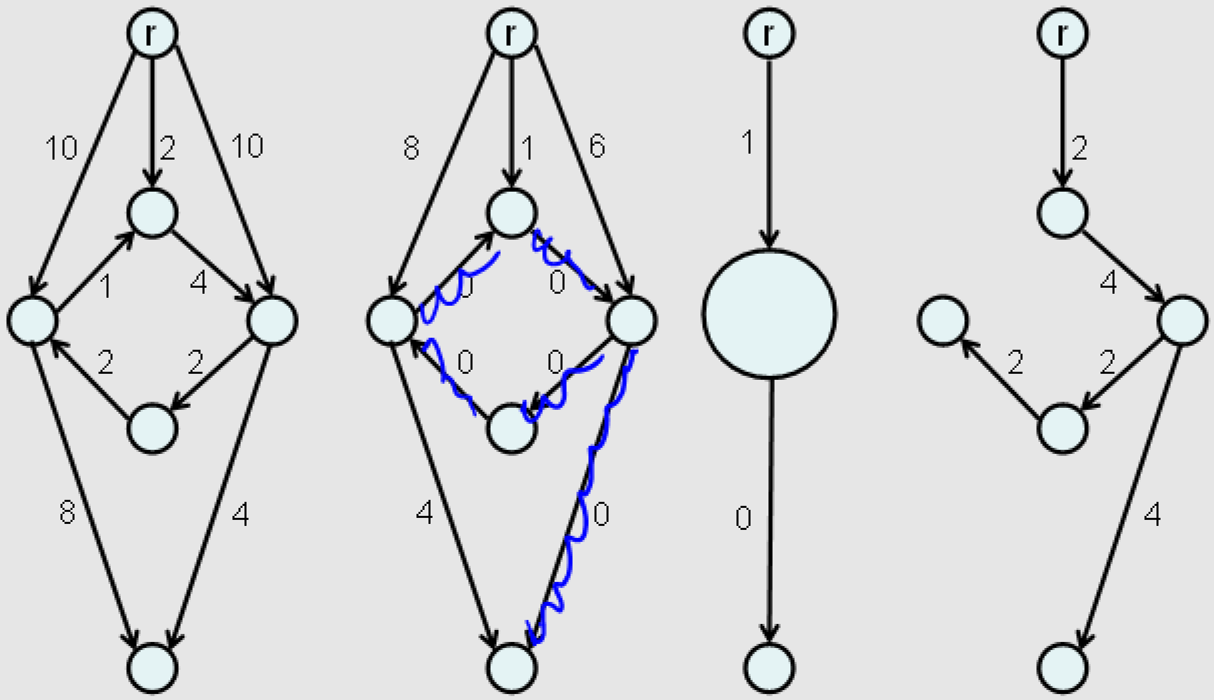
31

Finding a minimum branching

- Consider the graph that consists of the minimum cost edge coming in to each vertex
  - If this graph is a branching, then it is the minimum cost branching
  - Otherwise, the graph contains one or more cycles
    - Collapse the cycles in the original graph to super vertices
    - Reweight the graph and repeat the process

32

Finding a minimum branching

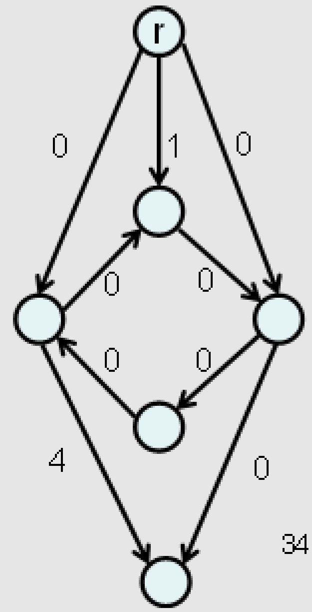


33

Correctness Proof

Lemma 4.38 Let C be a cycle in G consisting of edges of cost 0 with r not in C. There is an optimal branching rooted at r that has exactly one edge entering C.

- The lemma justifies using the edges of the cycle in the branching
- An induction argument is used to cover the multiple levels of compressing cycles



34